

ADVANCED SCIENTIFIC LIBRARY  
ASL C INTERFACE  
User's Guide  
<Shared Memory Parallel Functions>

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# PREFACE

This manual describes general concepts, functions, and specifications for use of the Advanced Scientific Library (ASL) C interface.

The manuals corresponding to this product consist of seven volumes, which are divided into the chapters shown below. This manual describes the shared memory parallel functions.

## Basic Functions Volume 1

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Storage Mode Conversion	Explanation of algorithms, method of using, and usage example of function related to storage mode conversion of array data.
3	Basic Matrix Algebra	Explanation of algorithms, method of using, and usage example of function related to basic calculations involving matrices.
4	Eigenvalues and Eigenvectors	Explanation of algorithms, method of using, and usage example of function related to <b>the standard eigenvalue problem</b> for real matrices, complex matrices, real symmetric matrices, Hermitian matrices, real symmetric band matrices, real symmetric tridiagonal matrices, real symmetric random sparse matrices, Hermitian random sparse matrices and <b>the generalized eigenvalue problem</b> for real matrices, real symmetric matrices, Hermitian matrices, real symmetric band matrices.

## Basic Functions Volume 2

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Simultaneous Linear Equations (Direct Method)	Explanation of algorithms, method of using, and usage example of function related to simultaneous linear equations corresponding to real matrices, complex matrices, positive symmetric matrices, real symmetric matrices, Hermitian matrices, real band matrices, positive symmetric band matrices, real tridiagonal matrices, real upper triangular matrices, and real lower triangular matrices.

Basic Functions Volume 3

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Fourier Transforms and their applications	Explanation of algorithms, method of using, and usage example of function related to one-, two- and three-dimensional complex Fourier transforms and real Fourier transforms, one-, two- and three-dimensional convolutions, correlations, and power spectrum analysis, wavelet transforms, and inverse Laplace transforms.

Basic Functions Volume 4

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Differential Equations and Their Applications	Explanation of algorithms, method of using, and usage example of function related to <b>ordinary differential equations initial value problems</b> for high-order simultaneous ordinary differential equations, implicit simultaneous ordinary differential equations, matrix type ordinary differential equations, stiff problem high-order simultaneous ordinary differential equations, simultaneous ordinary differential equations, first-order simultaneous ordinary differential equations, and high-order ordinary differential equations, and <b>ordinary differential equations boundary value problems</b> for high-order simultaneous ordinary differential equations, first-order simultaneous ordinary differential equations, high-order ordinary differential equations, high-order linear ordinary differential equations, and second-order linear ordinary differential equations, and <b>integral equations</b> for Fredholm's integral equations of second kind and Volterra's integral equations of first kind, and <b>partial differential equations</b> for two- and three-dimensional inhomogeneous Helmholtz equation.
3	Numerical Differentials	Explanation of algorithms, method of using, and usage example of function related to numerical differentials of one-variable functions and multi-variable functions.
4	Numerical Integration	Explanation of algorithms, method of using, and usage example of function related to numerical integration over a finite interval, semi-infinite interval, fully infinite interval, two-dimensional finite interval, and multi-dimensional finite interval.
5	Interpolations and Approximations	Explanation of algorithms, method of using, and usage example of function related to interpolations, surface interpolations, least squares approximations, least squares surface approximations, and Chebyshev's approximations.
6	Spline Functions	Explanation of algorithms, method of using, and usage example of function related to interpolation, smoothing, numerical derivatives, and numerical integrals using cubic splines, bicubic splines and B-splines.

Basic Functions Volume 5

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Special Functions	Explanation of algorithms, method of using, and usage example of function related to Bessel functions, modified Bessel functions, spherical Bessel functions, functions related to Bessel functions, Gamma functions, functions related to Gamma functions, elliptic functions, indefinite integrals of elementary functions, associated Legendre functions, orthogonal polynomials, and other special functions.
3	Sorting and Ranking	Explanation and usage examples of function related to sorting and ranking.
4	Roots of Equations	Explanation of algorithms, method of using, and usage example of function related to roots of algebraic equations, nonlinear equations, and simultaneous nonlinear equations.
5	Extremal Problems and Optimization	Explanation of algorithms, method of using, and usage example of function related to minimization of functions with no constraints, minimization of the sum of the squares of functions with no constraints, minimization of one-variable functions with constraints, minimization of multi-variable functions with constraints, and shortest path problem.

Basic Functions Volume 6

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Random Number Tests	Explanation and usage examples of function related to uniform random number tests, and distribution random number tests.
3	Probability Distributions	Explanation and usage examples of function related to continuous distributions and discrete distributions.
4	Basic Statistics	Explanation and usage examples of function related to basic statistics, variance-covariance and correlation.
5	Tests and Estimates	Explanation and usage examples of function related to interval estimates and tests.
6	Analysis of Variance and Design of Experiments	Explanation and usage examples of function related to one-way layout, two-way layout, multiple-way layout, randomized block design, Greco-Latin square method, cumulative Method.
7	Nonparametric Tests	Explanation and usage examples of function related to tests using $\chi^2$ distribution and tests using other distributions.
8	Multivariate Analysis	Explanation and usage examples of function related to principal component analysis, factor analysis, canonical correlation analysis, discriminant analysis, cluster analysis.
9	Time Series Analysis	Explanation and usage examples of function related to autocorrelation, cross correlation, autocovariance, cross covariance, smoothing and demand forecasting.
10	Regression analysis	Explanation and usage examples of function related to linear Regression and nonlinear Regression.

## Shared Memory Parallel Functions

Chapter	Title	Contents
1	Introduction	Explanation of the organization of this manual, how to view each item, and usage limitations.
2	Basic Matrix Algebra	Explanation of algorithms, method of using, and usage example of function related to obtain the product of real matrices and complex matrices.
3	Simultaneous Linear Equations (Direct Method)	Explanation of algorithms, method of using, and usage example of function related to simultaneous linear equations corresponding to real matrices, complex matrices, real symmetric matrices, and Hermitian matrices.
4	Simultaneous Linear Equations (Iteration Method)	Explanation of algorithms, method of using, and usage example of function related to simultaneous linear equations corresponding to real positive definite symmetric sparse matrices, real symmetric sparse matrices and real asymmetric sparse matrices.
5	Eigenvalues and Eigenvectors	Explanation of algorithms, method of using, and usage example of function related to the eigenvalue problem for real symmetric matrices and Hermitian matrices.
6	Fourier Transforms and their applications	Explanation of algorithms, method of using, and usage example of function related to one-, two- and three-dimensional complex Fourier transforms and real Fourier transforms, two- and three-dimensional convolutions, correlations, and power spectrum analysis.
7	Sorting	Explanation and usage examples of function related to sorting and ranking.

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### Remarks

- (1) This manual corresponds to ASL 1.1. All functions described in this manual are program products.
- (2) Proper nouns such as product names are registered trademarks or trademarks of individual manufacturers.
- (3) This library was developed by incorporating the latest numerical computational techniques. Therefore, to keep up with the latest techniques, if a newly added or improved function includes the function of an existing function may be removed.

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# Chapter 1

---

## INTRODUCTION

### 1.1 OVERVIEW

#### 1.1.1 Introduction to The Advanced Scientific Library ASL C interface

Table 1–1 lists correspondences among product categories, functions of ASL and supported hardware platforms. Interfaces of those functions that have the same name and that belong to the same version of ASL are common among hardware platforms.

Table 1–1 Classification of functions included in ASL

Classification of Functions	Volume
Basic functions	Vol. 1-6
Shared memory parallel functions	Vol. 7

#### 1.1.2 Distinctive Characteristics of ASL C interface

ASL C interface has the following distinctive characteristics.

- (1) Functions are optimized using compiler optimization to take advantage of corresponding system hardware features.
- (2) Special-purpose functions for handling matrices are provided so that the optimum processing can be performed according to the type of matrix (symmetric matrix, Hermitian matrix, or the like). Generally, processing performance can be increased and the amount of required memory can be conserved by using the special-purpose functions.
- (3) Functions are modularized according to processing procedures to improve reliability of each component function as well as the reliability and efficiency of the entire system.
- (4) Error information is easy to access after a function has been used since error indicator numbers have been systematically determined.

---

## 1.2 KINDS OF LIBRARIES

Numeric storage units of ASL C interface is 4-byte.

Table 1–2 Kinds of libraries providing ASL C interface

Size of variable(byte)		Declaration of arguments	Kind	Kind of library
integer	real			
4	8	int double	32bit integer Double-precision function	32bit integer library (link option: -lasl_openmp)
4	4	int float	32bit integer Single-precision function	
8	8	long double	64bit integer Double-precision function	64bit integer library (link option: -lasl_openmp_i64)
8	4	long float	64bit integer Single-precision function	

(\*1) Functions that appear in this documentation do not always support all of the four kinds of functions listed above. For those functions that do not support some of those function kinds, relevant notes will appear in the corresponding subsections.

(\*2) For compiling the program with functions in the 64-bit integer library, the option “-DASL\_LIB\_INT64” must be specified (See the Note (2) in 1.6).

---

## 1.3 ORGANIZATION

This section describes the organization of Chapters 2 and later.

### 1.3.1 Introduction

The first section of each chapter is a general introduction describing such information as the effective ways of using the functions, techniques employed, algorithms on which the functions are based, and notes.

### 1.3.2 Organization of Function Description

The second section of each chapter sequentially describes the following topics for each function.

- (1) Function
- (2) Usage
- (3) Arguments and return value
- (4) Restrictions
- (5) Error indicator (Return Value)
- (6) Notes
- (7) Example

Each item is described according to the following principles.

### 1.3.3 Contents of Each Item

(1) **Function**

Function briefly describes the purpose of the ASL C interface function.

(2) **Usage**

Usage describes the function name and the order of its arguments. In general, arguments are arranged as follows. When an argument is an address-passing variable, & is appended in front of the argument name.

```
ierr = function-name (input-arguments, input/output-arguments, output-arguments, isw, work);
```

isw is an input argument for specifying the processing procedure. ierr is a return value. In some cases, input/output arguments precede input arguments. The following general principles also apply.

- Array are placed as far to the left as possible according to their importance.
- The dimension of an array immediately follows the array name. If multiple arrays have the same dimension, the dimension is assigned as an argument of only the first array name. It is not assigned as an argument of subsequent array names.

(3) **Arguments and return value**

Arguments and return value are explained in the order described above in paragraph (2). The explanation format is as follows.

<u>Arguments and return value</u>	<u>Type</u>	<u>Size</u>	<u>Input/Output</u>	<u>Contents</u>
(a)	(b)	(c)	(d)	(e)



(a) Arguments and return value

Arguments and return value are explained in the order they are designated in the Usage paragraph.

(b) Type

Type indicates the data type of the argument. Any of the following codes may appear as the type.

**I** : Integer type

**D** : Double precision real

**R** : Real

**Z** : Double precision complex

**C** : Complex

There are 64-bit integer and 32-bit integer for integer type arguments. In a 32-bit (64-bit) integer type function, all the integer type arguments are 32-bit (64-bit) integer. In other words, kinds of libraries determine the sizes of integer type arguments (Refer to 1.4). In the user program, a 32-bit/64-bit integer type argument must be declared by `int`/`long`, respectively.

(c) Size

Size indicates the required size of the specified argument. If the size is greater than 1, the required area must be reserved in the program calling this function.

**1** : Indicates that argument is a variable.

**n** : Indicates that the argument is a vector (one-dimensional array) having `n` elements. The argument `n` indicating the size of this vector is defined immediately after the specified vector. However, if the size of a vector or array defined earlier, it is omitted following subsequently defined vectors or arrays. The size may be specified by only a numeric value or in the form of a product or sum such as  $3 \times n$  or  $n + m$ .

(d) Input/Output

Input/Output indicates whether the explanation of argument contents applies to input time or output time.

i. When only “Input” appears

When the control returns to the program using this function, information when the argument is input is preserved. The user must assign input-time information unless specifically instructed otherwise. When the argument is a variable, the variable value must be passed.

ii. When only “Output” appears

Results calculated within the function are output to the argument. No data is entered at input time. When the argument is a variable, the variable address must be passed.

iii. When both “Input” and “Output” appear

Argument contents change between the time control passes to the function and the time control returns from the function. The user must assign input-time information unless specifically instructed otherwise. When the argument is a variable, the variable address must be passed.

iv. When “Work” appears

Work indicates that the argument is an area used when performing calculations within the function. A work area having the specified size must be reserved in the program calling this function. The contents of the work area may have to be maintained so they can be passed along to the next calculation.

(e) Contents

Contents describes information held by the argument at input time or output time.

- A sample Argument description follows.

**Example**

The statement of the function (ASL\_dbgmlc, ASL\_rbgmlc) that obtains the LU decomposition and the condition number of a real matrix is as follows.

Double precision:

```
ierr = ASL_dbgmlc (a, lna, n, ipvt, &cond, w1);
```

Single precision:

```
ierr = ASL_rbgmlc (a, lna, n, ipvt, &cond, w1);
```

The explanation of the arguments and return value is as follows.

Table 1–3 Sample Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	Note $\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna×n	Input	Real matrix <i>A</i> (two-dimensional array)
				Output	The matrix <i>A</i> decomposed into the matrix <i>LU</i> where <i>U</i> is a unit upper triangular matrix and <i>L</i> is a lower triangular matrix.
2	lna	I	1	Input	Adjustable dimension size of array a
3	n	I	1	Input	Order <i>n</i> of matrix <i>A</i>
4	ipvt	I*	n	Output	Pivoting information ipvt[ <i>i</i> −1]: Number of the row exchanged with row <i>i</i> in the <i>i</i> -th step.
5	cond	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Output	Reciprocal of the condition number
6	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
7	ierr	I	1	Output	Error indicator (Return Value)

To use this function, arrays a, ipvt and w1 must first be allocated in the calling program so they can be used as arguments. a is a  $\begin{cases} \text{double-precision} \\ \text{single-precision} \end{cases}$  <sup>Note</sup> real array of size [lna × n], ipvt is an integer

array of size n and w1 is a  $\begin{cases} \text{double-precision} \\ \text{single-precision} \end{cases}$  real array of size n.

When the 64-bit integer version is used, all integer-type arguments (lna, n, ipvt and ierr) must be declared by using long, not int.

**Note** The entries enclosed in brace { } mean that the array should be declared double precision type when using function ASL\_dbgmlc and real type when using function ASL\_rbgmlc. Braces are used in this manner throughout the remainder of the text unless specifically stated otherwise.

Data must be stored in `a`, `lna` and `n` before this function is called. The LU decomposition and condition number of the assigned matrix are calculated with in the function, and the results are stored in array `a` and variable `cond`. In addition, pivoting information is stored in `ipvt` for use by subsequent functions.

`ierr` is a return value used to notify the user of invalid input data or an error that may occur during processing. If processing terminates normally, `ierr` is set to zero.

Since `w1` is a work area used only within the function, its contents at input and output time have no special meaning.

(4) **Restrictions**

Restrictions indicate limiting ranges for function arguments.

(5) **Error indicator (Return Value)**

Each function has been given an error indicator as a return value. This error indicator, which has uniformly been given the variable name `ierr`, is placed at the end of the arguments. If an error is detected within the function, a corresponding value is output to `ierr`. Error indicator values are divided into five levels.

Table 1–4 Classification of Return Values

Level	Return value	Meaning	Processing result
Normal	0	Processing is terminated normally.	Results are guaranteed.
Warning	1000~2999	Processing is terminated under certain conditions.	Results are conditionally guaranteed.
Fatal	3000~3499	Processing is aborted since an argument violated its restrictions.	Results are not guaranteed.
	3500~3999	Obtained results did not satisfy a certain condition.	Obtained results are returned (the results are not guaranteed).
	4000 or more	A fatal error was detected during processing. Usually, processing is aborted.	Results are not guaranteed.

(6) **Notes**

Notes describes ambiguous items and points requiring special attention when using the function.

(7) **Example**

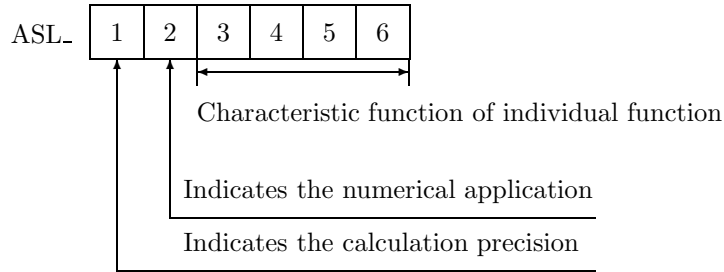
Here gives an example of how to use the function. Note that in some cases, multiple functions are combined in a single example. The output results are given in the 32-bit integer version, and may differ within the range of rounding error if the compiler or intrinsic functions are different.

In addition, when the 64-bit integer version library is used, the `long`-type conversion specification to be given to `printf` or `scanf` must be `%ld`. The source codes of examples in this document are included in User’s Guide. Input data, if required, is also included in it. To build up an executable files by compiling these example source codes, they should be linked with this product library.

## 1.4 FUNCTION NAMES

The functions name of ASL C interface shared memory parallel functions consists of ten characters with a prefix “ASL\_” and (six alphanumeric characters).

Figure 1–1 Function Name Components



**“1” in Figure 1–1 :** The following four letters are used to indicate the calculation precision.

- q Parallel function double precision real-type calculation
- p Parallel function single precision real-type calculation
- h Parallel function double precision complex-type calculation
- g Parallel function single precision complex-type calculation

In the Basic Functions Volumes 1 to 6, the following eight letters are used to indicate the calculation precision.

- d, w Double precision real-type calculation
- r, v Single precision real-type calculation
- z, j Double precision complex-type calculation
- c, i Single precision complex-type calculation

However, the complex type calculations listed above do not necessarily require complex arguments.

**“2” in Figure 1–1 :** Currently, the following letters letterererere are used to indicate the application field in the ASL C interface related products.

Letter	Application Field	Volume
a	Storage mode conversion	1
	Basic matrix algebra	1, 7
b	Simultaneous linear equations (direct method)	2, 7
c	Eigenvalues and eigenvectors	1, 7
f	Fourier transforms and their applications	3, 7
	Time series analysis	6
g	Spline function	4
h	Numeric integration	4
i	Special function	5
j	Random number tests	6

---

Letter	Application Field	Volume
k	Ordinary differential equation (initial value problems)	4
l	Roots of equations	5
m	Extremum problems and optimization	5
n	Approximation and regression analysis	4, 6
o	Ordinary differential equations (boundary value problems), integral equations and partial differential equations	4
p	Interpolation	4
q	Numerical differentials	4
s	Sorting and ranking	5, 7
x	Basic matrix algebra	1
	Simultaneous linear equations (iterative method)	7
1	Probability distributions	6
2	Basic statics	6
3	Tests and estimates	6
4	Analysis of variance and design of experiments	6
5	Nonparametric tests	6
6	Multivariate analysis	6

**“3–6” in Figure 1–1 :** These characters indicate the characteristic function of the individual function.

---

## 1.5 ASL C INTERFACE SHARED MEMORY PARALLEL FUNCTIONS

### 1.5.1 Overview of Shared Memory Parallel Functions

All the ASL C interface shared memory parallel functions are developed using OpenMP for a shared memory, multi-thread system model. Shared memory parallel functions run across plural processors. As a result, we may expect that an elapsed time will reduce which is required for analysis of large-scale problems.

There are two possible ways to run a program which uses ASL C interface in shared memory parallel as follows:

- (1) Use ASL C interface shared memory parallel functions from the user program. In this case, ASL C interface divides up and allocates internal processing among plural processors so that the allocated processing is executed in shared memory parallel.
- (2) Parallelize the user program itself by using multi-thread functions, and use either user functions and ASL C interface functions or plural ASL C interface functions in parallel. In this case, the processing of functions that are called in parallel will be allocated among plural processors and be executed in parallel.

For the method (2), use ASL C interface non-parallel functions (basic functions).

### 1.5.2 Performance Improvement Due to Parallel Functions

General conditions for improving program performance by using ASL C interface shared memory parallel functions are described below.

- (1) Execution environment

Scalability is nothing but an actual reduction in elapsed time, which is achieved through simultaneous use of plural processors. So you can not expect the scalability of shared memory parallel program when all the processors in your parallel system are busy most of time with many jobs.

- (2) Problem scale

Shared memory parallel processing functions require more overhead for dividing and synchronizing processing than the overhead required when shared memory parallel processing is not performed. Therefore, for small-sized problems, this overhead may surpass the time reduction achieved by parallel processing. In general, the larger the problem size, the greater the effectiveness of shared memory parallel processing since the influence of the surplus overhead becomes relatively small compared to the total processing time. Therefore, to obtain the benefits of shared memory parallel processing, the problem scale should be large.

### 1.5.3 General Notes Concerning the Use of shared memory Parallel Functions

- (1) Execution environment

Shared memory parallel functions can be used only with an operating system associated with a multi-core system. For a detailed explanation of the shared memory parallel processing (multi-thread), refer to the corresponding compiler manuals.

- (2) Thread parameter (nt)

Shared memory parallel function functions have the number of tasks parameter `nt` as an argument. This parameter specifies the number of subdivisions of the processing when parallelization is performed within ASL C interface.

(3) Setting the number of processor cores which are used

You can use the `OMP_NUM_THREADS` environment variable to control the number of processor cores which are used. The following command lines show the C shell syntax and Bourne shell syntax to use when setting the variable to 2 processor cores.

- C shell:

```
setenv OMP_NUM_THREADS 2 RETURN
```

- Bourne shell:

```
OMP_NUM_THREADS=2 RETURN  
export OMP_NUM_THREADS RETURN
```

---

## 1.6 NOTES

- (1) To use ASL C interface, the header file `asl.h` must be included.
- (2) For compiling the program with functions in ASL C interface 64-bit integer library, the compile option “`-DASL_LIB_INT64`” must be specified. This option will activate the prototype declaration for 64-bit integer functions in the header file `asl.h`, and without the option “`-DASL_LIB_INT64`”, those for 32-bit integer functions will be activated.
- (3) The name “(6 lowercase letters) following `ASL_`” is reserved by ASL C interface.
- (4) For using 64-bit integer library, you must use “`long`” for integer type declaration. Otherwise, use “`int`” for integer type declaration.
- (5) Use the functions of double precision version whenever possible. They not only provide higher precision solutions but also are more stable than single precision versions, in particular, for eigenvalue and eigenvector problems.
- (6) To suppress compiler operation exceptions, ASL C interface functions are set to so that they conform to the compiler parameter indications of a user’s main program. Therefore, the main program must suppress any operation exceptions.
- (7) The numerical calculation programs generally deal with operations on finite numbers of digits, so the precision of the results cannot exceed the number of operation digits being handled. For example, since the number of operation digits (in the mantissa part) for double-precision operations is on the order of 15 decimal digits, when using these floating point modes to calculate a value that mathematically becomes 1, an error on the order of  $10^{-15}$  may be introduced at any time. Of course, if multiple length arithmetic is emulated such as when performing operations on an arbitrary number of digits, this kind of error can be controlled. However, in this case, when constants such as  $\pi$  or function approximation constants, which are fixed in double-precision operations, for example, are also to be subject to calculations that depend on the length of the multiple length arithmetic operations, the calculation efficiency will be worse than for normal operations.
- (8) A solution cannot be obtained for a problem for which no solution exists mathematically. For example, a solution of simultaneous linear equations having a singular (or nearly singular) matrix for its coefficient matrix theoretically cannot be obtained with good precision mathematically. Numerical calculations cannot strictly distinguish between mathematically singular and nearly singular matrices. Of course, it is always possible to consider a matrix to be singular if the calculation value for the condition number is greater than or equal to an established criterion value.
- (9) Generally, if data is assigned that causes a floating point exception during calculations (such as a floating point overflow), a normal calculation result cannot be expected. However, a floating point underflow that occurs when adding residuals in an iterative calculation is an exception to this.
- (10) For problems that are handled using numerical calculations (specifically, problems that use iterative techniques as the calculation method), there are cases in which a solution cannot be obtained with good precision and cases in which no solution can be obtained at all, by a special-purpose function.
- (11) Depending on the problem being dealt with, there may be cases when there are multiple solutions, and the execution result differs in appearance according to the compiler used or the computer or OS under which



---

the program is executed. For example, when an eigenvalue problem is solved, the eigenvectors that are obtained may differ in appearance in this way.

- (12) The mark “DEPRECATED” denotes that the subroutine will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative practice instead.



## 2.2 BASIC MATRIX ALGEBRA

### 2.2.1 ASL\_qam1mu, ASL\_pam1mu

#### Multiplying Real Matrices (Two-Dimensional Array Type)

(1) **Function**

Obtain the product of two real matrices  $A$  and  $B$  (two-dimensional array type).

(2) **Usage**

Double precision:

ierr = ASL\_qam1mu (a, lma, nm, nn, b, lnb, nl, c, lmc, nt);

Single precision:

ierr = ASL\_pam1mu (a, lma, nm, nn, b, lnb, nl, c, lmc, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	lma×nn	Input	Real matrix $A$ (two-dimensional array type).
2	lma	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ (Number of rows in matrix $C$ ).
4	nn	I	1	Input	Number of columns in matrix $A$ (Number of rows in matrix $B$ ).
5	b	$\begin{cases} D* \\ R* \end{cases}$	lnb×nl	Input	Real matrix $B$ (two-dimensional array type).
6	lnb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ (Number of columns in matrix $C$ ).
8	c	$\begin{cases} D* \\ R* \end{cases}$	lmc×nl	Output	Product $A \cdot B$ of matrices $A, B$ (two-dimensional array type).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	nt	I	1	Input	Number of tasks to be generated.
11	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < mn \leq lnb$
- (c)  $nl > 0$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$mn = 1$	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	

(6) **Notes**

None

(7) **Example**

(a) Problem

$$A = \begin{bmatrix} 1 & 2 & 0 & -1 \\ -3 & -5 & 1 & 2 \\ 1 & 3 & 2 & -2 \\ 0 & 2 & 1 & -1 \end{bmatrix}$$

$$B = \begin{bmatrix} -3 & -1 & 1 & -1 \\ -3 & -1 & 0 & 1 \\ -4 & -1 & 1 & 0 \\ -10 & -3 & 1 & 1 \end{bmatrix}$$

Obtain  $C = AB$ .

(b) Input data

Matrix  $A$ , matrix  $B$ ,  $lma = 11$ ,  $lnb = 11$ ,  $lmc = 11$ ,  $nm = 4$ ,  $mn = 4$ ,  $nl = 4$  and  $nt = 2$ .

(c) Main program

```

/*      C interface example for ASL_qam1mu */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int ma=4;
    int mm=4;
    int nn=4;
    double *b;
    int nb=4;
    int nt=2;
    int ll=4;
    double *c;
    int mc=4;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "qam1mu.dat", "r" );
    if( fp == NULL )
    {

```

```

        printf( "file open error\n" );
        return -1;
    }

    printf( "    *** ASL_qam1mu ***\n" );
    printf( "\n    ** Input **\n\n" );

    a = ( double * )malloc((size_t)( sizeof(double) * (ma*nn) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    b = ( double * )malloc((size_t)( sizeof(double) * (nb*nn) ));
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    c = ( double * )malloc((size_t)( sizeof(double) * (mc*nn) ));
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    printf( "\tMatrix a\n\n" );
    for( i=0 ; i<ma ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            fscanf( fp, "%lf", &a[i+ma*j] );
            printf( "%8.3g ", a[i+ma*j] );
        }
        printf( "\n" );
    }

    printf( "\n\tMatrix b\n\n" );
    for( i=0 ; i<nb ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            fscanf( fp, "%lf", &b[i+nb*j] );
            printf( "%8.3g ", b[i+nb*j] );
        }
        printf( "\n" );
    }

    fclose( fp );

    ierr = ASL_qam1mu(a, ma, mm, nn, b, nb, ll, c, mc, nt);

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\n\tMatrix c\n\n" );
    for( i=0 ; i<mc ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "%8.3g ", c[i+mc*j] );
        }
        printf( "\n" );
    }

    free( a );
    free( b );
    free( c );

    return 0;
}

```

(d) Output results

```

*** ASL_qam1mu ***

** Input **

Matrix a
    1      2      0      -1
   -3     -5      1      2
    1      3      2     -2

```

```
      0      2      1     -1
Matrix b
     -3     -1      1     -1
     -3     -1      0      1
     -4     -1      1      0
    -10     -3      1      1

** Output **
ierr =      0
Matrix c
      1      0      0      0
      0      1      0      0
      0      0      1      0
      0      0      0      1
```

**2.2.2 ASL\_qam1mm, ASL\_pam1mm****Multiplying Real Matrices (Two-Dimensional Array Type) ( $C = C \pm AB$ )****(1) Function**

Obtain the product of real matrix  $A$  and real matrix  $B$  ( $C = C \pm AB$ ).

**(2) Usage**

Double precision:

ierr = ASL\_qam1mm (a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_pam1mm (a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

**(3) Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Real matrix $A$ (two-dimensional array type).
2	lma	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Real matrix $B$ (two-dimensional array type).
6	lnb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Initial real matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm]AB$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. isw = 1: Obtain $C = C + AB$ isw = 0: Obtain $C = AB$ isw = -1: Obtain $C = C - AB$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

**(4) Restrictions**

(a)  $0 < nm \leq lma, lmc$

(b)  $0 < nm \leq lnb$

(c)  $nl > 0$

(d)  $isw \in \{0, 1, -1\}$

(e)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) Notes

None

(7) Example

(a) Problem

$$A = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 2 & 2 & 2 \\ 3 & 3 & 3 & 3 & 3 \\ 4 & 4 & 4 & 4 & 4 \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 2 & 2 & 2 & 2 \\ 3 & 3 & 3 & 3 & 3 & 3 \\ 4 & 4 & 4 & 4 & 4 & 4 \\ 5 & 5 & 5 & 5 & 5 & 5 \end{bmatrix}$$

Obtain  $C = AB$ .

(b) Input data

Matrices  $A$  and  $B$ ,  $lma = 11$ ,  $lmb = 11$ ,  $lmc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$ ,  $isw = 0$  and  $nt = 2$ .

(c) Main program

```

/*      C interface example for ASL_qam1mm */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lma=11;
    int nm=4;
    int nn=5;
    double *b;
    int lnb=11;
    int nl=6;
    double *c;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_qam1mm ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n\n", lma, lnb, lmc );
    printf( "\tnm  =%2d  nn  =%2d  nl  =%2d\n\n", nm,  nn,  nl  );
    printf( "\tisw=%2d  nt  =%2d\n\n", isw, nt );

    a = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }
}

```



```

}
b = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lnb * nl ) ) );
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

c = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lmc * nl ) ) );
if( c == NULL )
{
    printf( "no enough memory for array c\n" );
    return -1;
}

printf( "\tMatrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", a[i+lma*j]=i+1 );
    }
    printf( "\n" );
}
printf( "\n\tMatrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", b[i+lmb*j]=i+1 );
    }
    printf( "\n" );
}

ierr = ASL_qam1mm( a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt );

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tMatrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", c[i+lmc*j] );
    }
    printf( "\n" );
}

free( a );
free( b );
free( c );

return 0;
}

```

(d) Output results

\*\*\* ASL\_qam1mm \*\*\*

\*\* Input \*\*

lma=11 lnb=11 lmc=11

nm = 4 nn = 5 nl = 6

isw= 0 nt = 2

Matrix A

1	1	1	1	1
2	2	2	2	2
3	3	3	3	3
4	4	4	4	4

Matrix B

1	1	1	1	1	1
2	2	2	2	2	2
3	3	3	3	3	3
4	4	4	4	4	4
5	5	5	5	5	5

\*\* Output \*\*

```
ierr =      0
Matrix C
  15      15      15      15      15      15
  30      30      30      30      30      30
  45      45      45      45      45      45
  60      60      60      60      60      60
```

### 2.2.3 ASL\_qam1mt, ASL\_pam1mt

#### Multiplying Real Matrices (Two-Dimensional Array Type) ( $C = C \pm AB^T$ )

(1) **Function**

Obtain the product of real matrix  $A$  and real matrix  $B$  ( $C = [C \pm]AB^T$ )

(2) **Usage**

Double precision:

ierr = ASL\_qam1mt (a, lma, nm, nn, b, llb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_pam1mt (a, lma, nm, nn, b, llb, nl, c, lmc, isw, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Real matrix $A$ (two-dimensional array type).
2	lma	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	$llb \times nn$	Input	Real transposed matrix $B$ (two-dimensional array type).
6	llb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Initial real matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm]AB^T$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + AB^T$ $isw = 0$ : Obtain $C = AB^T$ $isw = -1$ : Obtain $C = C - AB^T$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < nl \leq llb$
- (c)  $nn > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) **Notes**

None

(7) **Example**

(a) Problem

$$A = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 2 & 2 & 2 \\ 3 & 3 & 3 & 3 & 3 \\ 4 & 4 & 4 & 4 & 4 \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \end{bmatrix}$$

Obtain  $C = AB^T$ .

(b) Input data

Matrices  $A$  and  $B$ ,  $lma = 11$ ,  $llb = 11$ ,  $lmc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$ ,  $isw = 0$  and  $nt = 2$ .

(c) Main program

```

/*      C interface example for ASL_qam1mt */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lma=11;
    int nm=4;
    int nn=5;
    double *b;
    int llb=11;
    int nl=5;
    double *c;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_qam1mt ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d  llb=%2d  lmc=%2d\n\n", lma, llb, lmc );
    printf( "\tnm  =%2d  nn  =%2d  nl  =%2d\n\n", nm,  nn,  nl  );
    printf( "\tisw=%2d  nt  =%2d\n\n", isw, nt );

    a = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

```

```

    }
    b = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( llb * nn ) ) );
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    c = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lmc * nl ) ) );
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    printf( "\tMatrix A\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "%8.3g ", a[i+lma*j]=i+1 );
        }
        printf( "\n" );
    }
    printf( "\n\tMatrix B(Transposed Storage)\n" );
    for( i=0 ; i<nn ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", b[j+llb*i]=j+1 );
        }
        printf( "\n" );
    }

    ierr = ASL_qam1mt( a, lma, nm, nn, b, llb, nl, c, lmc, isw, nt );

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\n\tMatrix C\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", c[i+lmc*j] );
        }
        printf( "\n" );
    }

    free( a );
    free( b );
    free( c );

    return 0;
}

```

(d) Output results

```

*** ASL_qam1mt ***

** Input **

lma=11  llb=11  lmc=11
nm = 4  nn = 5  nl = 5
isw= 0  nt = 2

Matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4

Matrix B(Transposed Storage)
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

** Output **

```

```
ierr =      0
Matrix C
   5      10      15      20      25
  10      20      30      40      50
  15      30      45      60      75
  20      40      60      80     100
```

## 2.2.4 ASL\_qam1tm, ASL\_pam1tm

### Multiplying Real Matrices (Two-Dimensional Array Type) ( $C = C \pm A^T B$ )

#### (1) Function

Obtain the product of real matrix  $A$  and real matrix  $B$  ( $C = [C \pm] A^T B$ )

#### (2) Usage

Double precision:

ierr = ASL\_qam1tm (a, lna, nm, nn, b, lnb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_pam1tm (a, lna, nm, nn, b, lnb, nl, c, lmc, isw, nt);

#### (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	lna×nm	Input	Real transposed matrix $A$ (two-dimensional array type).
2	lna	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	lnb×nl	Input	Real matrix $B$ (two-dimensional array type).
6	lnb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	lmc×nl	Input	Initial real matrix $C$ (If isw = ±1) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm] A^T B$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. isw = 1: Obtain $C = C + A^T B$ isw = 0: Obtain $C = A^T B$ isw = -1: Obtain $C = C - A^T B$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

#### (4) Restrictions

- (a)  $0 < nm \leq lmc$
- (b)  $0 < nn \leq lna, lnb$
- (c)  $nl > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) **Notes**

None

(7) **Example**

(a) Problem

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 1 & 1 & 1 & 1 & 1 \\ 2 & 2 & 2 & 2 & 2 & 2 \\ 3 & 3 & 3 & 3 & 3 & 3 \\ 4 & 4 & 4 & 4 & 4 & 4 \\ 5 & 5 & 5 & 5 & 5 & 5 \end{bmatrix}$$

Obtain  $C = A^T B$ .

(b) Input data

Matrices  $A$  and  $B$ ,  $lna = 11$ ,  $lnb = 11$ ,  $lnc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$ ,  $isw = 0$  and  $nt = 2$ .

(c) Main program

```

/*      C interface example for ASL_qam1tm */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lna=11;
    int nm=5;
    int nn=5;
    double *b;
    int lnb=11;
    int nl=4;
    double *c;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_qam1tm ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlna=%2d  lnb=%2d  lmc=%2d\n\n", lna, lnb, lmc );
    printf( "\tnm  =%2d  nn  =%2d  nl  =%2d\n\n", nm,  nn,  nl  );
    printf( "\tisw=%2d  nt  =%2d\n\n", isw, nt );

    a = ( double * )malloc((size_t)( sizeof(double) * (lna*nm) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }
}

```



```

    }
    b = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lnb * nl ) ) );
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    c = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lmc * nl ) ) );
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    printf( "\tMatrix A(Transposed Storage)\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "%8.3g ", a[ j+lna*i ] = j+1 );
        }
        printf( "\n" );
    }
    printf( "\n\tMatrix B\n" );
    for( i=0 ; i<nn ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", b[ i+lnb*j ] = i+1 );
        }
        printf( "\n" );
    }

    ierr = ASL_qam1tm( a, lna, nm, nn, b, lnb, nl, c, lmc, isw, nt );

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\n\tMatrix C\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", c[ i+lmc*j ] );
        }
        printf( "\n" );
    }

    free( a );
    free( b );
    free( c );

    return 0;
}

```

(d) Output results

```

*** ASL_qam1tm ***

** Input **

lna=11  lnb=11  lmc=11
nm = 5  nn = 5  nl = 4
isw= 0  nt = 2

Matrix A(Transposed Storage)
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Matrix B
  1      1      1      1
  2      2      2      2
  3      3      3      3
  4      4      4      4
  5      5      5      5

** Output **

```

```
ierr =      0
Matrix C
  55      55      55      55
  55      55      55      55
  55      55      55      55
  55      55      55      55
  55      55      55      55
```

### 2.2.5 ASL\_qam1tt, ASL\_pam1tt

#### Multiplying Real Matrices (Two-Dimensional Array Type) ( $C = C \pm A^T B^T$ )

(1) **Function**

Obtain the product of real matrix  $A$  and real matrix  $B$  ( $C = [C \pm] A^T B^T$ )

(2) **Usage**

Double precision:

ierr = ASL\_qam1tt (a, lna, nm, nn, b, llb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_pam1tt (a, lna, nm, nn, b, llb, nl, c, lmc, isw, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	lna×nm	Input	Real transposed matrix $A$ (two-dimensional array type).
2	lna	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	llb×nn	Input	Real transposed matrix $B$ (two-dimensional array type).
6	llb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	lmc×nl	Input	Initial real matrix $C$ (If isw = ±1) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm] A^T B^T$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. isw = 1: Obtain $C = C + A^T B^T$ isw = 0: Obtain $C = A^T B^T$ isw = -1: Obtain $C = C - A^T B^T$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < nm \leq lmc$
- (b)  $0 < nn \leq lna$
- (c)  $0 < nl \leq llb$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) **Notes**

None

(7) **Example**

(a) Problem

$$A = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \\ 1 & 2 & 3 & 4 & 5 \end{bmatrix}$$

$$B = \begin{bmatrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \\ 1 & 2 & 3 & 4 & 5 & 6 \end{bmatrix}$$

Obtain  $C = A^T B^T$ .

(b) Input data

Matrices  $A$  and  $B$ ,  $l_{na} = 11$ ,  $l_{lb} = 11$ ,  $l_{nc} = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$ ,  $isw = 0$  and  $nt = 2$ .

(c) Main program

```

/*      C interface example for ASL_qam1tt */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lna=11;
    int nm=5;
    int nn=5;
    double *b;
    int llb=11;
    int nl=4;
    double *c;
    int lcm=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_qam1tt ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlna=%2d  llb=%2d  lcm=%2d\n\n", lna, llb, lcm );
    printf( "\tnm =%2d  nn =%2d  nl =%2d\n\n", nm, nn, nl );
    printf( "\tisw=%2d  nt =%2d\n\n", isw, nt );

    a = ( double * )malloc((size_t)( sizeof(double) * (lna*nm) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }
}

```

```

    }
    b = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( llb * nn ) ) );
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    c = ( double * ) malloc( ( size_t ) ( sizeof( double ) * ( lmc * nl ) ) );
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    printf( "\tMatrix A(Transposed Storage)\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "%8.3g ", a[ j+lna*i ] = j+1 );
        }
        printf( "\n" );
    }
    printf( "\n\tMatrix B(Transposed Storage)\n" );
    for( i=0 ; i<nn ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", b[ j+llb*i ] = j+1 );
        }
        printf( "\n" );
    }

    ierr = ASL_qam1tt( a, lna, nm, nn, b, llb, nl, c, lmc, isw, nt );

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\n\tMatrix C\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", c[ i+lmc*j ] );
        }
        printf( "\n" );
    }

    free( a );
    free( b );
    free( c );

    return 0;
}

```

(d) Output results

```

*** ASL_qam1tt ***

** Input **

lna=11  llb=11  lmc=11
nm = 5  nn = 5  nl = 4
isw= 0  nt = 2

Matrix A(Transposed Storage)
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Matrix B(Transposed Storage)
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4

** Output **

```

```
ierr =      0
Matrix C
  15      30      45      60
  15      30      45      60
  15      30      45      60
  15      30      45      60
  15      30      45      60
```

## 2.2.6 ASL\_ham1mm, ASL\_gam1mm

### Multiplying Complex Matrices (Two-Dimensional Array Type) (Real Argument Type) ( $C = C \pm AB$ )

(1) **Function**

Obtain the product of two complex matrices (Two-dimensional Array Type) ( $C = [C \pm]AB$ ).

(2) **Usage**

Double precision:

```
ierr = ASL_ham1mm (ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);
```

Single precision:

```
ierr = ASL_gam1mm (ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Real part of complex matrix $A$ (two-dimensional array type).
2	ai	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Imaginary part of complex matrix $A$ (two-dimensional array type).
3	lma	I	1	Input	Adjustable dimension of array ar and ai.
4	nm	I	1	Input	Number of rows in matrix $A$ .
5	nn	I	1	Input	Number of columns in matrix $A$ .
6	br	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Real part of complex matrix $B$ (two-dimensional array type).
7	bi	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Imaginary part of complex matrix $B$ (two-dimensional array type).
8	lnb	I	1	Input	Adjustable dimension of array br and bi.
9	nl	I	1	Input	Number of columns in matrix $B$ .
10	cr	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Real part of Initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]AB$ ).
11	ci	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Imaginary part of initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]AB$ ).
12	lmc	I	1	Input	Adjustable dimension of array cr and ci.
13	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + AB$ $isw = 0$ : Obtain $C = AB$ $isw = -1$ : Obtain $C = C - AB$
14	nt	I	1	Input	Number of tasks to be generated.
15	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < nn \leq lnb$
- (c)  $nl > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$



## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	nm was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

(a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i \\ 2+i & 2+2i & 2+3i & 2+4i \\ 3+i & 3+2i & 3+3i & 3+4i \\ 4+i & 4+2i & 4+3i & 4+4i \\ 5+i & 5+2i & 5+3i & 5+4i \end{bmatrix}$$

Obtain  $C = AB$ .

(b) Input data

Matrix  $A$ , matrix  $B$ , lma = 11, lnb = 11, lmc = 11, nm = 4, nn = 5, nl = 4, isw = 0 and nt = 2.

(c) Main program

```

/*      C interface example for ASL_ham1mm */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lma=11;
    int nm=4;
    int nn=5;
    double *br;
    double *bi;
    int lnb=11;
    int nl=4;
    double *cr;
    double *ci;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_ham1mm ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n\n", lma, lnb, lmc );
    printf( "\tnm=%2d  nn=%2d  nl=%2d\n\n", nm, nn, nl );
    printf( "\tisw=%2d  nt=%2d\n\n", isw, nt );

    ar = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
    if( ar == NULL )

```

```

{
    printf( "no enough memory for array ar\n" );
    return -1;
}

ai = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * (lnb*nl) ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * (lnb*nl) ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

cr = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( cr == NULL )
{
    printf( "no enough memory for array cr\n" );
    return -1;
}

ci = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( ci == NULL )
{
    printf( "no enough memory for array ci\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", ar[i+lma*j]=i+1 );
    }
    printf( "\n" );
}

printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", ai[i+lma*j]=j+1 );
    }
    printf( "\n" );
}

printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", br[i+lnb*j]=i+1 );
    }
    printf( "\n" );
}

printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", bi[i+lnb*j]=j+1 );
    }
    printf( "\n" );
}

ierr = ASL_ham1mm(ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);

printf( "\n    ** Output **\n" );
printf( "\tierr = %6d\n", ierr );

```

```

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cr[i+lmc*j] );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", ci[i+lmc*j] );
    }
    printf( "\n" );
}

free( ar );
free( ai );
free( br );
free( bi );
free( cr );
free( ci );

return 0;
}

```

(d) Output results

```

*** ASL_ham1mm ***

** Input **

lma=11  lnb=11  lmc=11
nm = 4  mn = 5  nl = 4

isw= 0  nt= 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1
  2      2      2      2
  3      3      3      3
  4      4      4      4
  5      5      5      5

Imaginary part of matrix B
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4

** Output **

ierr =      0

Real part of matrix C
  0      -15     -30     -45
 15       0     -15     -30
 30      15      0     -15
 45      30      15      0

Imaginary part of matrix C
 60      65      70      75
 65      75      85      95
 70      85     100     115
 75      95     115     135

```

### 2.2.7 ASL\_ham1mh, ASL\_gam1mh

#### Multiplying Complex Matrices (Two-Dimensional Array Type) (Real Argument Type) ( $C = C \pm AB^*$ )

(1) **Function**

Obtain the product of two complex matrices (Two-dimensional Array Type) ( $C = [C \pm]AB^*$ )

(2) **Usage**

Double precision:

ierr = ASL\_ham1mh (ar, ai, lma, nm, nn, br, bi, llb, nl, cr, ci, lmc, isw, nt);

Single precision:

ierr = ASL\_gam1mh (ar, ai, lma, nm, nn, br, bi, llb, nl, cr, ci, lmc, isw, nt);

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lma \times nm$	Input	Real part of complex matrix $A$ (two-dimensional array type).
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lma \times nm$	Input	Imaginary part of complex matrix $A$ (two-dimensional array type).
3	lma	I	1	Input	Adjustable dimension of array ar and ai.
4	nm	I	1	Input	Number of rows in matrix $A$ .
5	nn	I	1	Input	Number of columns in matrix $A$ .
6	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$llb \times nn$	Input	Real part of complex matrix $B$ (two-dimensional array type).
7	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$llb \times nn$	Input	Imaginary part of complex matrix $B$ (two-dimensional array type).
8	llb	I	1	Input	Adjustable dimension of array br and bi.
9	nl	I	1	Input	Number of rows in matrix $B$ .
10	cr	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lmc \times nl$	Input	Real part of initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]AB^*$ ).
11	ci	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lmc \times nl$	Input	Imaginary part of initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]AB^*$ ).
12	lmc	I	1	Input	Adjustable dimension of array cr and ci.
13	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + AB^*$ $isw = 0$ : Obtain $C = AB^*$ $isw = -1$ : Obtain $C = C - AB^*$
14	nt	I	1	Input	Number of tasks to be generated.
15	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < nl \leq llb$
- (c)  $nm > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	nm was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) Notes

None

(7) Example

(a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

Obtain  $C = AB^*$ .

(b) Input data

Matrix  $A$ , matrix  $B$ , lma = 11, llb = 11, lmc = 11, nm = 4, nn = 5, nl = 4, isw = 0 and nt = 2.

(c) Main program

```

/*      C interface example for ASL_ham1mh */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lma=11;
    int nm=4;
    int nn=5;
    double *br;
    double *bi;
    int lnb=11;
    int nl=4;
    double *cr;
    double *ci;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_ham1mh ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d   lnb=%2d   lmc=%2d\n\n", lma, lnb, lmc );
    printf( "\tnm =%2d   nn =%2d   nl =%2d\n\n", nm, nn, nl );
    printf( "\tisw=%2d   nt=%2d\n\n", isw, nt );

    ar = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
    if( ar == NULL )
    {
        printf( "no enough memory for array ar\n" );
    }

```

```

    }    return -1;

ai = ( double * )malloc((size_t)( sizeof(double) * (lma*nn) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * (lnb*nn) ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * (lnb*nn) ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

cr = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( cr == NULL )
{
    printf( "no enough memory for array cr\n" );
    return -1;
}

ci = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( ci == NULL )
{
    printf( "no enough memory for array ci\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", ar[i+lma*j]=i+1 );
    }
    printf( "\n" );
}
printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", ai[i+lma*j]=j+1 );
    }
    printf( "\n" );
}
printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", br[i+lnb*j]=i+1 );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", bi[i+lnb*j]=j+1 );
    }
    printf( "\n" );
}

ierr = ASL_ham1mh(ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\t ierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )

```

```

{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cr[i+lmc*j] );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", ci[i+lmc*j] );
    }
    printf( "\n" );
}

free( ar );
free( ai );
free( br );
free( bi );
free( cr );
free( ci );

return 0;
}

```

(d) Output results

```

*** ASL_ham1mh ***

** Input **

lma=11  lnb=11  lmc=11
nm = 4  nn = 5  nl = 4
isw= 0  nt= 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4

Imaginary part of matrix B
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

** Output **

ierr =      0

Real part of matrix C
  60      65      70      75
  65      75      85      95
  70      85      100     115
  75      95      115     135

Imaginary part of matrix C
  0      15      30      45
 -15      0      15      30
 -30     -15      0      15
 -45     -30     -15      0

```



### 2.2.8 ASL\_ham1hm, ASL\_gam1hm

#### Multiplying Complex Matrices (Two-Dimensional Array Type) (Real Argument Type) ( $C = C \pm A^*B$ )

(1) **Function**

Obtain the product of complex matrix  $A$  and complex matrix  $B$  ( $C = [C \pm]A^*B$ )

(2) **Usage**

Double precision:

```
ierr = ASL_ham1hm (ar, ai, lna, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);
```

Single precision:

```
ierr = ASL_gam1hm (ar, ai, lna, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{cases} D^* \\ R^* \end{cases}$	$lna \times nm$	Input	Real part of complex matrix $A$ (two-dimensional array type).
2	ai	$\begin{cases} D^* \\ R^* \end{cases}$	$lna \times nm$	Input	Imaginary part of complex matrix $A$ (two-dimensional array type).
3	lna	I	1	Input	Adjustable dimension of array ar and ai.
4	nm	I	1	Input	Number of rows in matrices $AR$ and $AI$ .
5	nn	I	1	Input	Number of columns in matrices $AR$ and $AI$ .
6	br	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Real part of complex matrix $B$ (two-dimensional array type).
7	bi	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Imaginary part of complex matrix $B$ (two-dimensional array type).
8	lnb	I	1	Input	Adjustable dimension of array br and bi.
9	nl	I	1	Input	Number of columns in matrix $B$ .
10	cr	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Real part of initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Real part of product ( $C = [C \pm]A^*B$ ).
11	ci	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Imaginary part of initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Imaginary part of product ( $C = [C \pm]A^*B$ ).
12	lmc	I	1	Input	Adjustable dimension of array cr and ci.
13	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + A^*B$ $isw = 0$ : Obtain $C = A^*B$ $isw = -1$ : Obtain $C = C - A^*B$
14	nt	I	1	Input	Number of tasks to be generated.
15	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lmc$
- (b)  $0 < nn \leq lna, lnb$
- (c)  $nl > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	nm was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

(a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i \\ 2+i & 2+2i & 2+3i & 2+4i \\ 3+i & 3+2i & 3+3i & 3+4i \\ 4+i & 4+2i & 4+3i & 4+4i \end{bmatrix}$$

Obtain  $C = A*B$ .

(b) Input data

Matrix  $A$ , matrix  $B$ ,  $lma = 11$ ,  $lmb = 11$ ,  $lmc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$  and  $isw = 0$ .

(c) Main program

```

/*      C interface example for ASL_ham1hm */

#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lma=11;
    int nm=5;
    int nn=4;
    double *br;
    double *bi;
    int lnb=11;
    int nl=4;
    double *cr;
    double *ci;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_ham1hm ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d   lnb=%2d   lmc=%2d\n\n", lma, lnb, lmc );
    printf( "\tnm =%2d   nn =%2d   nl =%2d\n\n", nm, nn, nl );
    printf( "\tisw=%2d   nt=%2d\n\n", isw, nt );

    ar = ( double * )malloc((size_t)( sizeof(double) * (lma*nm) ));
    if( ar == NULL )
    {
        printf( "no enough memory for array ar\n" );
    }

```

```

    } return -1;

ai = ( double * )malloc((size_t)( sizeof(double) * (lma*nm) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * (lnb*nl) ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * (lnb*nl) ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

cr = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( cr == NULL )
{
    printf( "no enough memory for array cr\n" );
    return -1;
}

ci = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( ci == NULL )
{
    printf( "no enough memory for array ci\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", ar[i+lma*j]=i+1 );
    }
    printf( "\n" );
}
printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", ai[i+lma*j]=j+1 );
    }
    printf( "\n" );
}
printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", br[i+lnb*j]=i+1 );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", bi[i+lnb*j]=j+1 );
    }
    printf( "\n" );
}

ierr = ASL_ham1hm(ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\t ierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )

```

```

    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", cr[i+lmc*j] );
        }
        printf( "\n" );
    }
    printf( "\n\tImaginary part of matrix C\n" );
    for( i=0 ; i<nm ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<nl ; j++ )
        {
            printf( "%8.3g ", ci[i+lmc*j] );
        }
        printf( "\n" );
    }

    free( ar );
    free( ai );
    free( br );
    free( bi );
    free( cr );
    free( ci );

    return 0;
}

```

(d) Output results

```

*** ASL_ham1hm ***

** Input **

lma=11  lnb=11  lmc=11
nm = 5  nn = 4  nl = 4
isw= 0  nt= 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1
  2      2      2      2
  3      3      3      3
  4      4      4      4

Imaginary part of matrix B
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4

** Output **

ierr =      0

Real part of matrix C
  34      38      42      46
  38      46      54      62
  42      54      66      78
  46      62      78      94
  50      70      90      110

Imaginary part of matrix C
  0      10      20      30
 -10      0      10      20
 -20     -10      0      10
 -30     -20     -10      0
 -40     -30     -20     -10

```

**2.2.9 ASL\_ham1hh, ASL\_gam1hh****Multiplying Complex Matrices (Two-Dimensional Array Type) (Real Argument Type) ( $C = C \pm A^*B^*$ )****(1) Function**

Obtain the product of complex matrix  $A$  and complex matrix  $B$  ( $C = [C \pm]A^*B^*$ )

**(2) Usage**

Double precision:

```
ierr = ASL_ham1hh (ar, ai, lna, nm, nn, br, bi, llb, nl, cr, ci, lmc, isw, nt);
```

Single precision:

```
ierr = ASL_gam1hh (ar, ai, lna, nm, nn, br, bi, llb, nl, cr, ci, lmc, isw, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{lna} \times \text{nm}$	Input	Real part of complex matrix $A$ (two-dimensional array).
2	ai	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{lna} \times \text{nm}$	Input	Imaginary part of complex matrix $A$ (two-dimensional array).
3	lna	I	1	Input	Adjustable dimension of array ar and ai.
4	nm	I	1	Input	Number of rows in matrix $A$ .
5	nn	I	1	Input	Number of columns in matrix $A$ .
6	br	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{llb} \times \text{nn}$	Input	Real part of complex matrix $A$ (two-dimensional array type).
7	bi	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{llb} \times \text{nn}$	Input	Imaginary part of complex matrix $B$ (two-dimensional array type).
8	llb	I	1	Input	Adjustable dimension of array br and bi.
9	nl	I	1	Input	Number of columns in matrix $B$ .
10	cr	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{lmc} \times \text{nl}$	Input	Real part of initial complex matrix $C$ (If $\text{isw} = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]A^*B^*$ ).
11	ci	$\begin{cases} D^* \\ R^* \end{cases}$	$\text{lmc} \times \text{nl}$	Input	Imaginary part of initial complex matrix $C$ (If $\text{isw} = \pm 1$ ) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]A^*B^*$ ).
12	lmc	I	1	Input	Adjustable dimension of array cr and ci.
13	isw	I	1	Input	Processing switch. $\text{isw} = 1$ : Obtain $C = C + A^*B^*$ $\text{isw} = 0$ : Obtain $C = A^*B^*$ $\text{isw} = -1$ : Obtain $C = C - A^*B^*$
14	nt	I	1	Input	Number of tasks to be generated.
15	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < \text{nm} \leq \text{lmc}$
- (b)  $0 < \text{nn} \leq \text{lna}$
- (c)  $0 < \text{nl} \leq \text{lnb}$
- (d)  $\text{isw} \in \{0, 1, -1\}$
- (e)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	mn was equal to 1	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

(a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \\ 5+i & 5+2i & 5+3i & 5+4i & 5+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

Obtain  $C = A^*B^*$ .

(b) Input data

Matrix  $A$ , matrix  $B$ , lna = 11, llb = 11, lnc = 11, nm = 4, nn = 5, nl = 6 and isw = 0.

(c) Main program

```

/*      C interface example for ASL_ham1hh */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lma=11;
    int nm=5;
    int nn=5;
    double *br;
    double *bi;
    int lnb=11;
    int nl=4;
    double *cr;
    double *ci;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;
    int i,j;

    printf( "      *** ASL_ham1hh ***\n" );
    printf( "\n      ** Input **\n\n" );

    printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n\n", lma, lnb, lmc );
    printf( "\tnm =%2d  nn =%2d  nl =%2d\n\n", nm, nn, nl );
    printf( "\tisw=%2d  nt=%2d\n\n", isw, nt );

    ar = ( double * )malloc((size_t)( sizeof(double) * (lma*nm) ));
    if( ar == NULL )

```



```

{
    printf( "no enough memory for array ar\n" );
    return -1;
}

ai = ( double * )malloc((size_t)( sizeof(double) * (lma*nm) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * (lnb*nn) ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * (lnb*nn) ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

cr = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( cr == NULL )
{
    printf( "no enough memory for array cr\n" );
    return -1;
}

ci = ( double * )malloc((size_t)( sizeof(double) * (lmc*nl) ));
if( ci == NULL )
{
    printf( "no enough memory for array ci\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", ar[i+lma*j]=i+1 );
    }
    printf( "\n" );
}

printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", ai[i+lma*j]=j+1 );
    }
    printf( "\n" );
}

printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", br[i+lnb*j]=i+1 );
    }
    printf( "\n" );
}

printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", bi[i+lnb*j]=j+1 );
    }
    printf( "\n" );
}

}

ierr = ASL_ham1hh(ar, ai, lma, nm, nn, br, bi, lnb, nl, cr, ci, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

```

```

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cr[i+lmc*j] );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", ci[i+lmc*j] );
    }
    printf( "\n" );
}

free( ar );
free( ai );
free( br );
free( bi );
free( cr );
free( ci );

return 0;
}

```

(d) Output results

```

*** ASL_ham1hh ***

** Input **

lma=11  lnb=11  lmc=11
nm = 5  nn = 5  nl = 4
isw= 0  nt= 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
  5      5      5      5      5
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4

Imaginary part of matrix B
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

** Output **

ierr =      0

Real part of matrix C
  0      15      30      45
 -15      0      15      30
 -30     -15      0      15
 -45     -30     -15      0
 -60     -45     -30     -15

Imaginary part of matrix C
 -60     -65     -70     -75
 -65     -75     -85     -95
 -70     -85     -100    -115
 -75     -95     -115    -135
 -80    -105    -130    -155

```

**2.2.10 ASL\_han1mm, ASL\_gan1mm****Multiplying Complex Matrices (Two-Dimensional Array Type) (Complex Argument Type) ( $C = C \pm AB$ )****(1) Function**

Obtain the product of two complex matrices (Two-dimensional Array Type) ( $C = [C \pm]AB$ ).

**(2) Usage**

Double precision:

ierr = ASL\_han1mm (a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_gan1mm (a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

**(3) Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Complex matrix $A$ (two-dimensional array type).
2	lma	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	$lnb \times nl$	Input	Complex matrix $B$ (two-dimensional array type).
6	lnb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm]AB$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + AB$ $isw = 0$ : Obtain $C = AB$ $isw = -1$ : Obtain $C = C - AB$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < nn \leq lnb$
- (c)  $nl > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	nn was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

(6) **Notes**

None

(7) **Example**

## (a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i \\ 2+i & 2+2i & 2+3i & 2+4i \\ 3+i & 3+2i & 3+3i & 3+4i \\ 4+i & 4+2i & 4+3i & 4+4i \\ 5+i & 5+2i & 5+3i & 5+4i \end{bmatrix}$$

Obtain  $C = AB$ .

## (b) Input data

Matrix  $A$ , matrix  $B$ ,  $lma = 11$ ,  $lnb = 11$ ,  $lnc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 4$ ,  $isw = 0$  and  $nt = 2$ .

## (c) Main program

```

/*      C interface example for ASL_han1mm */

#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lma=11;
    int nm=4;
    int nn=5;
    double _Complex *b;
    int lnb=11;
    int nl=4;
    double _Complex *c;
    int lmc=11;

```

```

int isw=0;
int nt=2;
int ierr;
int i,j;

printf( "    *** ASL_han1mm ***\n" );
printf( "\n    ** Input **\n\n" );

printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n", lma, lnb, lmc );
printf( "\tnm =%2d  nn =%2d  nl =%2d\n", nm, nn, nl );
printf( "\tisz=%2d  nt =%2d\n", isw, nt );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lma*nn) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lnb*nl) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lmc*nl) ));
if( c == NULL )
{
    printf( "no enough memory for array c\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        a[i+lma*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(a[i+lma*j]) );
    }
    printf( "\n" );
}

printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", cimag(a[i+lma*j]) );
    }
    printf( "\n" );
}

printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        b[i+lnb*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(b[i+lnb*j]) );
    }
    printf( "\n" );
}

printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(b[i+lnb*j]) );
    }
    printf( "\n" );
}

ierr = ASL_han1mm(a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )

```

```

{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", creal(c[i+lmc*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(c[i+lmc*j]) );
    }
    printf( "\n" );
}

free( a );
free( b );
free( c );

return 0;
}

```

(d) Output results

```

*** ASL_han1mm ***

** Input **

lma=11  lnb=11  lmc=11
nm = 4  nn = 5  nl = 4
isw= 0  nt = 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1
  2      2      2      2
  3      3      3      3
  4      4      4      4
  5      5      5      5

Imaginary part of matrix B
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4

** Output **

ierr =      0

Real part of matrix C
  0      -15     -30     -45
 15       0      -15     -30
 30      15       0      -15
 45      30      15       0

Imaginary part of matrix C
 60      65      70      75
 65      75      85      95
 70      85     100     115
 75      95     115     135

```

**2.2.11 ASL\_han1mh, ASL\_gan1mh****Multiplying Complex Matrices (Two-Dimensional Array Type) (Complex Argument Type) ( $C = C \pm AB^*$ )****(1) Function**

Obtain the product of two complex matrices (Two-dimensional Array Type) ( $C = [C \pm]AB^*$ )

**(2) Usage**

Double precision:

```
ierr = ASL_han1mh (a, lma, nm, nn, b, llb, nl, c, lmc, isw, nt);
```

Single precision:

```
ierr = ASL_gan1mh (a, lma, nm, nn, b, llb, nl, c, lmc, isw, nt);
```

**(3) Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D^* \\ R^* \end{cases}$	$lma \times nn$	Input	Complex matrix $A$ (two-dimensional array type).
2	lma	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} D^* \\ R^* \end{cases}$	$llb \times nn$	Input	Complex matrix $B$ (two-dimensional array type).
6	llb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of rows in matrix $B$ .
8	c	$\begin{cases} D^* \\ R^* \end{cases}$	$lmc \times nl$	Input	Initial complex matrix $C$ (If $isw = \pm 1$ ) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm]AB^*$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. $isw = 1$ : Obtain $C = C + AB^*$ $isw = 0$ : Obtain $C = AB^*$ $isw = -1$ : Obtain $C = C - AB^*$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lma, lmc$
- (b)  $0 < nl \leq llb$
- (c)  $nm > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$nm$ was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

## (a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 1+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

Obtain  $C = AB^*$ .

## (b) Input data

Matrix  $A$ , matrix  $B$ ,  $lma = 11$ ,  $llb = 11$ ,  $lmc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 4$ ,  $isw = 0$  and  $nt = 2$ .

## (c) Main program

```

/*      C interface example for ASL_han1mh */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lma=11;
    int nm=4;
    int nn=5;
    double _Complex *b;
    int lnb=11;
    int nl=4;
    double _Complex *c;
    int lmc=11;
    int isw=0;
    int nt=2;
    int ierr;

```



```

int i,j;

printf( "    *** ASL_han1mh ***\n" );
printf( "\n    ** Input **\n\n" );

printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n\n", lma, lnb, lmc );
printf( "\tnm  =%2d  nn  =%2d  nl  =%2d\n\n", nm, nn, nl );
printf( "\tismw=%2d  nt  =%2d\n\n", isw, nt );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lma*nn) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lnb*nn) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lmc*nl) ));
if( c == NULL )
{
    printf( "no enough memory for array c\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        a[i+lma*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(a[i+lma*j]) );
    }
    printf( "\n" );
}
printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", cimag(a[i+lma*j]) );
    }
    printf( "\n" );
}
printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        b[i+lmb*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(b[i+lmb*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", cimag(b[i+lmb*j]) );
    }
    printf( "\n" );
}

ierr = ASL_han1mh(a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )

```

```

    {
        printf( "%8.3g ", creal(c[i+lmc*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(c[i+lmc*j]) );
    }
    printf( "\n" );
}

free( a );
free( b );
free( c );

return 0;
}

```

(d) Output results

```

*** ASL_han1mh ***
** Input **
lma=11  lnb=11  lmc=11
nm = 4  nn = 5  nl = 4
isw= 0  nt = 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix B
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

** Output **
ierr =      0

Real part of matrix C
  60      65      70      75
  65      75      85      95
  70      85      100     115
  75      95      115     135

Imaginary part of matrix C
  0      15      30      45
 -15     0      15      30
 -30    -15     0      15
 -45    -30    -15     0

```

**2.2.12 ASL\_han1hm, ASL\_gan1hm****Multiplying Complex Matrices (Two-Dimensional Array Type) (Complex Argument Type) ( $C = C \pm A*B$ )****(1) Function**

Obtain the product of complex matrix  $A$  and complex matrix  $B$  ( $C = [C \pm]A*B$ )

**(2) Usage**

Double precision:

ierr = ASL\_han1hm (a, lna, nm, nn, b, lnb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_gan1hm (a, lna, nm, nn, b, lnb, nl, c, lmc, isw, nt);

**(3) Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} Z* \\ C* \end{cases}$	lna×nm	Input	Complex Matrix $A$ (two-dimensional array type).
2	lna	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns of matrix $A$ .
5	b	$\begin{cases} Z* \\ C* \end{cases}$	lnb×nl	Input	Complex matrix $B$ (two-dimensional array type).
6	lnb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns of matrix $B$ .
8	c	$\begin{cases} Z* \\ C* \end{cases}$	lmc×nl	Input	Initial complex matrix $C$ (If isw = ±1) (two-dimensional array type).
				Output	Product of complex matrices ( $C = [C \pm]A*B$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. isw = 1: Obtain $C = C + A*B$ isw = 0: Obtain $C = A*B$ isw = -1: Obtain $C = C - A*B$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lcm$
- (b)  $0 < nn \leq lna, lnb$
- (c)  $nl > 0$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	nn was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

## (a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i \\ 2+i & 2+2i & 2+3i & 2+4i \\ 3+i & 3+2i & 3+3i & 3+4i \\ 4+i & 4+2i & 4+3i & 4+4i \end{bmatrix}$$

Obtain  $C = A*B$ .

## (b) Input data

Matrix  $A$ , Matrix  $B$ ,  $lna = 11$ ,  $lnb = 11$ ,  $lnc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$  and  $isw = 0$ .

## (c) Main program

```

/*      C interface example for ASL_han1hm */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lma=11;
    int nm=5;
    int nn=4;
    double _Complex *b;
    int lnb=11;
    int nl=4;
    double _Complex *c;
    int lcm=11;
    int isw=0;
    int nt=2;
    int ierr;

```

```

int i,j;

printf( "    *** ASL_han1hm ***\n" );
printf( "\n    ** Input **\n\n" );

printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n\n", lma, lnb, lmc );
printf( "\tnm =%2d  nn =%2d  nl =%2d\n\n", nm, nn, nl );
printf( "\tism=%2d  nt =%2d\n\n", isw, nt );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lma*nm) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lnb*nl) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lmc*nl) ));
if( c == NULL )
{
    printf( "no enough memory for array c\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        a[i+lma*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(a[i+lma*j]) );
    }
    printf( "\n" );
}
printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", cimag(a[i+lma*j]) );
    }
    printf( "\n" );
}
printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        b[i+lmb*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(b[i+lmb*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(b[i+lmb*j]) );
    }
    printf( "\n" );
}

ierr = ASL_han1hm(a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\t(ierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )

```

```

    {
        printf( "%8.3g ", creal(c[i+lmc*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(c[i+lmc*j]) );
    }
    printf( "\n" );
}

free( a );
free( b );
free( c );

return 0;
}

```

(d) Output results

```

*** ASL_han1hm ***
** Input **
lma=11  lnb=11  lmc=11
nm = 5  nn = 4  nl = 4
isw= 0  nt = 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      2      3
  2      2      3      4
  3      3      4      4
  4      4      4      4

Imaginary part of matrix B
  1      2      3      4
  1      2      3      4
  1      2      3      4
  1      2      3      4

** Output **
ierr =      0

Real part of matrix C
  34      38      42      46
  38      46      54      62
  42      54      66      78
  46      62      78      94
  50      70      90      110

Imaginary part of matrix C
  0      10      20      30
 -10      0      10      20
 -20     -10      0      10
 -30     -20     -10      0
 -40     -30     -20     -10

```

**2.2.13 ASL\_han1hh, ASL\_gan1hh****Multiplying Complex Matrices (Two-Dimensional Array Type) (Complex Argument Type) ( $C = C \pm A^*B^*$ )****(1) Function**

Obtain the product of complex matrix  $A$  and complex matrix  $B$  ( $C = [C \pm]A^*B^*$ )

**(2) Usage**

Double precision:

ierr = ASL\_han1hh (a, lna, nm, nn, b, llb, nl, c, lmc, isw, nt);

Single precision:

ierr = ASL\_gan1hh (a, lna, nm, nn, b, llb, nl, c, lmc, isw, nt);

**(3) Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} Z^* \\ C^* \end{cases}$	lna×nm	Input	Complex matrix $A$ (two-dimensional array type).
2	lna	I	1	Input	Adjustable dimension of array a.
3	nm	I	1	Input	Number of rows in matrix $A$ .
4	nn	I	1	Input	Number of columns in matrix $A$ .
5	b	$\begin{cases} Z^* \\ C^* \end{cases}$	llb×nn	Input	Complex matrix $B$ (two-dimensional array type).
6	llb	I	1	Input	Adjustable dimension of array b.
7	nl	I	1	Input	Number of columns in matrix $B$ .
8	c	$\begin{cases} Z^* \\ C^* \end{cases}$	lmc×nl	Input	Initial complex matrix $C$ (If isw = ±1) (two-dimensional array type).
				Output	Product of real matrices ( $C = [C \pm]A^*B^*$ ).
9	lmc	I	1	Input	Adjustable dimension of array c.
10	isw	I	1	Input	Processing switch. isw = 1: Obtain $C = C + A^*B^*$ isw = 0: Obtain $C = A^*B^*$ isw = -1: Obtain $C = C - A^*B^*$
11	nt	I	1	Input	Number of tasks to be generated.
12	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < nm \leq lcm$
- (b)  $0 < nn \leq lna$
- (c)  $0 < nl \leq lnb$
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	nn was equal to 1.	Processing continues.
3000	Restriction (a), (b) or (c) was not satisfied.	Processing is aborted.
3010	Restriction (d) was not satisfied.	
3020	Restriction (e) was not satisfied.	

## (6) Notes

None

## (7) Example

## (a) Problem

$$A = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \\ 5+i & 5+2i & 5+3i & 5+4i & 5+5i \end{bmatrix}$$

$$B = \begin{bmatrix} 1+i & 1+2i & 1+3i & 1+4i & 1+5i \\ 2+i & 2+2i & 2+3i & 2+4i & 2+5i \\ 3+i & 3+2i & 3+3i & 3+4i & 3+5i \\ 4+i & 4+2i & 4+3i & 4+4i & 4+5i \end{bmatrix}$$

Obtain  $C = A^*B^*$ .

## (b) Input data

Matrix  $A$ , matrix  $B$ ,  $lna = 11$ ,  $llb = 11$ ,  $lnc = 11$ ,  $nm = 4$ ,  $nn = 5$ ,  $nl = 6$  and  $isw = 0$ .

## (c) Main program

```

/*      C interface example for ASL_han1hh */

#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lma=11;
    int nm=5;
    int nn=5;
    double _Complex *b;
    int lnb=11;
    int nl=4;
    double _Complex *c;
    int lcm=11;

```



```

int isw=0;
int nt=2;
int ierr;
int i,j;

printf( "    *** ASL_han1hh ***\n" );
printf( "\n    ** Input **\n\n" );

printf( "\tlma=%2d  lnb=%2d  lmc=%2d\n", lma, lnb, lmc );
printf( "\tnm =%2d  nn =%2d  nl =%2d\n", nm, nn, nl );
printf( "\tisw=%2d  nt =%2d\n", isw, nt );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lma*nm) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lnb*nn) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lmc*nl) ));
if( c == NULL )
{
    printf( "no enough memory for array c\n" );
    return -1;
}

printf( "\tReal part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        a[i+lma*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(a[i+lma*j]) );
    }
    printf( "\n" );
}

printf( "\tImaginary part of matrix A\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nm ; j++ )
    {
        printf( "%8.3g ", cimag(a[i+lma*j]) );
    }
    printf( "\n" );
}

printf( "\n\tReal part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        b[i+lmb*j] = (double)(i+1) + (double)(j+1) * _Complex_I;
        printf( "%8.3g ", creal(b[i+lmb*j]) );
    }
    printf( "\n" );
}

printf( "\n\tImaginary part of matrix B\n" );
for( i=0 ; i<nl ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "%8.3g ", cimag(b[i+lmb*j]) );
    }
    printf( "\n" );
}

ierr = ASL_han1hh(a, lma, nm, nn, b, lnb, nl, c, lmc, isw, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tReal part of matrix C\n" );
for( i=0 ; i<nm ; i++ )

```

```

{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", creal(c[i+lmc*j]) );
    }
    printf( "\n" );
}
printf( "\n\tImaginary part of matrix C\n" );
for( i=0 ; i<nm ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<nl ; j++ )
    {
        printf( "%8.3g ", cimag(c[i+lmc*j]) );
    }
    printf( "\n" );
}

free( a );
free( b );
free( c );

return 0;
}

```

(d) Output results

```

*** ASL_han1hh ***
** Input **

lma=11  lnb=11  lmc=11
nm = 5  nn = 5  nl = 4
isw= 0  nt = 2

Real part of matrix A
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4
  5      5      5      5      5
Imaginary part of matrix A
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

Real part of matrix B
  1      1      1      1      1
  2      2      2      2      2
  3      3      3      3      3
  4      4      4      4      4

Imaginary part of matrix B
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5
  1      2      3      4      5

** Output **

ierr =      0

Real part of matrix C
  0      15      30      45
 -15      0      15      30
 -30     -15      0      15
 -45     -30     -15      0
 -60     -45     -30     -15

Imaginary part of matrix C
 -60     -65     -70     -75
 -65     -75     -85     -95
 -70     -85     -100    -115
 -75     -95     -115    -135
 -80    -105    -130    -155

```

## Chapter 3

---

# SIMULTANEOUS LINEAR EQUATIONS (DIRECT METHOD)

### 3.1 INTRODUCTION

This chapter describes functions that solve simultaneous linear equations and obtain the determinant value and inverse matrix of a matrix.

**Function described in this chapter divides up and allocates internal processing among threads and executes allocated processing in parallel.**

In this library, functions having the following functions are provided individually for each set of matrix characteristics and storage mode.

- (1) Perform triangular decomposition of coefficient matrix, then solve simultaneous linear equations.
- (2) Perform triangular decomposition of coefficient matrix.
- (3) Perform triangular decomposition of coefficient matrix and obtain condition number.
- (4) Solve simultaneous linear equations after triangular decomposition of coefficient matrix
- (5) Obtain determinant value and inverse matrix.

You can freely combine the various types of functions (1) through (5) to suit your processing needs. This enables you to perform efficient processing by eliminating unnecessary calculation steps.

For explanations of functions (4) and (5) (See Chapter 2 of <Basic Functions Vol. 2> .)

### 3.1.1 Methods of using functions

Methods of using functions are described below using a real matrix (two-dimensional array type) as an example.

(1) Simultaneous linear equations

(a) Using ASL\_qbgmsl

$$\text{ierr} = \text{ASL\_qbgmsl}(A, \dots, \mathbf{b}, \dots);$$

Performs a triangular decomposition of coefficient matrix  $A$  and solves  $A\mathbf{x} = \mathbf{b}$ .

(b) Using ASL\_qbgmlu and  $\left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\}$

$$\text{ierr} = \text{ASL\_qbgmlu}(A, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\} (A, \dots, \mathbf{b}, \dots);$$

ASL\_qbgmlu performs a triangular decomposition of coefficient matrix  $A$ , and  $\left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\}$  (See Section 2.2.5 of <Basic Functions Vol. 2>) solves  $A\mathbf{x} = \mathbf{b}$ .

(c) Obtaining the condition number in addition to solving simultaneous linear equations

$$\text{ierr} = \text{ASL\_qbgmlc}(A, \dots, \&\text{cond}, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\} (A, \dots, \mathbf{b}, \dots);$$

ASL\_qbgmlc calculates the condition number and performs a triangular decomposition of coefficient matrix  $A$ , and  $\left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\}$  (See Section 2.2.5 of <Basic Functions Vol. 2>) solves  $A\mathbf{x} = \mathbf{b}$ .

(2) Determinant and inverse matrix

$$\text{ierr} = \text{ASL\_qbgmlu}(A, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmidi} \\ \text{ASL\_rbgmidi} \end{array} \right\} (A, \dots, \text{det}, \dots);$$

ASL\_qbgmlu performs a triangular decomposition of matrix  $A$ , and  $\left\{ \begin{array}{l} \text{ASL\_dbgmidi} \\ \text{ASL\_rbgmidi} \end{array} \right\}$  (See Section 2.2.7 of <Basic Functions Vol. 2>) obtains the determinant and inverse matrix.

(3) Improving the solution

(a) Using ASL\_qbgmsl

$$A_2 \leftarrow A$$

$$\mathbf{b}_2 \leftarrow \mathbf{b}$$

$$\text{ierr} = \text{ASL\_qbgmsl}(A_2, \dots, \mathbf{b}_2, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmlx} \\ \text{ASL\_rbgmlx} \end{array} \right\} (A, \dots, A_2, \dots, \mathbf{b}, \dots, \mathbf{b}_2, \dots);$$

The function shown above improves the solution obtained by using ASL\_qbgmsl.

(b) Using ASL\_qbgmlu and  $\left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\}$

$$A_2 \leftarrow A$$

$$\mathbf{b}_2 \leftarrow \mathbf{b}$$

$$\text{ierr} = \text{ASL\_qbgmlu}(A_2, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\} (A_2, \dots, \mathbf{b}_2, \dots);$$

$$\text{ierr} = \left\{ \begin{array}{l} \text{ASL\_dbgmlx} \\ \text{ASL\_rbgmlx} \end{array} \right\} (A, \dots, A_2, \dots, \mathbf{b}, \dots, \mathbf{b}_2, \dots);$$

ASL\_qbgmlu performs a triangular decomposition of matrix  $A$ ,  $\left\{ \begin{array}{l} \text{ASL\_dbgmls} \\ \text{ASL\_rbgmls} \end{array} \right\}$  (See Section 2.2.5 of <Basic Functions Vol. 2> ) solves  $A\mathbf{x} = \mathbf{b}$ , and  $\left\{ \begin{array}{l} \text{ASL\_dbgmlx} \\ \text{ASL\_rbgmlx} \end{array} \right\}$  (See Section 2.2.8 of <Basic Functions Vol. 2> ) improves the solution.

### 3.1.2 Notes

- (1) Since the parallel processing overhead significantly affects the computation cost if the order of the matrix is small, performance may be worse than when a non-parallel function is used.
- (2) To solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$ , you could use the mathematical formula  $\mathbf{x} = A^{-1}\mathbf{b}$ . However, it would be ill-advised to solve these equations by obtaining the inverse matrix  $A^{-1}$  and multiplying it by the constant vector. For example, for a real matrix (two-dimensional array type), if you compare this method to one in which you obtain the solution by performing a triangular decomposition of the coefficient matrix, you would find that for  $n$  variables the inverse matrix method requires approximately  $n^3$  multiplications, while the triangular decomposition method requires approximately  $n^3/3$  multiplications. Clearly, the triangular decomposition method is preferable. Therefore, you should obtain the inverse matrix  $A^{-1}$  only if you actually need the inverse matrix itself.
- (3) If you need to perform calculations many times for the same matrix such as when solving multiple sets of simultaneous linear equations where only the constant vector differs, it is more efficient to first perform the triangular decomposition once and then use that result repetitively thereafter.

**Example :**

To solve the equations:

$$\begin{aligned} A\mathbf{x}_1 &= \mathbf{b}_1 \\ A\mathbf{x}_2 &= \mathbf{b}_2 \end{aligned}$$

execute either:

$$\begin{aligned} \text{ierr} &= \text{ASL\_qbgmsl}(A, \dots, \mathbf{b}_1, \dots); \\ \text{ierr} &= \begin{cases} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{cases} (A, \dots, \mathbf{b}_2, \dots); \end{aligned}$$

or

$$\begin{aligned} \text{ierr} &= \text{ASL\_qbgmlu}(A, \dots); \\ \text{ierr} &= \begin{cases} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{cases} (A, \dots, \mathbf{b}_1, \dots); \\ \text{ierr} &= \begin{cases} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{cases} (A, \dots, \mathbf{b}_2, \dots); \end{aligned}$$

ASL\_qbgmsl or ASL\_qbgmlu performs the triangular decomposition of coefficient matrix  $A$ , and this result is only referred thereafter without its contents being changed.

- (4) Two functions are provided for performing triangular decomposition. One obtains the condition number and the other does not. The function that obtains the condition number requires many more calculations just to obtain the condition number. For a matrix of order  $n$ , it requires approximately  $n^2$  more multiplications than the function that does not obtain the condition number. Therefore, unless you specifically need the condition number, you can save execution time by performing triangular decomposition without obtaining the condition number.

### 3.1.3 Algorithms Used

#### 3.1.3.1 Solution of Simultaneous Linear Equations

Assume that the following simultaneous linear equations are to be solved :

$$\begin{cases} a_{11}x_1 + a_{12}x_2 + \cdots + a_{1n}x_n = b_1 \\ a_{21}x_1 + a_{22}x_2 + \cdots + a_{2n}x_n = b_2 \\ \vdots \\ a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nn}x_n = b_n \end{cases}$$

If we define the coefficient matrix  $A = (a_{ij})$  for  $(i, j = 1, 2, \dots, n)$ , the right-hand side vector  $\mathbf{b} = (b_i)$  for  $(i = 1, 2, \dots, n)$ , and the solution vector  $\mathbf{x} = (x_i)$  for  $(i = 1, 2, \dots, n)$ , then the simultaneous linear equations shown above can be represented in matrix format as:

$$A\mathbf{x} = \mathbf{b}$$

Now, if we decompose (LU decomposition) the coefficient matrix  $A = (a_{ij})$  for  $(i, j = 1, 2, \dots, n)$  into  $PA = LU$  as the product of the lower triangular matrix  $L = (l_{ij})$  for  $(i, j = 1, 2, \dots, n)$  and upper triangular matrix  $U = (u_{ij})$  for  $(i, j = 1, 2, \dots, n)$ , the solution vector  $\mathbf{x} = (x_i)$  for  $(i = 1, 2, \dots, n)$  is obtained by sequentially solving the following matrix equations:

$$\begin{cases} L\mathbf{y} = P\mathbf{b} \\ U\mathbf{x} = \mathbf{y} \end{cases}$$

where, the matrix  $P$  is an  $n \times n$  row exchange matrix corresponding to  $n$  row exchanges, and the vector  $\mathbf{y} = (y_i)$  for  $(i = 1, 2, \dots, n)$  is a work vector that is assigned intermediate calculation results. These matrix equations can be solved easily by using forward substitution and back substitution since the coefficient matrices are triangular matrices.

#### 3.1.3.2 LU Decomposition (Gauss Method)

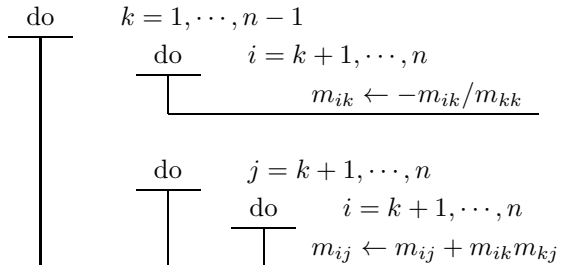
Consider the decomposition of the regular matrix  $A = (a_{ij})$  for  $(i, j = 1, 2, \dots, n)$  into  $PA = LU$  as the product of the lower triangular matrix  $L = (l_{ij})$  for  $(i, j = 1, 2, \dots, n)$  and upper triangular matrix  $U = (u_{ij})$  for  $(i, j = 1, 2, \dots, n)$ . This kind of decomposition can be done at any time for a regular matrix. For simplicity, let  $A' = PA = (a'_{ij})$  for  $(i, j = 1, 2, \dots, n)$ . Then, from the definition of matrix multiplication, we have:

$$a'_{ij} = \begin{cases} \sum_{k=1}^{j-1} l_{ik}u_{ki} + u_{ij} & (i \leq j) \\ \sum_{k=1}^{j-1} l_{ik}u_{ki} + l_{ij}u_{jj} & (i > j) \end{cases}$$

To uniquely determine the decomposition, let  $l_{ii} = 1$  for  $(i = 1, 2, \dots, n)$ . By transforming this expression, various algorithms for calculating the LU decomposition are obtained. Consider an algorithm in which the matrix  $A'$  is assigned as the initial values of the  $n \times n$  matrix  $M = (m_{ij})$  for  $(i, j = 1, 2, \dots, n)$ , and after the calculations are performed, the each element of the lower triangular matrix  $L = (l_{ij})$  for  $(i, j = 1, 2, \dots, n)$  and upper triangular matrix  $U = (u_{ij})$  for  $(i, j = 1, 2, \dots, n)$  is assigned for the element of that matrix respectively as follows:

$$M = \begin{bmatrix} u_{11} & u_{12} & u_{13} & \cdots & u_{1n} \\ -l_{21} & u_{22} & u_{23} & \cdots & u_{2n} \\ -l_{31} & -l_{32} & u_{33} & \cdots & u_{3n} \\ \vdots & \vdots & \ddots & \ddots & \vdots \\ -l_{n1} & -l_{n2} & \cdots & -l_{nn-1} & u_{nn} \end{bmatrix}$$

This algorithm (which is called the kji-SAXPY format Gauss method) is represented as follows.

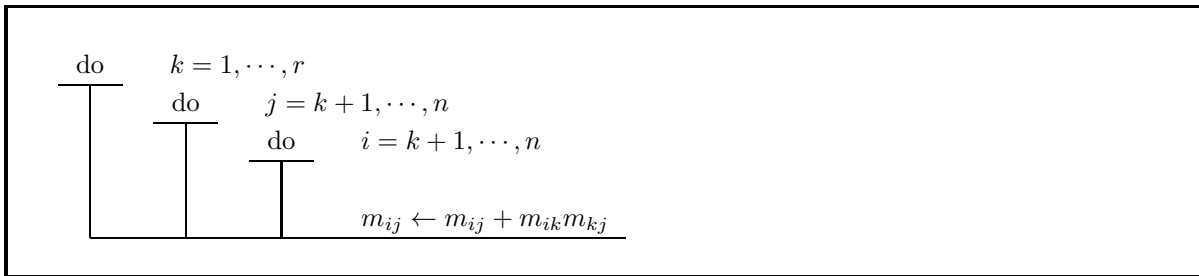


To prevent division by an element  $m_{kk} \simeq 0$  when implementing the algorithm shown above, select the element  $m_{pk}$  for which the following equation holds:

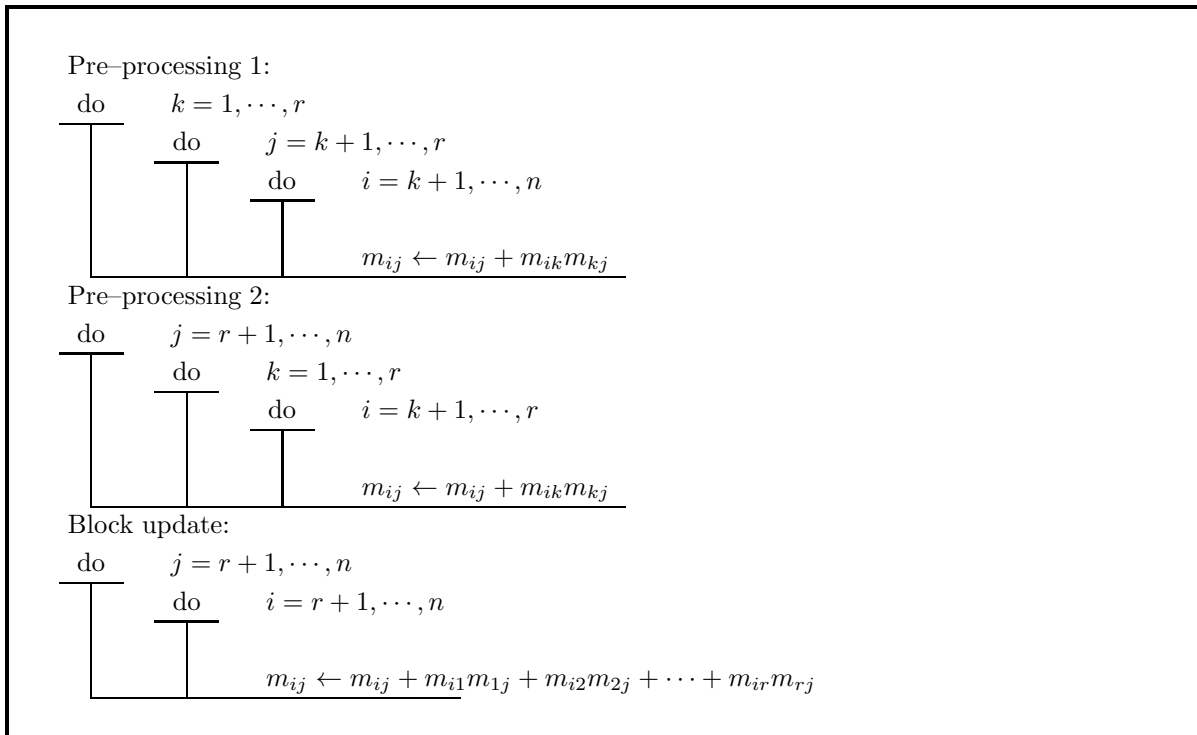
$$|m_{pk}| = \max_{k \leq i \leq n} |m_{ik}|$$

and perform partial pivoting to exchange rows  $p$  and  $k$  before the division. The matrix that expresses the accumulation of these row exchange operations is the row exchange matrix  $P$ .

This library reduce the number of substitutions in order to increase calculation speed by taking  $r$  loops for  $k$  at a time and repeatedly applying the following algorithm transformation (blocking).







Also, by modifying  $r$ , speed is increased by executing this kind of transformation recursively even for pre-processing 1.

Parallelization of this algorithm is described below.

The algorithm mentioned above can be parallelized by taking the pre-processing 2 and block update portions of the algorithm and dividing up the loop for  $j$  of these portions according to the number of parallel processing tasks, and executing these loop subdivisions in parallel. This library increase efficiency by using task division that also takes into account pre-processing 1 for the next block update. (See reference bibliography (1).)

### **3.1.4 Reference Bibliography**

- (1) Robert, Y. and Sguazzero, P. , “The LU decomposition algorithm and its efficient FORTRAN implementation on IBM 3090 Vector Multiprocessor”, IBM Tech. Rep. , ICE-0006, (1987).

## 3.2 REAL MATRIX (TWO-DIMENSIONAL ARRAY TYPE)

### 3.2.1 ASL\_qbgmsm

#### Simultaneous Linear Equations with Multiple Right-Hand Sides (Real Matrix)

(1) **Function**

ASL\_qbgmsm uses Gauss' method to solve the simultaneous linear equations  $A\mathbf{x}_i = \mathbf{b}_i (i = 1, 2, \dots, m)$  having real matrix  $A$  (two-dimensional array type) as coefficient matrix. That is, when the  $n \times m$  matrix  $B$  is defined by  $B = [\mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m]$ , the function obtains  $[\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m] = A^{-1}B$ .

(2) **Usage**

Double precision:

ierr = ASL\_qbgmsm (ab, lna, n, m, ipvt, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ab	D*	$lna \times (n + m)$	Input	Matrix (real matrix, two-dimensional array type) consisting of coefficient matrix $A$ and right-hand side vectors $\mathbf{b}_i$ [ $A, \mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m$ ]
				Output	Matrix (real matrix, two-dimensional array type) consisting of the factored matrix $A'$ of coefficient matrix $A$ and solution vectors $\mathbf{x}_i$ [ $A', \mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m$ ] (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array ab
3	n	I	1	Input	Order of matrix $A$
4	m	I	1	Input	Number of right-hand side vectors, $m$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step. (See Note (a))
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

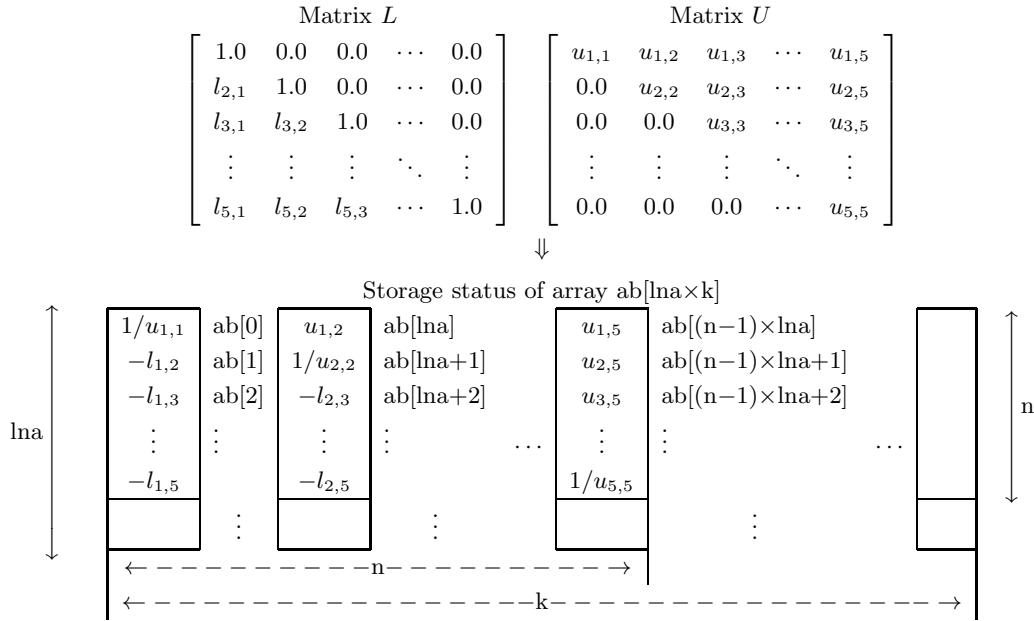
- (a)  $0 < n \leq \text{lna}$
- (b)  $0 < m$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ier value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$	$\text{ab}[\text{lna} * (n + i - 1)] \leftarrow \text{ab}[\text{lna} * (n + i - 1)] / \text{ab}[0]$ ( $i = 1, 2, \dots, m$ ) is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) **Notes**

- (a) This function perform partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (b) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array  $\text{ab}$  with the sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array  $\text{ab}$ . In addition, the reciprocals of the diagonal components of  $U$  are stored.

Figure 3-1 Storage Status of Matrices  $L$  and  $U$ **Remarks**

a.  $l_{na} \geq n$  and  $n+m \leq k$  must be hold.

(c) Shared memory parallel function for single-precision is not supported.

**(7) Example**

(a) ProblemSolve the following simultaneous linear equations.

$$\begin{bmatrix} 2 & 4 & -1 & 6 \\ -1 & -5 & 4 & 2 \\ 1 & 2 & 3 & 1 \\ 3 & 5 & -1 & -3 \end{bmatrix} \begin{bmatrix} x_{1,1} & x_{1,2} \\ x_{2,1} & x_{2,2} \\ x_{3,1} & x_{3,2} \\ x_{4,1} & x_{4,2} \end{bmatrix} = \begin{bmatrix} 36 & 11 \\ 15 & 0 \\ 22 & 7 \\ -6 & 4 \end{bmatrix}$$

(b) Input data

Array  $ab$  in which coefficient matrix  $A$  and constant vectors  $\mathbf{b}_1$  and  $\mathbf{b}_2$  are stored,  $l_{na}=11$ ,  $n=4$  and  $m=2$ .

(c) Main program

```

/*      C interface example for ASL_qbgmsm */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ab;
    int lna=11, lma=5;
    int n;
    int m;
    int *ipvt;
    int nt=2;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "qbgmsm.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_qbgmsm ***\n" );

```

```

printf( "\n    ** Input **\n\n" );
fscanf( fp, "%d", &n );
fscanf( fp, "%d", &m );

ab = ( double * )malloc((size_t)( sizeof(double) * (lna*(lna+lma)) ));
if( ab == NULL )
{
    printf( "no enough memory for array ab\n" );
    return -1;
}

ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( ipvt == NULL )
{
    printf( "no enough memory for array ipvt\n" );
    return -1;
}

printf( "\n\tCoefficient Matrix\n\n");
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ab[i+lna*j] );
        printf( "%8.3g ", ab[i+lna*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vectors\n\n");
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<m ; j++ )
    {
        fscanf( fp, "%lf", &ab[i+lna*(n+j)] );
        printf( "%8.3g ", ab[i+lna*(n+j)] );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_qbgmsm(ab, lna, n, m, ipvt, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<m ; j++ )
    {
        printf( "%8.3g ", ab[i+lna*(n+j)] );
    }
    printf( "\n" );
}

free( ab );
free( ipvt );

return 0;
}

```

(d) Output results

\*\*\* ASL\_qbgmsm \*\*\*

\*\* Input \*\*

Coefficient Matrix

2	4	-1	6
-1	-5	4	2
1	2	3	1
3	5	-1	-3

Constant Vectors

36	11
15	0
22	7
-6	4

```
  ** Output **  
  ierr =      0  
  Solution  
      1      1  
      2      1  
      4      1  
      5      1
```

### 3.2.2 ASL\_qbgmsl Simultaneous Linear Equations (Real Matrix)

(1) **Function**

ASL\_qbgmsl uses the Gauss method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the real matrix  $A$  (two-dimensional array type) as coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_qbgmsl (a, lna, n, b, ipvt, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int} \text{ as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	D*	lna × n	Input	Coefficient matrix $A$ (real matrix, two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (b) and (c))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	D*	n	Input	Constant vector $\mathbf{b}$
				Output	Solution vector $\mathbf{x}$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step. (See Note (b))
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$

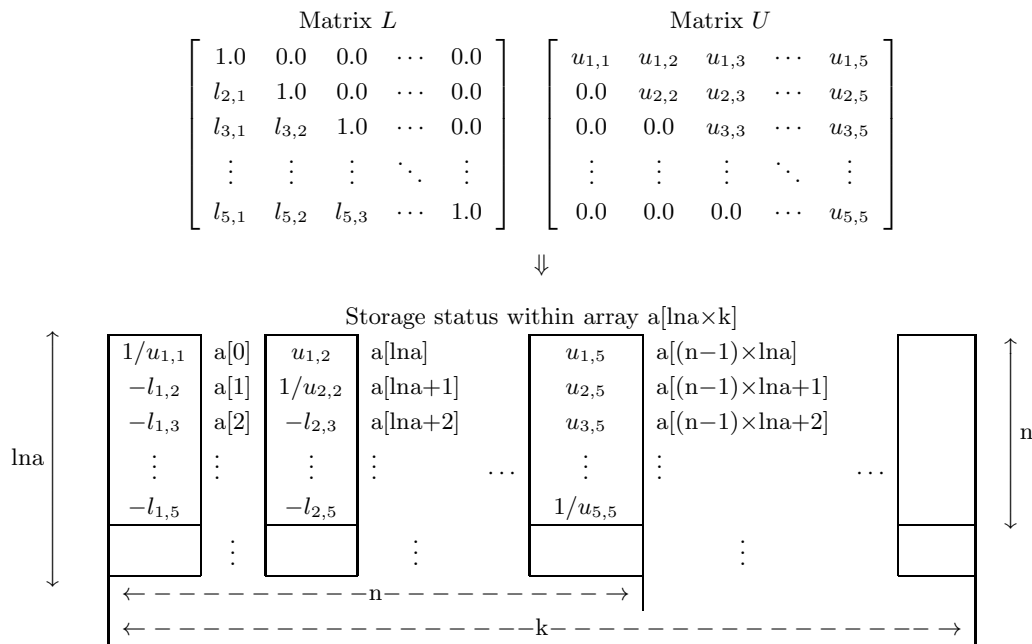


(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$	$b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) Notes

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, the solution is obtained more efficiently by directly using the function 3.2.1 ASL\_qbgmsm to perform the calculations. However, when 3.2.1 ASL\_qbgmsm cannot be used such as when all of the right-hand side vectors  $\mathbf{b}$  are not known in advance, call this function only once and then call function 2.2.5  $\left\{ \begin{array}{l} \text{ASL\_dbgmsl} \\ \text{ASL\_rbgmsl} \end{array} \right\}$  (in <Basic Functions Vol. 2>) the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculations by performing the LU decomposition of matrix  $A$  only once.
- (b) This function performs partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $ipvt[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array  $\mathbf{a}$  with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array  $\mathbf{a}$ . Also, reciprocals are stored for the diagonal components of  $U$ .
- (d) Shared memory parallel function for single-precision is not supported.



**Remarks**

- a.  $l_{na} \geq n$  and  $n \leq k$  must hold.

Figure 3–2 Storage Status of Matrices  $L$  and  $U$

(7) **Example**

(a) **Problem**

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 2 & 4 & -1 & 6 \\ -1 & -5 & 4 & 2 \\ 1 & 2 & 3 & 1 \\ 3 & 5 & -1 & -3 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 36 \\ 15 \\ 22 \\ -6 \end{bmatrix}$$

(b) **Input data**

Coefficient matrix  $A$ ,  $l_{na}=11$ ,  $n=4$ , constant vector  $\mathbf{b}$  and  $nt=2$ .

(c) **Main Program**

```

/*      C interface example for ASL_qbgmsl */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lna;
    int n;
    double *b;
    int *ipvt;
    int ierr;
    int nt = 2;
    int i,j;
    FILE *fp;

    fp = fopen( "qbgmsl.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_qbgmsl ***\n" );

```

```

printf( "\n    ** Input **\n\n" );
fscanf( fp, "%d", &lna );
fscanf( fp, "%d", &n );

a = ( double * )malloc((size_t)( sizeof(double) * (lna*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double * )malloc((size_t)( sizeof(double) * n ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( ipvt == NULL )
{
    printf( "no enough memory for array ipvt\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\t nt = %6d\n", nt );
printf( "\n\tCoefficient Matrix\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &a[i+lna*j] );
        printf( "%8.3g ", a[i+lna*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector\n\n" );
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &b[i] );
    printf( "\t%8.3g\n", b[i] );
}

fclose( fp );

ierr = ASL_qbgmsl(a, lna, n, b, ipvt, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = %8.3g\n", i,b[i] );
}

free( a );
free( b );
free( ipvt );

return 0;
}

```

(d) Output results

\*\*\* ASL\_qbgmsl \*\*\*

\*\* Input \*\*

n = 4  
 nt = 2

Coefficient Matrix

2	4	-1	6
-1	-5	4	2
1	2	3	1
3	5	-1	-3

Constant Vector

36
15
22
-6

**\*\* Output \*\***

ierr = 0

Solution

x[	0]	=	1
x[	1]	=	2
x[	2]	=	4
x[	3]	=	5

### 3.2.3 ASL\_qbgmlu LU Decomposition of a Real Matrix

(1) **Function**

ASL\_qbgmlu uses the Gauss method to perform an LU decomposition of the real matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

`ierr = ASL_qbgmlu (a, lna, n, ipvt, nt);`

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	D*	lna×n	Input	Real matrix $A$ (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step. (See Note (b))
5	nt	I	1	Input	Number of tasks to be generated
6	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n is equal to 1	Contents of array a are not changed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
4000+i	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular	

(6) **Notes**

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array a with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array a. Also, reciprocals are stored for the diagonal components of  $U$ . (See Fig. 3–2.)
- (b) This function performs partial pivoting. Pivoting information is stored in array ipvt for use by subsequent functions. If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in ipvt[ $i-1$ ]. In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Shared memory parallel function for single-precision is not supported.

### 3.2.4 ASL\_qbgmlc

#### LU Decomposition and Condition Number of a Real Matrix

(1) **Function**

ASL\_qbgmlc uses the Gauss method to perform an LU decomposition and obtain the condition number of the real matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

`ierr = ASL_qbgmlc (a, lna, n, ipvt, &cond, w1, nt);`

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	D*	lna×n	Input	Real matrix $A$ (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step (See Note (b))
5	cond	D*	1	Output	Reciprocal of the condition number
6	w1	D*	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n is equal to 1	Contents of array a are not changed and $\text{cond} \leftarrow 1.0$ is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
4000+i	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular.	Processing is aborted. The condition number is not obtained.

(6) **Notes**

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array a with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array a. Also, reciprocals are stored for the diagonal components of  $U$ . (See Fig. 3–2.)
- (b) This function partial performs pivoting. Pivoting information is stored in array ipvt for use by subsequent functions. If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i-1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Although the condition number is defined by  $\|A\| \cdot \|A^{-1}\|$ , the one obtained by this function is an approximation.
- (d) Shared memory parallel function for single-precision is not supported.



---

### 3.3 COMPLEX MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (REAL ARGUMENT TYPE)

#### 3.3.1 ASL\_hbgmsm

Simultaneous Linear Equations with Multiple Right-Hand Sides (Complex Matrix)

(1) **Function**

ASL\_hbgmsm uses Gauss' method to solve the simultaneous linear equations  $A\mathbf{x}_i = \mathbf{b}_i (i = 1, 2, \dots, m)$  having complex matrix  $A$  (two-dimensional array type) as coefficient matrix. That is, when the  $n \times m$  matrix  $B$  is defined by  $B = [\mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m]$ , the function obtains  $[\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m] = A^{-1}B$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hbgmsm (abr, abi, lna, n, m, ipvt, w1, nt);
```

Single precision:

Nothing

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	abr	D*	$\text{lna} \times (\text{n} + \text{m})$	Input	Real part of matrix (complex matrix, two-dimensional array type) consisting of coefficient matrix $A$ and right-hand side vectors $\mathbf{b}_i$ [ $A, \mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m$ ]
				Output	Real part of matrix (complex matrix, two-dimensional array type) consisting of the factored matrix $A'$ of coefficient matrix $A$ and solution vectors $\mathbf{x}_i$ [ $A', \mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m$ ] (See Notes (a) and (b))
2	abi	D*	$\text{lna} \times (\text{n} + \text{m})$	Input	Imaginary part of matrix (complex matrix, two-dimensional array type) consisting of coefficient matrix $A$ and right-hand side vectors $\mathbf{b}_i$ [ $A, \mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m$ ]
				Output	Imaginary part of matrix (complex matrix, two-dimensional array type) consisting of the factored matrix $A'$ of coefficient matrix $A$ and solution vectors $\mathbf{x}_i$ [ $A', \mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m$ ] (See Notes (a) and (b))
3	lna	I	1	Input	Adjustable dimension of arrays abr and abi
4	n	I	1	Input	Order of matrix $A$
5	m	I	1	Input	Number of right-hand side vectors, $m$
6	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step. (See Note (a))
7	w1	D*	n	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

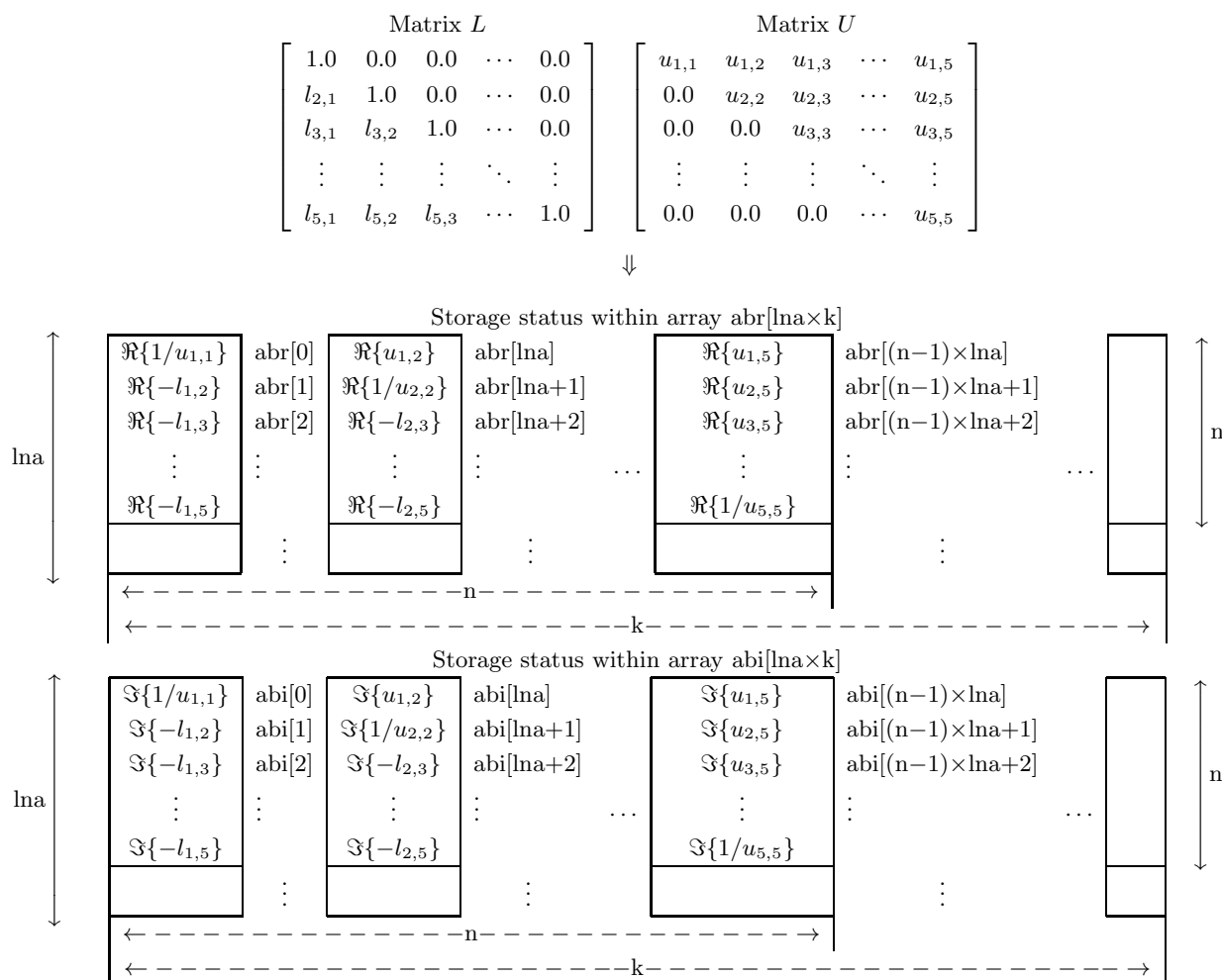
- (a)  $0 < \text{n} \leq \text{lna}$
- (b)  $0 < \text{m}$
- (c)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1	$\text{abr}[\text{lna}*(\text{n}+\text{i}-1)] \leftarrow (\text{abr}[\text{lna}*(\text{n}+\text{i}-1)] \times \text{abr}[0] + \text{abi}[\text{lna}*(\text{n}+\text{i}-1)] \times \text{abi}[0]) / (\text{abr}[0]^2 + \text{abi}[0]^2),$ $\text{abi}[\text{lna}*(\text{n}+\text{i}-1)] \leftarrow (\text{abi}[\text{lna}*(\text{n}+\text{i}-1)] \times \text{abr}[0] - \text{abr}[\text{lna}*(\text{n}+\text{i}-1)] \times \text{abi}[0]) / (\text{abr}[0]^2 + \text{abi}[0]^2)$ (i=1, 2, ..., m)
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
4000+i	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of coefficient matrix $A$ . $A$ is singular.	

## (6) Notes

- (a) This function perform partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (b) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array  $\text{abr}$  and  $\text{abi}$  with the sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array  $\text{abr}$  and  $\text{abi}$ . In addition, the reciprocals of the diagonal components of  $U$  are stored.
- (c) Shared memory parallel function for single-precision is not supported.

Figure 3-3 Storage Status of Matrices  $L$  and  $U$ 

## (7) Example

## (a) Problem

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 4+2i & 3+9i & 4+i & 7+9i \\ 5+7i & 4i & 4+7i & 2+5i \\ 9+3i & 6+2i & 9+5i & 8+5i \\ 1+5i & 7+9i & 3+5i & 2+4i \end{bmatrix} \begin{bmatrix} x_{1,1} & x_{1,2} & x_{1,3} & x_{1,4} \\ x_{2,1} & x_{2,2} & x_{2,3} & x_{2,4} \\ x_{3,1} & x_{3,2} & x_{3,3} & x_{3,4} \\ x_{4,1} & x_{4,2} & x_{4,3} & x_{4,4} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

## (b) Input data

Array  $abr$  and  $abi$  in which coefficient matrix  $A$  and constant vectors  $\mathbf{b}_1$  and  $\mathbf{b}_2$  are stored,  $l_{na}=11$ ,  $n=4$  and  $m=4$ .

## (c) Main program

```
/*      C interface example for ASL_hbgmsm */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
```

```

{
double *abr;
double *abi;
int lna=11, lma=5;
int n;
int m;
int *ipvt;
double *w1;
int ierr;
int i,j;
int nt=2;
FILE *fp;

fp = fopen( "hbgmsm.dat", "r" );
if( fp == NULL )
{
printf( "file open error\n" );
return -1;
}

printf( " *** ASL_hbgmsm ***\n" );
printf( "\n ** Input **\n\n" );

fscanf( fp, "%d", &n );
fscanf( fp, "%d", &m );
printf( "\t n = %6d m = %6d\n", n, m );

abr = ( double * )malloc((size_t)( sizeof(double) * (lna*(lna+lma)) ));
if( abr == NULL )
{
printf( "no enough memory for array abr\n" );
return -1;
}

abi = ( double * )malloc((size_t)( sizeof(double) * (lna*(lna+lma)) ));
if( abi == NULL )
{
printf( "no enough memory for array abi\n" );
return -1;
}

ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( ipvt == NULL )
{
printf( "no enough memory for array ipvt\n" );
return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * n ));
if( w1 == NULL )
{
printf( "no enough memory for array w1\n" );
return -1;
}

printf( "\n\tCoefficient Matrix\n\n");
for( i=0 ; i<n ; i++ )
{
printf( "\t" );
for( j=0 ; j<n ; j++ )
{
fscanf( fp, "%lf", &abr[i+lna*j] );
fscanf( fp, "%lf", &abi[i+lna*j] );
printf( "(%8.3g,%8.3g)", abr[i+lna*j], abi[i+lna*j] );
}
printf( "\n" );
}

printf( "\n\tConstant Vectors\n\n");
for( i=0 ; i<n ; i++ )
{
printf( "\t" );
for( j=0 ; j<m ; j++ )
{
fscanf( fp, "%lf", &abr[i+lna*(n+j)] );
fscanf( fp, "%lf", &abi[i+lna*(n+j)] );
printf( "(%8.3g,%8.3g)",abr[i+lna*(n+j)],abi[i+lna*(n+j)] );
}
printf( "\n" );
}

fclose( fp );

ierr = ASL_hbgmsm(abr, abi, lna, n, m, ipvt, w1, nt);

printf( "\n ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution\n\n" );

```

```

for( i=0 ; i<n ; i++ )
{
    printf( "\\t" );
    for( j=0 ; j<m ; j++ )
    {
        printf( "(%8.3g,%8.3g)", abr[i+lna*(n+j)],abi[i+lna*(n+j)] );
    }
    printf( "\\n" );
}

free( abr );
free( abi );
free( ipvt );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_hbgmsm ***

** Input **

n =      4 m =      4

Coefficient Matrix

(      4,      2)(      3,      9)(      4,      1)(      7,      9)
(      6,      7)(      0,      4)(      4,      7)(      2,      5)
(      9,      3)(      6,      2)(      9,      5)(      8,      5)
(      1,      5)(      7,      9)(      3,      5)(      2,      4)

Constant Vectors

(      1,      0)(      0,      0)(      0,      0)(      0,      0)
(      0,      0)(      1,      0)(      0,      0)(      0,      0)
(      0,      0)(      0,      0)(      1,      0)(      0,      0)
(      0,      0)(      0,      0)(      0,      0)(      1,      0)

** Output **

ierr =      0

Solution

( 0.0133, -0.073)( 0.181, -0.247)( -0.184, 0.178)( -0.104, -0.056)
( -0.0178, -0.0189)( -0.068, -0.0696)( -0.0128, 0.1)( 0.0415, -0.0657)
( -0.0353, 0.138)( -0.0585, 0.17)( 0.133, -0.241)( 0.131, 0.0191)
( 0.0494, -0.0686)(-0.00961, 0.13)( 0.0885, -0.0709)( -0.0462, 0.0662)

```

### 3.3.2 ASL\_hbgmsl Simultaneous Linear Equation (Complex Matrix)

(1) **Function**

ASL\_hbgmsl uses the Gauss method to solve the simultaneous equations  $Ax = b$  having the coefficient matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

ierr = ASL\_hbgmsl (ar, ai, lna, n, br, bi, ipvt, w, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	D*	lna×n	Input	Real part of coefficient matrix $A$ (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (b) and (c))
2	ai	D*	lna×n	Input	Imaginary part of coefficient matrix $A$ (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (b) and (c))
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	br	D*	n	Input	Real part of constant vector $b$
				Output	Real part of solution vector $x$
6	bi	D*	n	Input	Imaginary part of constant vector $b$
				Output	Imaginary part of solution vector $x$
7	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row $i$ in the $i$ -th processing step (See Note (b))
8	w	D*	n	Work	Work area
9	nt	I	1	Input	Number of tasks to be generated
10	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$	$\text{br}[0] \leftarrow \{\text{br}[0] \times \text{ar}[0] + \text{bi}[0] \times \text{ai}[0]\} / \{\text{ar}[0]^2 + \text{ai}[0]^2\}$ $\text{bi}[1] \leftarrow \{\text{bi}[0] \times \text{ar}[0] - \text{br}[0] \times \text{ai}[0]\} / \{\text{ar}[0]^2 + \text{ai}[0]^2\}$ are performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of the coefficient matrix $A$ . $A$ is singular.	

(6) Notes

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, the solution is obtained more efficiently by directly using the function 3.3.1 ASL\_hbgmsm to perform the calculations. However, when 3.3.1 ASL\_hbgmsm cannot be used such as when all of the right-hand side vectors  $\mathbf{b}$  are not known in advance, call this function only once and then call function 2.3.5  $\left\{ \begin{array}{l} \text{ASL\_zbgmsl} \\ \text{ASL\_cbgmsl} \end{array} \right\}$  (in <Basic Functions Vol. 2>) the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculations by performing the LU decomposition of matrix  $A$  only once.
- (b) This function performs partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) The unit lower triangular matrix  $L$  is stored in the lower triangular portions of arrays  $\text{ar}$  and  $\text{ai}$  with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portions. However, since the diagonal components of  $L$  always are 1.0, they are not stored in arrays  $\text{ar}$  and  $\text{ai}$ . Also, reciprocals are stored for the diagonal components of  $U$ . In Fig. 3–4,  $\Re\{z\}$  and  $\Im\{z\}$  denote a real part and an imaginary part of a complex number  $z$ , respectively.
- (d) Shared memory parallel function for single-precision is not supported.



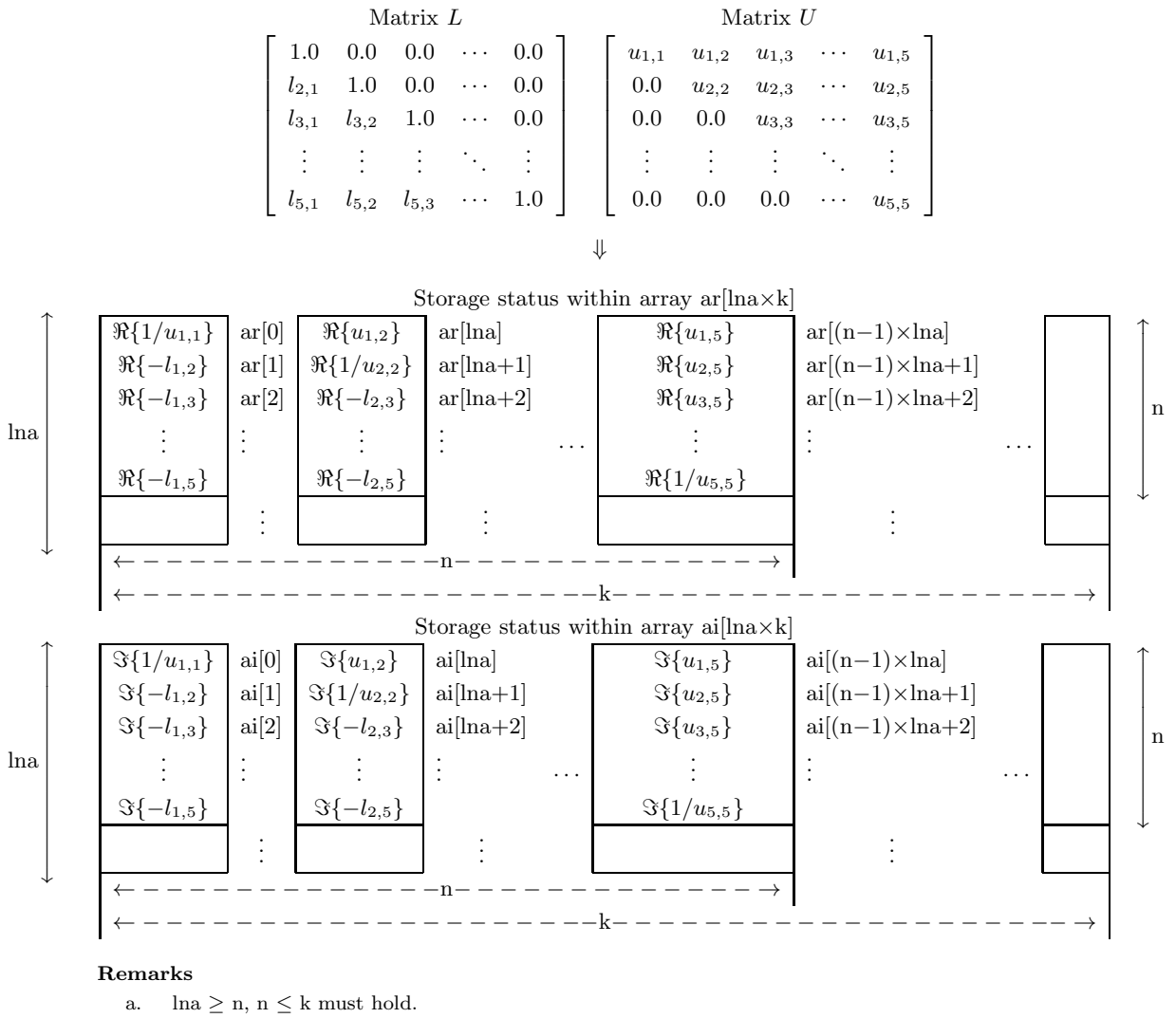


Figure 3-4 Storage Status of Matrices  $L$  and  $U$

(7) Example

(a) Problem

Obtain the condition number by solving the following simultaneous linear equations.

$$\begin{bmatrix} 5 + 8i & 7 + i & 6 + 3i & 1 + 2i \\ 1 + i & 9 + 5i & 4 + i & 5 \\ 4i & 3 + 3i & 4 + 2i & 6 + 9i \\ 7 + 8i & 6 & 7 + 6i & 10 + 4i \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 3 + 20i \\ -6 + 7i \\ -6i \\ 13i \end{bmatrix}$$

(b) Input data

Coefficient matrix real part ar and imaginary part ai, l<sub>na</sub>=11, n=4 and constant vector  $\mathbf{b}$ .

(c) Main program

```

/*      C interface example for ASL_hbgmsl */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{

```

```

double *ar;
double *ai;
int na;
int n;
int nt = 2;
double *br;
double *bi;
int *kpvt;
double *w;
int ierr;
int i,j;
FILE *fp;

fp = fopen( "hbgmsl.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_hbgmsl ***\n" );
printf( "\n    ** Input **\n\n" );

fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );

ar = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( ar == NULL )
{
    printf( "no enough memory for array ar\n" );
    return -1;
}

ai = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * n ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * n ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

w = ( double * )malloc((size_t)( sizeof(double) * n ));
if( w == NULL )
{
    printf( "no enough memory for array w\n" );
    return -1;
}

kpvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( kpvt == NULL )
{
    printf( "no enough memory for array kpvt\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\t nt = %6d\n", nt );
printf( "\n\tCoefficient Matrix (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ar[i+na*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ai[i+na*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<n ; j++ )
    {

```

```

        printf( "(%8.3g , %8.3g) ", ar[i+na*j],ai[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &br[i] );
}
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &bi[i] );
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t(%8.3g , %8.3g) \n", br[i], bi[i] );
}

fclose( fp );

ierr = ASL_hbgmsl(ar, ai, na, n, br, bi, kpvt, w, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = (%8.3g , %8.3g) \n", i, br[i], bi[i] );
}

free( ar );
free( ai );
free( br );
free( bi );
free( w );
free( kpvt );

return 0;
}

```

(d) Output results

```

*** ASL_hbgmsl ***

** Input **

n =      4
nt =     2

Coefficient Matrix (Real, Imaginary)

(      5 ,      8) (      7 ,      1) (      6 ,      3) (      1 ,      2)
(      1 ,      1) (      9 ,      5) (      4 ,      1) (      5 ,      0)
(      0 ,      4) (      3 ,      3) (      4 ,      2) (      6 ,      9)
(      7 ,      8) (      6 ,      0) (      7 ,      6) (     10 ,      4)

Constant Vector (Real, Imaginary)

(      3 ,     20)
(     -6 ,      7)
(      0 ,     -6)
(      0 ,     13)

** Output **

ierr =      0

Solution (Real, Imaginary)

x[      0] = (      1 ,      1)
x[      1] = (-2.22e-16 ,      1)
x[      2] = (      1 , -5e-16)
x[      3] = (     -1 ,     -1)

```

### 3.3.3 ASL\_hbgmlu LU Decomposition of a Complex Matrix

(1) **Function**

ASL\_hbgmlu uses the Gauss method to perform an LU decomposition of the complex matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

ierr = ASL\_hbgmlu (ar, ai, lna, n, ipvt, w, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int} \text{ as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	D*	lna×n	Input	Real part of complex matrix $A$ (two-dimensional array type)
				Output	Real part of upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
2	ai	D*	lna×n	Input	Imaginary part of complex matrix $A$ (two-dimensional array type)
				Output	Imaginary part of upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i int the i-th processing step. (See Note (b))
6	w	D*	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \ln a$
- (b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$	The contents of arrays ar and ai are unchanged.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular.	

(6) **Notes**

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portions of arrays ar and ai with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portions. However, since the diagonal components of  $L$  always are 1.0, they are not stored in arrays ar and ai. Also, reciprocals are stored for the diagonal components of  $U$ . (Refer to Fig. 3–4.)
- (b) This function performs partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $ipvt[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Shared memory parallel function for single-precision is not supported.

### 3.3.4 ASL\_hbgmlc

#### LU Decomposition and Condition Number of a Complex Matrix

(1) **Function**

ASL\_hbgmlc uses the Gauss method to perform an LU decomposition and obtain the condition number of the complex matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

```
ierr = ASL_hbgmlc (ar, ai, lna, n, ipvt, &cond, w, nt);
```

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int} \text{ as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	D*	lna × n	Input	Real part of complex matrix $A$ (two-dimensional array type)
				Output	Real part of upper triangular matrix $L$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
2	ai	D*	lna × n	Input	Imaginary part of complex matrix $A$ (two-dimensional array type)
				Output	Imaginary part of upper triangular matrix $L$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row $i$ in the $i$ -th processing step (See Note (b))
6	cond	D*	1	Output	Reciprocal of the condition number
7	w	D*	2 × n	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \text{lna}$
- (b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$	The contents of arrays ar and ai are unchanged. cond $\leftarrow$ 1.0 is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular.	Processing is aborted. The condition number is not obtained.

(6) **Notes**

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portions of arrays ar and ai with sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portions. However, since the diagonal components of  $L$  always are 1.0, they are not stored in arrays ar and ai. Also, reciprocals are stored for the diagonal components of  $U$ . (Refer to Fig. 3–4.)
- (b) This function performs partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Although the condition number is defined by  $\|A\| \cdot \|A^{-1}\|$ , the one obtained by this function is an approximation.
- (d) Shared memory parallel function for single-precision is not supported.

### 3.4 COMPLEX MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (COMPLEX ARGUMENT TYPE)

#### 3.4.1 ASL\_hbgnsn

#### Simultaneous Linear Equations with Multiple Right-Hand Sides (Complex Matrix)

(1) **Function**

ASL\_hbgnsn uses Gauss' method to solve the simultaneous linear equations  $A\mathbf{x}_i = \mathbf{b}_i (i = 1, 2, \dots, m)$  having complex matrix  $A$  (two-dimensional array type) as coefficient matrix. That is, when the  $n \times m$  matrix  $B$  is defined by  $B = [\mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m]$ , the function obtains  $[\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m] = A^{-1}B$ .

(2) **Usage**

Double precision:

ierr = ASL\_hbgnsn (ab, lna, n, m, ipvt, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ab	Z*	lna × (n + m)	Input	Matrix (complex matrix, two-dimensional array type) consisting of coefficient matrix $A$ and right-hand side vectors $\mathbf{b}_i$ [ $A, \mathbf{b}_1, \mathbf{b}_2, \dots, \mathbf{b}_m$ ]
				Output	Matrix (complex matrix, two-dimensional array type) consisting of the factored matrix $A'$ of coefficient matrix $A$ and solution vectors $\mathbf{x}_i$ [ $A', \mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m$ ] (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array ab
3	n	I	1	Input	Order of matrix $A$
4	m	I	1	Input	Number of right-hand side vectors, $m$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of row exchanged with row i in the i-th processing step. (See Note (a))
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)



(4) **Restrictions**

- (a)  $0 < n \leq \text{lna}$
- (b)  $0 < m$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$\text{ab}[\text{lna} * (n + i - 1)] \leftarrow \text{ab}[\text{lna} * (n + i - 1)] / \text{ab}[0]$ ( $i = 1, 2, \dots, m$ ) is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) **Notes**

- (a) This function perform partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (b) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array  $\text{ab}$  with the sign changed, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array  $\text{ab}$ . In addition, the reciprocals of the diagonal components of  $U$  are stored. (See Figure 3–1 in Section 3.2.1).
- (c) Shared memory parallel function for single-precision is not supported.

(7) **Example**(a) **Problem**

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 4 + 2i & 3 + 9i & 4 + i & 7 + 9i \\ 5 + 7i & 4i & 4 + 7i & 2 + 5i \\ 9 + 3i & 6 + 2i & 9 + 5i & 8 + 5i \\ 1 + 5i & 7 + 9i & 3 + 5i & 2 + 4i \end{bmatrix} \begin{bmatrix} x_{1,1} & x_{1,2} & x_{1,3} & x_{1,4} \\ x_{2,1} & x_{2,2} & x_{2,3} & x_{2,4} \\ x_{3,1} & x_{3,2} & x_{3,3} & x_{3,4} \\ x_{4,1} & x_{4,2} & x_{4,3} & x_{4,4} \end{bmatrix} = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

(b) **Input data**

Array  $\text{ab}$  in which coefficient matrix  $A$  and constant vectors  $\mathbf{b}_1$ ,  $\mathbf{b}_2$ ,  $\mathbf{b}_3$  and  $\mathbf{b}_4$  are stored,  $\text{lna}=11$ ,  $n=4$ ,  $m=4$  and  $\text{nt}=2$ .

(c) Main program

```

/*      C interface example for ASL_hbgnsnm */

#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *ab;
    int lna=11, lma=5;
    int n;
    int m;
    int *ipvt;
    int nt = 2;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "hbgnsnm.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hbgnsnm ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &n );
    fscanf( fp, "%d", &m );
    printf( "\t n = %6d m = %6d\n", n, m );

    ab = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lna*(lna+lma)) ));
    if( ab == NULL )
    {
        printf( "no enough memory for array ab\n" );
        return -1;
    }

    ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
    if( ipvt == NULL )
    {
        printf( "no enough memory for array ipvt\n" );
        return -1;
    }

    printf( "\n\tCoefficient Matrix\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<n ; j++ )
        {
            double tmp_re, tmp_im;
            fscanf( fp, "%lf", &tmp_re );
            fscanf( fp, "%lf", &tmp_im );
            ab[i+lna*j] = tmp_re + tmp_im * _Complex_I;
            printf( "(%.3g,%.3g)", creal(ab[i+lna*j]), cimag(ab[i+lna*j]) );
        }
        printf( "\n" );
    }

    printf( "\n\tConstant Vectors\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<m ; j++ )
        {
            double tmp_re, tmp_im;
            fscanf( fp, "%lf", &tmp_re );
            fscanf( fp, "%lf", &tmp_im );
            ab[i+lna*(n+j)] = tmp_re + tmp_im * _Complex_I;
            printf( "(%.3g,%.3g)", creal(ab[i+lna*(n+j)]), cimag(ab[i+lna*(n+j)]) );
        }
        printf( "\n" );
    }

    fclose( fp );

    ierr = ASL_hbgnsnm(ab, lna, n, m, ipvt, nt);

    printf( "\n      ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );
    printf( "\n\tSolution\n\n" );
    for( i=0 ; i<n ; i++ )
    {

```

```

    printf( "\t" );
    for( j=0 ; j<m ; j++ )
    {
        printf( "%8.3g,%8.3g", creal(ab[i+lna*(n+j)]),cimag(ab[i+lna*(n+j)]) );
    }
    printf( "\n" );
}

free( ab );
free( ipvt );

return 0;
}

```

(d) Output results

```

*** ASL_hbgmsm ***
** Input **
n =      4 m =      4
Coefficient Matrix
(      4,      2)(      3,      9)(      4,      1)(      7,      9)
(      6,      7)(      0,      4)(      4,      7)(      2,      5)
(      9,      3)(      6,      2)(      9,      5)(      8,      5)
(      1,      5)(      7,      9)(      3,      5)(      2,      4)
Constant Vectors
(      1,      0)(      0,      0)(      0,      0)(      0,      0)
(      0,      0)(      1,      0)(      0,      0)(      0,      0)
(      0,      0)(      0,      0)(      1,      0)(      0,      0)
(      0,      0)(      0,      0)(      0,      0)(      1,      0)
** Output **
ierr =      0
Solution
(  0.0133, -0.073)(  0.181, -0.247)( -0.184,  0.178)( -0.104, -0.056)
( -0.0178, -0.0189)( -0.068, -0.0696)( -0.0128,  0.1)(  0.0415, -0.0657)
( -0.0353,  0.138)( -0.0585,  0.17)(  0.133, -0.241)(  0.131,  0.0191)
(  0.0494, -0.0686)(-0.00961,  0.13)(  0.0885, -0.0709)( -0.0462,  0.0662)

```

### 3.4.2 ASL\_hbgnsl Simultaneous Linear Equations (Complex Matrix)

(1) **Function**

ASL\_hbgnsl uses the Gauss method or the Crout method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the complex matrix  $A$  (two-dimensional array type) as coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_hbgnsl (a, lna, n, b, ipvt, nt);

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	Z*	lna×n	Input	Coefficient matrix (complex matrix, two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ (See Notes (b) and (c))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	Z*	n	Input	Constant vector $\mathbf{b}$
				Output	Solution $\mathbf{x}$
5	ipvt	I*	n	Output	Pivoting information ipvt[i-1]: Number of the row exchanged with row i in the i-th processing step. (See Note (b))
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$  and  $\text{nt} \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LU decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) Notes

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, the solution is obtained more efficiently by directly using the function 3.4.1 ASL\_hbgns1 to perform the calculations. However, when 3.4.1 ASL\_hbgns1 cannot be used such as when all of the right-hand side vectors  $\mathbf{b}$  are not known in advance, call this function only once and then call function 2.4.5  $\left\{ \begin{matrix} \text{ASL\_zbgns1} \\ \text{ASL\_cbgns1} \end{matrix} \right\}$  the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculation by performing the LU decomposition of matrix  $A$  only once.
- (b) This function performs partial pivoting when obtaining the LU decomposition of coefficient matrix  $A$ . If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i-1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array  $\mathbf{a}$  with a minus sign added to each element, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array  $\mathbf{a}$ . Also, reciprocals are stored for the diagonal components of  $U$ . (See Figure 3–2 in Section 3.2.2).
- (d) Shared memory parallel function for single-precision is not supported.

(7) Example

(a) Problem

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 5 + 8i & 7 + i & 6 + 3i & 1 + 2i \\ 1 + i & 9 + 5i & 4 + i & 5 \\ 4i & 3 + 3i & 4 + 2i & 6 + 9i \\ 7 + 8i & 6 & 7 + 6i & 10 + 4i \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 3 + 20i \\ -6 + 7i \\ -6i \\ 13i \end{bmatrix}$$

(b) Input data

Coefficient matrix  $A$ ,  $\text{lna} = 11$ ,  $n = 4$ , constant vector  $\mathbf{b}$  and  $\text{nt}=2$ .

(c) Main program

```

/*      C interface example for ASL_hbgns1 */

#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lna;
    int n;
    double _Complex *b;
    int *ipvt;
    int nt = 2;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "hbgns1.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hbgns1 ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &lna );
    fscanf( fp, "%d", &n );

    a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lna*n) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * n ));
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
    if( ipvt == NULL )
    {
        printf( "no enough memory for array ipvt\n" );
        return -1;
    }

    printf( "\t n = %6d\n", n );
    printf( "\n\tCoefficient Matrix   (Real, Imaginary)\n\n");
    for( i=0 ; i<n ; i++ )
    {
        for( j=0 ; j<n ; j++ )
        {
            double tmp_re;
            fscanf( fp, "%lf", &tmp_re );
            a[i+lna*j] = tmp_re;
        }
    }
    for( i=0 ; i<n ; i++ )
    {
        for( j=0 ; j<n ; j++ )
        {
            double tmp_im;
            fscanf( fp, "%lf", &tmp_im );
            a[i+lna*j] = a[i+lna*j] + tmp_im * _Complex_I;
        }
    }
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<n ; j++ )
        {
            printf( "(%8.3g , %8.3g) ", creal(a[i+lna*j]),cimag(a[i+lna*j]) );
        }
        printf( "\n" );
    }

    for( i=0 ; i<n ; i++ )
    {
        double tmp_re;
        fscanf( fp, "%lf", &tmp_re );
    }
}

```

```

        b[i] = tmp_re;
    }
    for( i=0 ; i<n ; i++ )
    {
        double tmp_im;
        fscanf( fp, "%lf", &tmp_im );
        b[i] = b[i] + tmp_im * _Complex_I;
    }

    printf( "\n\tConstant Vector (Real, Imaginary)\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t(%8.3g , %8.3g)\n", creal(b[i]), cimag(b[i]) );
    }

    fclose( fp );

    ierr = ASL_hbgns1(a, lna, n, b, ipvt, nt);

    printf( "\n      ** Output **\n\n" );
    printf( "\t(ierr = %6d)\n", ierr );

    printf( "\n\tSolution (Real, Imaginary)\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t x[%6d] =( %8.3g , %8.3g)\n", i, creal(b[i]), cimag(b[i]) );
    }

    free( a );
    free( b );
    free( ipvt );

    return 0;
}

```

(d) Output results

```

*** ASL_hbgns1 ***

** Input **

n =      4

Coefficient Matrix (Real, Imaginary)

(      5 ,      8) (      7 ,      1) (      6 ,      3) (      1 ,      2)
(      1 ,      1) (      9 ,      5) (      4 ,      1) (      5 ,      0)
(      0 ,      4) (      3 ,      3) (      4 ,      2) (      6 ,      9)
(      7 ,      8) (      6 ,      0) (      7 ,      6) (     10 ,      4)

Constant Vector (Real, Imaginary)

(      3 ,      20)
(     -6 ,      7)
(      0 ,     -6)
(      0 ,     13)

** Output **

ierr =      0

Solution (Real, Imaginary)

x[      0] =(      1 ,      1)
x[      1] =( -1.67e-16 ,      1)
x[      2] =(      1 , -2.78e-16)
x[      3] =(     -1 ,     -1)

```

### 3.4.3 ASL\_hbgnu LU Decomposition of a Complex Matrix

(1) **Function**

ASL\_hbgnu uses the Gauss method or the Crout method to perform an LU decomposition of the complex matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

`ierr = ASL_hbgnu (a, lna, n, ipvt, nt);`

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int} \text{ as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	Z*	lna×n	Input	Complex matrix $A$ (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ . (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	n	Output	Pivoting information ipvt[i-1]: Number of the row exchanged with row i in the i-th processing step. (See Note (b))
5	nt	I	1	Input	Number of tasks to be generated
6	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$  and  $\text{nt} \geq 1$



(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	The contents of array a are unchanged.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
4000+i	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular.	

(6) **Notes**

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array a with a minus sign added to each element, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array a. Also, reciprocals are stored for the diagonal components of  $U$ . (See Fig. 3–2 in Section 3.2.2.)
- (b) This function performs partial pivoting. Pivoting information is stored in array ipvt for use by subsequent functions. If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in ipvt[ $i - 1$ ]. In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Shared memory parallel function for single-precision is not supported.

### 3.4.4 ASL\_hbgnlc

#### LU Decomposition and Condition Number of a Complex Matrix

(1) **Function**

ASL\_hbgnlc uses the Gauss method or the Crout method to perform an LU decomposition and obtain the condition number of the complex matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

```
ierr = ASL_hbgnlc (a, lna, n, ipvt, &cond, w1, nt);
```

Single precision:

Nothing

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	Z*	lna×n	Input	Complex matrix (two-dimensional array type)
				Output	Upper triangular matrix $U$ and lower triangular matrix $L$ when $A$ is decomposed into $A = LU$ (See Notes (a) and (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	n	Output	Pivoting information ipvt[i-1]: Number of the row exchanged with row i in the i-th processing step. (See Note (b))
5	cond	D*	1	Output	Reciprocal of the condition number
6	w1	Z*	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$  and  $\text{nt} \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	The contents of array a are unchanged. $\text{cond} \leftarrow 1.0$ is performed.
2100	There existed the diagonal element which was close to zero in the LU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
$4000+i$	The pivot became 0.0 in the $i$ -th processing step. $A$ is singular.	

(6) Notes

- (a) The unit lower triangular matrix  $L$  is stored in the lower triangular portion of array a with a minus sign added to each element, and the upper triangular matrix  $U$  is stored in the upper triangular portion. However, since the diagonal components of  $L$  always are 1.0, they are not stored in array a. Also, reciprocals are stored for the diagonal components of  $U$ . (See Fig. 3–2 in Section 3.2.2.)
- (b) This function performs partial pivoting. Pivoting information is stored in array ipvt for use by subsequent functions. If the pivot row in the  $i$ -th step is row  $j$  ( $i \leq j$ ), then  $j$  is stored in  $\text{ipvt}[i - 1]$ . In addition, among the column elements corresponding to row  $i$  and row  $j$  of matrix  $A$ , elements from column 1 to column  $n$  actually are exchanged at this time.
- (c) Although the condition number is defined by  $\|A\| \cdot \|A^{-1}\|$ , an approximate value is obtained by this function.
- (d) Shared memory parallel function for single-precision is not supported.

## 3.5 REAL SYMMETRIC MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE)

### 3.5.1 ASL\_qbpsl, ASL\_pbpsl

#### Simultaneous Linear Equations (Real Symmetric Matrix)

(1) **Function**

ASL\_qbpsl or ASL\_pbpsl uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the real symmetric matrix  $A$  (two-dimensional array type) (upper triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_qbpsl (a, lna, n, b, ipvt, wk, nt);

Single precision:

ierr = ASL\_pbpsl (a, lna, n, b, ipvt, wk, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	lna×n	Input	Coefficient matrix $A$ (real symmetric matrix, two-dimensional array type, upper triangular type)
				Output	Upper triangular matrix $L^T$ when $A$ is decomposed into $A = LDL^T$ (See Note (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	$\begin{cases} D* \\ R* \end{cases}$	n	Input	Constant vector $\mathbf{b}$
				Output	Solution $\mathbf{x}$
5	ipvt	I*	N	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (c))
6	wk	$\begin{cases} D* \\ R* \end{cases}$	n	Work	Work Area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

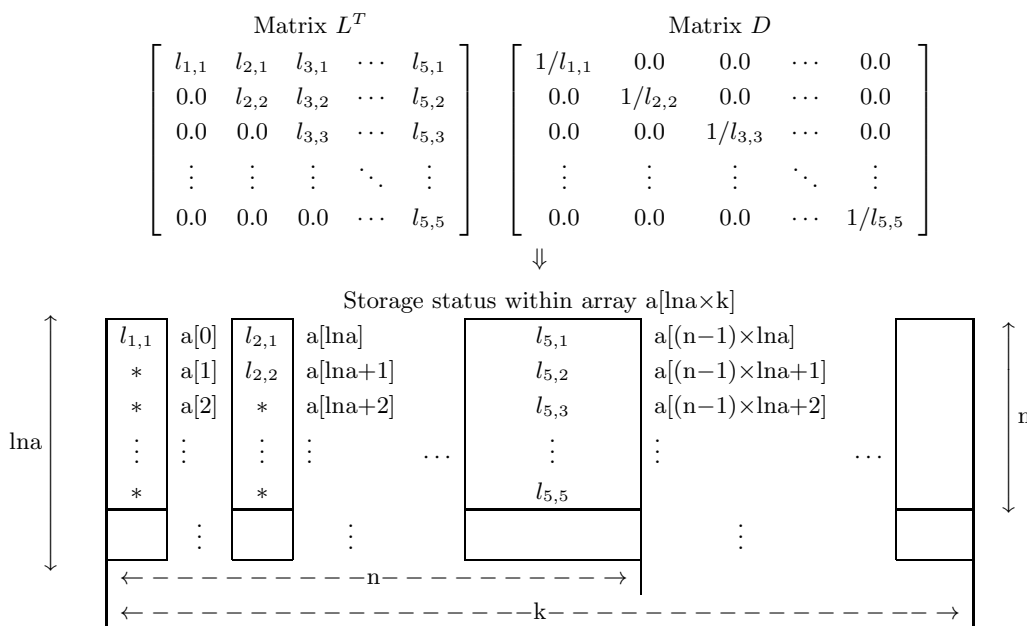
- (a)  $0 < n \leq \text{lna}$
- (b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the $\text{LDL}^T$ decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the $\text{LDL}^T$ decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) **Notes**

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector differs, call this function only once and then call function <Basic Functions Vol. 2> 2.6.4  $\left\{ \begin{array}{l} \text{ASL\_dbspls} \\ \text{ASL\_rbspls} \end{array} \right\}$  you to eliminate unnecessary calculations by performing the  $\text{LDL}^T$  decomposition of matrix  $A$  only once.
- (b) The upper triangular matrix  $L^T$  is stored in array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^T$ , they are not stored in array a. The matrix  $L$  is the transpose of matrix  $L^T$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of matrix  $L^T$  as components.  
 This function uses only the upper triangular portion of array a.



**Remarks**

- a.  $lna \geq n$  and  $n \leq k$  must hold.
- b. Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3–5 Storage Status of Matrix  $L^T$  and Contents of Matrix  $D$

- (c) This function performs partial pivoting when obtaining the  $LDL^T$  decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the  $i$ -th step is row(column)  $j$  ( $i < j$ ), then  $j$  is stored in  $ipvt[i-1]$ . In addition, among the column(row) elements corresponding to row(column)  $i$  and row(column)  $j$  of matrix  $A$ , elements from column(row)  $i$  to column(row)  $n$  actually are exchanged at this time.

**(7) Example**

- (a) Problem

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 5 & 4 & 1 & 1 \\ 4 & 5 & 1 & 1 \\ 1 & 1 & 4 & 2 \\ 1 & 1 & 2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 1 \\ -1 \\ 4 \\ -4 \end{bmatrix}$$

- (b) Input data

Coefficient matrix  $A$ ,  $lna = 11$ ,  $n = 4$  and constant vector  $\mathbf{b}$ .

- (c) Main Program

```

/*      C interface example for ASL_qbpsl */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int na;
    int n;
    double *b;
    int *ipvt;
    double *wk;
    int nt = 2;
    int ierr;

```

```
int i,j;
FILE *fp;
fp = fopen( "qbspsl.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_qbspsl ***\n" );
printf( "\n    ** Input **\n\n" );

fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );
a = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double * )malloc((size_t)( sizeof(double) * n ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( ipvt == NULL )
{
    printf( "no enough memory for array ipvt\n" );
    return -1;
}

wk = ( double * )malloc((size_t)( sizeof(double) * n ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\n\tCoefficient Matrix\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &a[i+na*j] );
        printf( "%8.3g ", a[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector\n\n" );
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &b[i] );
    printf( "\t%8.3g\n", b[i] );
}

fclose( fp );

ierr = ASL_qbspsl(a, na, n, b, ipvt, wk, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n\n", ierr );

printf( "\tSolution \n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = %8.3g\n", i, b[i] );
}

free( a );
free( b );
free( ipvt );
free( wk );
```

```
    return 0;  
}
```

(d) Output results

```
*** ASL_qbspsl ***  
** Input **  
n =      4  
Coefficient Matrix  
      5      4      1      1  
      4      5      1      1  
      1      1      4      2  
      1      1      2      4  
Constant Vector  
      1  
     -1  
      4  
     -4  
** Output **  
ierr =      0  
Solution  
x[  0] =      1  
x[  1] =     -1  
x[  2] =      2  
x[  3] =     -2
```



### 3.5.2 ASL\_qbspud, ASL\_pbspud LDL<sup>T</sup> Decomposition of a Real Symmetric Matrix

(1) **Function**

ASL\_qbspud or ASL\_pbspud uses the modified Cholesky method to perform an LDL<sup>T</sup> decomposition of the real symmetric matrix  $A$  (two-dimensional array type).

(2) **Usage**

Double precision:

ierr = ASL\_qbspud (a, lna, n, ipvt, wk, nt);

Single precision:

ierr = ASL\_pbspud (a, lna, n, ipvt, wk, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} \text{D*} \\ \text{R*} \end{cases}$	lna × n	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Upper triangular matrix $L^T$ when $A$ is decomposed into $A = LDL^T$ (See Note (a))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	N	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (b))
5	wk	$\begin{cases} \text{D*} \\ \text{R*} \end{cases}$	n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of array a are not changed.
2100	There existed the diagonal element which was close to zero in the LDL <sup>T</sup> decomposition of the coefficient matrix <i>A</i> . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000+ <i>i</i>	The pivot became 0.0 in the <i>i</i> -th processing step of the LDL <sup>T</sup> decomposition of coefficient matrix <i>A</i> . <i>A</i> is singular.	

(6) **Notes**

- (a) The upper triangular matrix  $L^T$  is stored in array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^T$ , they are not stored in array a. (See Section 3.5.1, Figure 3–5.)
- (b) This function performs partial pivoting when obtaining the LDL<sup>T</sup> decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the *i*-th step is row(column) *j* ( $i < j$ ), then *j* is stored in `ipvt[i-1]`. In addition, among the column(row) elements corresponding to row(column) *i* and row(column) *j* of matrix  $A$ , elements from column(row) *i* to column(row) *n* actually are exchanged at this time.

## 3.6 REAL SYMMETRIC MATRIX (TWO-DIMENSIONAL ARRAY TYPE, LOWER TRIANGULAR TYPE) (NO PIVOTING)

### 3.6.1 ASL\_qbsnsl, ASL\_pbsnsl

#### Simultaneous Linear Equations (Real Symmetric Matrix) (No Pivoting)

(1) **Function**

ASL\_qbsnsl or ASL\_pbsnsl uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the real symmetric matrix  $A$  (two-dimensional array type, lower triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_qbsnsl (a, lna, n, b, nt);
```

Single precision:

```
ierr = ASL_pbsnsl (a, lna, n, b, nt);
```

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	lna×n	Input	Coefficient matrix $A$ (real symmetric matrix, two-dimensional array type, lower triangular type)
				Output	lower triangular matrix $U^T$ when $A$ is decomposed into $A = U^T D U$ (See Note (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	$\begin{cases} D* \\ R* \end{cases}$	n	Input	Constant vector $\mathbf{b}$
				Output	Solution $\mathbf{x}$
5	nt	I	1	Input	Number of tasks to be generated
6	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

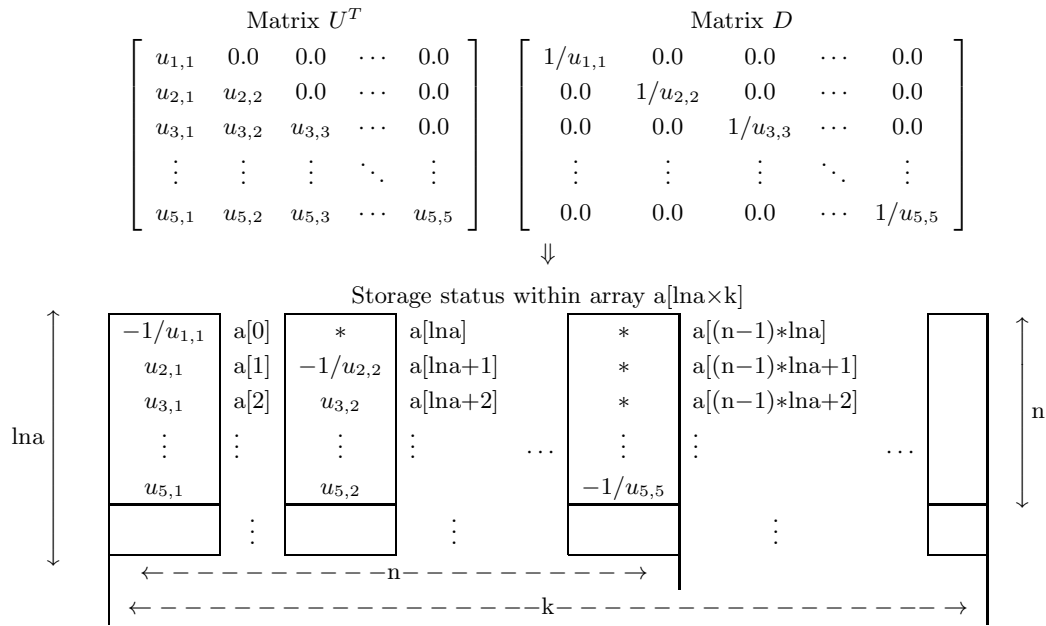
(b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the $U^T D U$ decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the $U^T D U$ decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) **Notes**

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector differs, call this function only once and then call function <Basic Functions Vol. 2> 2.8.3  $\left\{ \begin{matrix} \text{ASL\_dbsnsl} \\ \text{ASL\_rbsnsl} \end{matrix} \right\}$  you to eliminate unnecessary calculations by performing the  $U^T D U$  decomposition of matrix  $A$  only once.
- (b) The lower triangular matrix  $U^T$  is stored in array a. For the diagonal components of  $U^T$ , their reciprocals are stored in array a with the sign changed. Since the diagonal matrix  $D$  and the upper triangular matrix  $U$  are calculated from  $U^T$ , they are not stored in array a. The matrix  $U$  is the transpose of matrix  $U^T$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of matrix  $U^T$  as components.
- This function uses only the lower triangular portion of array a.



- Remarks**
- lma ≥ n and n ≤ k must hold.
  - Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3-6 Storage Status of Matrix  $U^T$  and Contents of Matrix  $D$

(7) Example

- (a) Problem  
 Solve the following simultaneous linear equations.

$$\begin{bmatrix} 5 & 4 & 1 & 1 \\ 4 & 5 & 1 & 1 \\ 1 & 1 & 4 & 2 \\ 1 & 1 & 2 & 4 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 1 \\ -1 \\ 4 \\ -4 \end{bmatrix}$$

- (b) Input data  
 Coefficient matrix  $A$ , lma = 11, n = 4 and constant vector  $b$ .
- (c) Main Program

```

/* C interface example for ASL_qbsnsl */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int na;
    int n;
    double *b;
    int ierr;

    int i,j;
    int nt = 2;

    FILE *fp;

    fp = fopen( "qbsnsl.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }
}
    
```

```

printf( "      *** ASL_qbsnsl ***\n" );
printf( "\n      ** Input **\n\n" );

fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );
a = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double * )malloc((size_t)( sizeof(double) * n ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\n\tCoefficient Matrix\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &a[i+na*j] );
        printf( "%8.3g ", a[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector\n\n" );
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &b[i] );
    printf( "\t%8.3g\n", b[i] );
}

fclose( fp );

ierr = ASL_qbsnsl(a, na, n, b, nt);

printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n\n", ierr );

printf( "\tSolution \n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = %8.3g\n", i, b[i] );
}

free( a );
free( b );

return 0;
}

```

(d) Output results

```

*** ASL_qbsnsl ***

** Input **

n =      4

Coefficient Matrix

      5      4      1      1
      4      5      1      1
      1      1      4      2
      1      1      2      4

Constant Vector

      1
     -1
      4
     -4

** Output **

ierr =      0

Solution

```

$$\begin{array}{rcl} x[0] & = & 1 \\ x[1] & = & -1 \\ x[2] & = & 2 \\ x[3] & = & -2 \end{array}$$

### 3.6.2 ASL\_qbsnud, ASL\_pbsnud

#### $U^T$ DU Decomposition of a Real Symmetric Matrix (No Pivoting)

(1) **Function**

ASL\_qbsnud or ASL\_pbsnud uses the modified Cholesky method to perform an  $U^T$ DU decomposition of the real symmetric matrix  $A$  (two-dimensional array type) (lower triangular type).

(2) **Usage**

Double precision:

ierr = ASL\_qbsnud (a, lna, n, nt);

Single precision:

ierr = ASL\_pbsnud (a, lna, n, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	lna × n	Input	Real symmetric matrix $A$ (two-dimensional array type) (lower triangular type)
				Output	Lower triangular matrix $U^T$ when $A$ is decomposed into $A = U^T D U$ (See Note (a))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	nt	I	1	Input	Number of tasks to be generated
5	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq \text{lna}$

(b)  $\text{nt} \geq 1$



(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of array a are not changed.
2100	There existed the diagonal element which was close to zero in the $U^T$ DU decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000+i	The pivot became 0.0 in the $i$ -th processing step of the $U^T$ DU decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) Notes

- (a) The lower triangular matrix  $U^T$  is stored in array a. For the diagonal components of  $U^T$ , their reciprocals are stored in array a with the sign changed. Since the diagonal matrix  $D$  and the upper triangular matrix  $U$  are calculated from  $U^T$ , they are not stored in array a. (See Section 3.6.1, Figure 3–6.)

---

## 3.7 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE)

### 3.7.1 ASL\_hbhpsl, ASL\_gbhpsl

#### Simultaneous Linear Equations (Hermitian Matrix)

(1) **Function**

ASL\_hbhpsl or ASL\_gbhpsl uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having a Hermitian matrix (two-dimensional array type) (upper triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_hbhpsl (ar, ai, lna, n, br, bi, ipvt, w1, nt);
```

Single precision:

```
ierr = ASL_gbhpsl (ar, ai, lna, n, br, bi, ipvt, w1, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna × n	Input	Real part of coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Real part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna × n	Input	Imaginary part of coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Imaginary part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Real part of constant vector $\mathbf{b}$
				Output	Real part of solution $\mathbf{x}$
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Imaginary part of constant vector $\mathbf{b}$
				Output	Imaginary part of solution $\mathbf{x}$
7	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (c))
8	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
9	nt	I	1	Input	Number of tasks to be generated
10	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

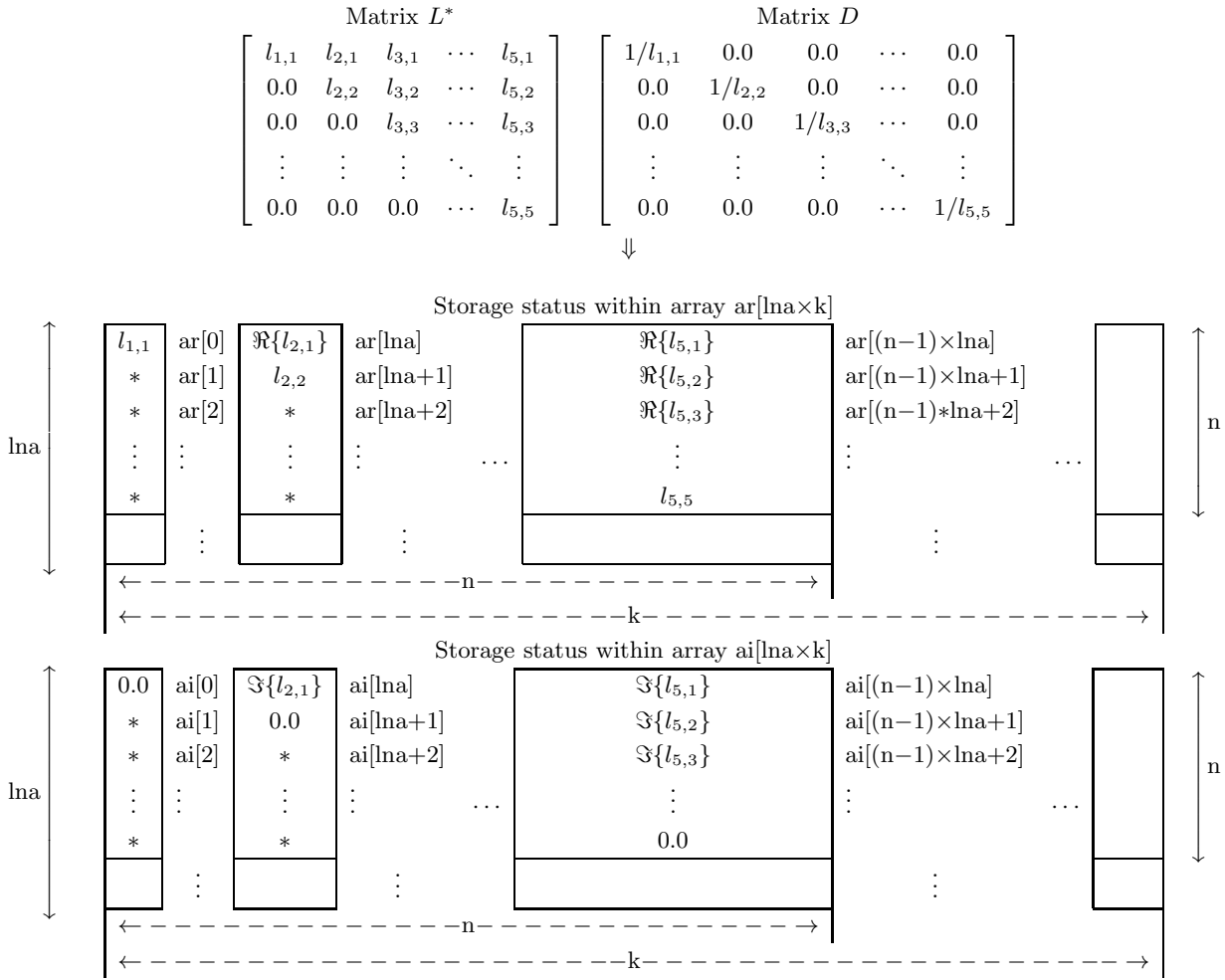
- (a)  $0 < n \leq \text{lna}$
- (b)  $\text{nt} \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$ .	Contents of arrays ar and ai are not changed. $b[0] \leftarrow b[0]/ar[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LDL* decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) Notes

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, call this function only once and then call function <Basic Functions Vol. 2> 2.9.4  $\left\{ \begin{array}{l} \text{ASL\_zbhpls} \\ \text{ASL\_cbhpls} \end{array} \right\}$  the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculation by performing the LDL\* decomposition of matrix  $A$  only once.
- (b) The upper triangular matrix  $L^*$  is stored in the upper triangular portions of arrays ar and ai. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in arrays ar and ai. The matrix  $L$  is the adjoint matrix of the matrix  $L^*$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of the matrix  $L^*$  as its components. This function uses only the upper triangular portions of arrays ar and ai.



**Remarks**

- a.  $l_{na} \geq n$  and  $n \leq k$  must hold.
- b. Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3–7 Storage Status of Matrix  $L^*$  and Contents of Matrix  $D$

(c) This function performs partial pivoting when obtaining the LDL\* decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the  $i$ -th step is row(column)  $j$  ( $i < j$ ), then  $j$  is stored in  $ipvt[i-1]$ . In addition, among the column(row) elements corresponding to row(column)  $i$  and row(column)  $j$  of matrix  $A$ , elements from column(row)  $i$  to column(row)  $n$  actually are exchanged at this time.

**(7) Example**

(a) Problem

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 9 & 7+3i & 2+5i & 1+i \\ 7-3i & 10 & 3+2i & 2+4i \\ 2-5i & 3-2i & 8 & 5+i \\ 1-i & 2-4i & 5-i & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 10+6i \\ 11+2i \\ 4+6i \\ 4+6i \end{bmatrix}$$

(b) Input data

Coefficient matrix real part ar and Imaginary part ai,  $l_{na} = 11$ ,  $n = 4$  and constant vector b.

(c) Main Program

```

/*      C interface example for ASL_hbhpsl */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int na;
    int n;
    double *br;
    double *bi;
    int *ipvt;
    double *w1;
    int nt = 2;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "hbhpsl.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hbhpsl ***\n" );
    printf( "\n      ** Input **\n\n" );
    fscanf( fp, "%d", &na );
    fscanf( fp, "%d", &n );

    ar = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
    if( ar == NULL )
    {
        printf( "no enough memory for array ar\n" );
        return -1;
    }

    ai = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
    if( ai == NULL )
    {
        printf( "no enough memory for array ai\n" );
        return -1;
    }

    br = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( br == NULL )
    {
        printf( "no enough memory for array br\n" );
        return -1;
    }

    bi = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( bi == NULL )
    {
        printf( "no enough memory for array bi\n" );
        return -1;
    }

    ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
    if( ipvt == NULL )
    {
        printf( "no enough memory for array ipvt\n" );
        return -1;
    }

    w1 = ( double * )malloc((size_t)( sizeof(double) * (n*2) ));
    if( w1 == NULL )
    {
        printf( "no enough memory for array w1\n" );
        return -1;
    }

    printf( "\t n = %6d\n\n", n );
    printf( "\tCoefficient Matrix (Real, Imaginary)\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        for( j=0 ; j<n ; j++ )
        {
            fscanf( fp, "%lf", &ar[i+na*j] );
        }
    }
    for( i=0 ; i<n ; i++ )
    {
        for( j=0 ; j<n ; j++ )
        {
            fscanf( fp, "%lf", &ai[i+na*j] );
        }
    }
}

```

```

    }
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t" );
        for( j=0 ; j<i ; j++ )
        {
            printf( "          " );
        }
        for( j=i ; j<n ; j++ )
        {
            printf( "(%8.3g , %8.3g) ", ar[i+na*j],ai[i+na*j] );
        }
        printf( "\n" );
    }

    printf( "\n\tConstant Vector (Real, Imaginary)\n\n");
    for( i=0 ; i<n ; i++ )
    {
        fscanf( fp, "%lf", &br[i] );
    }
    for( i=0 ; i<n ; i++ )
    {
        fscanf( fp, "%lf", &bi[i] );
    }
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t(%8.3g , %8.3g) \n", br[i],bi[i] );
    }

    fclose( fp );

    ierr = ASL_hbhpsl(ar, ai, na, n, br, bi, ipvt, w1, nt);

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\n\tSolution (Real, Imaginary)\n\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t x[%6d] = (%8.3g , %8.3g)\n", i, br[i],bi[i] );
    }

    free( ar );
    free( ai );
    free( br );
    free( bi );
    free( ipvt );
    free( w1 );

    return 0;
}

```

(d) Output results

```

*** ASL_hbhpsl ***

** Input **

n =      4

Coefficient Matrix (Real, Imaginary)

(      9 ,      0) (      7 ,      3) (      2 ,      5) (      1 ,      1)
(      0 ,     10) (     10 ,      0) (      3 ,      2) (      2 ,      4)
(      0 ,      0) (      0 ,      8) (      8 ,      0) (      5 ,      1)
(      0 ,      0) (      0 ,      0) (      0 ,      6) (      6 ,      0)

Constant Vector (Real, Imaginary)

(     10 ,      6)
(     11 ,      2)
(      4 ,      6)
(      4 ,      6)

** Output **

ierr =      0

Solution (Real, Imaginary)

x[      0] = (      1 ,      0)
x[      1] = (      1 ,      0)
x[      2] = (-6.63e-17 ,      1)
x[      3] = (2.09e-17 ,      1)

```

### 3.7.2 ASL\_hbhpud, ASL\_gbhpud LDL\* Decomposition of a Hermitian Matrix

(1) **Function**

ASL\_hbhpud or ASL\_gbhpud uses the modified Cholesky method to perform an LDL\* decomposition of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type).

(2) **Usage**

Double precision:

ierr = ASL\_hbhpud (ar, ai, lna, n, ipvt, w1, nt);

Single precision:

ierr = ASL\_gbhpud (ar, ai, lna, n, ipvt, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Real part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
2	ai	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Imaginary part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
3	lna	I	1	Input	Adjustable dimension of array ar and ai
4	n	I	1	Input	Order of matrix $A$
5	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (b))
6	w1	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$2 \times n$	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)



(4) **Restrictions**

- (a)  $0 < n \leq \ln a$
- (b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of arrays ar and ai are not changed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	The pivot became 0.0 in the $i$ -th processing step of the LDL* decomposition of coefficient matrix $A$ . $A$ is singular.	

(6) **Notes**

- (a) The upper triangular matrix  $L^*$  is stored in the upper triangular portions of arrays ar and ai. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in arrays ar and ai. This function uses only the upper triangular portions of arrays ar and ai. (See Sections 3.7.1 Figure 3–7.)
- (b) This function performs partial pivoting when obtaining the LDL\* decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the  $i$ -th step is row(column)  $j$  ( $i < j$ ), then  $j$  is stored in  $ipvt[i-1]$ . In addition, among the column(row) elements corresponding to row(column)  $i$  and row(column)  $j$  of matrix  $A$ , elements from column(row)  $i$  to column(row)  $n$  actually are exchanged at this time.

---

### 3.8 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE) (NO PIVOTING)

#### 3.8.1 ASL\_hbhrs1, ASL\_gbhrs1 Simultaneous Linear Equations (Hermitian Matrix) (No Pivoting)

(1) **Function**

ASL\_hbhrs1 or ASL\_gbhrs1 uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having a Hermitian matrix (two-dimensional array type) (upper triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_hbhrs1 (ar, ai, lna, n, br, bi, w1, nt);
```

Single precision:

```
ierr = ASL_gbhrs1 (ar, ai, lna, n, br, bi, w1, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna×n	Input	Real part of coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Real part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna×n	Input	Imaginary part of coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Imaginary part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Real part of constant vector $\mathbf{b}$
				Output	Real part of solution $\mathbf{x}$
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Imaginary part of constant vector $\mathbf{b}$
				Output	Imaginary part of solution $\mathbf{x}$
7	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

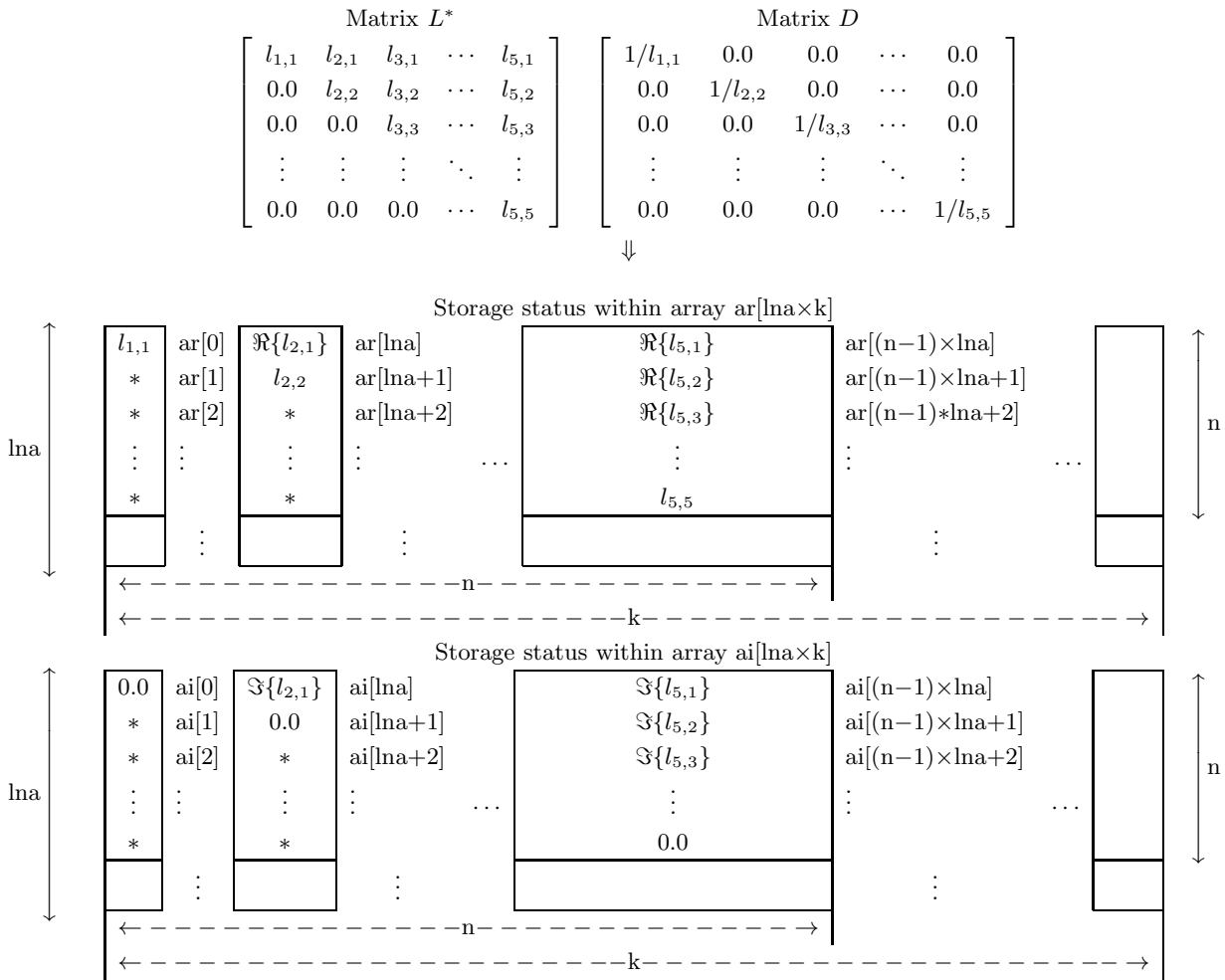
- (a)  $0 < n \leq \text{lna}$
- (b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$ .	Contents of arrays ar and ai are not changed. $b[0] \leftarrow b[0]/ar[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	A diagonal element became equal to 0.0 in the $i$ -th processing step of the LDL* decomposition of coefficient matrix $A$ . $A$ is nearly singular.	

(6) **Notes**

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, call this function only once and then call function <Basic Functions Vol. 2> 2.10.4  $\left\{ \begin{array}{l} \text{ASL_zbhrs1} \\ \text{ASL_cbhrs1} \end{array} \right\}$  the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculation by performing the LDL\* decomposition of matrix  $A$  only once.
- (b) The upper triangular matrix  $L^*$  is stored in the upper triangular portions of arrays ar and ai. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in arrays ar and ai. The matrix  $L$  is the adjoint matrix of the matrix  $L^*$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of the matrix  $L^*$  as its components. This function uses only the upper triangular portions of arrays ar and ai. (See Fig. 3–8.)



**Remarks**

- a.  $l_{na} \geq n$  and  $n \leq k$  must hold.
- b. Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3–8 Storage Status of Matrix  $L^*$  and Contents of Matrix  $D$

(7) **Example**

(a) **Problem**

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 9 & 7 + 3i & 2 + 5i & 1 + i \\ 7 - 3i & 10 & 3 + 2i & 2 + 4i \\ 2 - 5i & 3 - 2i & 8 & 5 + i \\ 1 - i & 2 - 4i & 5 - i & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 10 + 6i \\ 11 + 2i \\ 4 + 6i \\ 4 + 6i \end{bmatrix}$$

(b) **Input data**

Coefficient matrix real part ar and Imaginary part ai,  $l_{na} = 11, n = 4$  and constant vector b.

(c) **Main Program**

```

/*      C interface example for ASL_hbhrs1 */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()

```

```

{
double *ar;
double *ai;
int na;
int n;
double *br;
double *bi;
double *w1;
int nt = 2;
int ierr;
int i,j;
FILE *fp;

fp = fopen( "hbhrs1.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_hbhrs1 ***\n" );
printf( "\n    ** Input **\n\n" );
fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );

ar = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( ar == NULL )
{
    printf( "no enough memory for array ar\n" );
    return -1;
}

ai = ( double * )malloc((size_t)( sizeof(double) * (na*n) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

br = ( double * )malloc((size_t)( sizeof(double) * n ));
if( br == NULL )
{
    printf( "no enough memory for array br\n" );
    return -1;
}

bi = ( double * )malloc((size_t)( sizeof(double) * n ));
if( bi == NULL )
{
    printf( "no enough memory for array bi\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * (n*2) ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\t n = %6d\n\n", n );
printf( "\tCoefficient Matrix (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ar[i+na*j] );
    }
}
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ai[i+na*j] );
    }
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "          " );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", ar[i+na*j],ai[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector (Real, Imaginary)\n\n");

```

```

for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &br[i] );
}
for( i=0 ; i<n ; i++ )
{
    fscanf( fp, "%lf", &bi[i] );
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t(%8.3g , %8.3g) \n", br[i],bi[i] );
}
fclose( fp );

ierr = ASL_hbhrsl(ar, ai, na, n, br, bi, w1, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = (%8.3g , %8.3g)\n", i, br[i],bi[i] );
}

free( ar );
free( ai );
free( br );
free( bi );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_hbhrsl ***

** Input **

n =      4

Coefficient Matrix (Real, Imaginary)

(      9 ,      0) (      7 ,      3) (      2 ,      5) (      1 ,      1)
(      0 ,      0) (     10 ,      0) (      3 ,      2) (      2 ,      4)
(      0 ,      0) (      0 ,      0) (      8 ,      0) (      5 ,      1)
(      0 ,      0) (      0 ,      0) (      0 ,      6) (      6 ,      0)

Constant Vector (Real, Imaginary)

(     10 ,      6)
(     11 ,      2)
(      4 ,      6)
(      4 ,      6)

** Output **

ierr =      0

Solution (Real, Imaginary)

x[      0] = (      1 , 1.48e-16)
x[      1] = (      1 , -3.9e-17)
x[      2] = (-1.02e-16 ,      1)
x[      3] = (8.34e-17 ,      1)

```

### 3.8.2 ASL\_hbhrud, ASL\_gbhrud

#### LDL\* Decomposition of a Hermitian Matrix (No Pivoting)

(1) **Function**

ASL\_hbhrud or ASL\_gbhrud uses the modified Cholesky method to perform an LDL\* decomposition of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type).

(2) **Usage**

Double precision:

ierr = ASL\_hbhrud (ar, ai, lna, n, w1, nt);

Single precision:

ierr = ASL\_gbhrud (ar, ai, lna, n, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$l_n \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Real part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$l_n \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Imaginary part of upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
3	lna	I	1	Input	Adjustable dimension of array ar and ai
4	n	I	1	Input	Order of matrix $A$
5	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq l_n$

(b)  $nt \geq 1$



(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of arrays ar and ai are not changed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000+i	A diagonal element became equal to 0.0 in the $i$ -th processing step. $A$ is nearly singular.	

(6) Notes

- (a) The upper triangular matrix  $L^*$  is stored in the upper triangular portions of arrays ar and ai. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in arrays ar and ai. This function uses only the upper triangular portions of arrays ar and ai. (See Fig. 3–8 in Section 3.8.1.)

### 3.9 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (COMPLEX ARGUMENT TYPE)

#### 3.9.1 ASL\_hbhfs1, ASL\_gbhfs1 Simultaneous Linear Equations (Hermitian Matrix)

(1) **Function**

ASL\_hbhfs1 or ASL\_gbhfs1 uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_hbhfs1 (a, lna, n, b, ipvt, w1, nt);

Single precision:

ierr = ASL\_gbhfs1 (a, lna, n, b, ipvt, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$lna \times n$	Input	Coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	n	Input	Constant vector $\mathbf{b}$
				Output	Solution $\mathbf{x}$
5	ipvt	$I^*$	n	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (c))
6	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

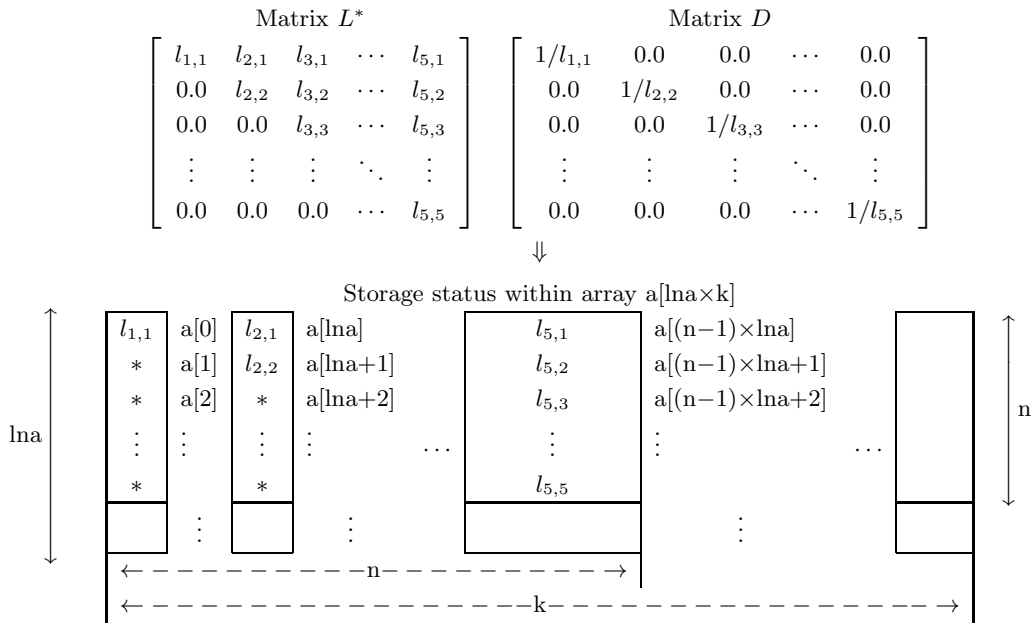
- (a)  $0 < n \leq \text{lna}$
- (b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$ .	Contents of array a are not changed. $b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix A. The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	A diagonal element became equal to 0.0 in the $i$ -th processing step of the LDL* decomposition of coefficient matrix A. A is nearly singular.	

(6) **Notes**

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, call this function only once and then call function <Basic Functions Vol. 2> 2.11.4  $\left\{ \begin{array}{l} \text{ASL_zbhfls} \\ \text{ASL_cbhfls} \end{array} \right\}$  the required number of times varying only the contents of b. This enables you to eliminate unnecessary calculations by performing the LDL\* decomposition of matrix A only once.
- (b) The upper triangular matrix  $L^*$  is stored in the upper triangular portion of array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in array a. The matrix  $L$  is the adjoint matrix of the matrix  $L^*$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of the matrix  $L^*$  as its components.



**Remarks**

- a.  $\text{lna} \geq n$  and  $n \leq k$  must hold.
- b. Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3–9 Storage Status of Matrix  $L^*$  and Contents of Matrix  $D$

(c) This function performs partial pivoting when obtaining the LDL\* decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the  $i$ -th step is row(column)  $j$  ( $i < j$ ), then  $j$  is stored in  $\text{ipvt}[i-1]$ . In addition, among the column(row) elements corresponding to row(column)  $i$  and row(column)  $j$  of matrix  $A$ , elements from column(row)  $i$  to column(row)  $n$  actually are exchanged at this time.

**(7) Example**

(a) Problem

Solve the following simultaneous linear equations.

$$\begin{bmatrix} 9 & 7+3i & 2+5i & 1+i \\ 7-3i & 10 & 3+2i & 2+4i \\ 2-5i & 3-2i & 8 & 5+i \\ 1-i & 2-4i & 5-i & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 10+6i \\ 11+2i \\ 4+6i \\ 4+6i \end{bmatrix}$$

(b) Input data

Coefficient matrix  $A$ ,  $\text{lna} = 11$ ,  $n = 4$  and constant vector  $\mathbf{b}$ .

(c) Main Program

```

/*      C interface example for ASL_hbhfs1 */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int na;
    int n;
    double _Complex *b;
    int *ipvt;
    double *w1;
    int nt=2;

```

```

int ierr;
int i,j;
FILE *fp;

fp = fopen( "hbhfs1.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_hbhfs1 ***\n" );
printf( "\n        ** Input **\n\n" );

fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (na*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * n ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

ipvt = ( int * )malloc((size_t)( sizeof(int) * n ));
if( ipvt == NULL )
{
    printf( "no enough memory for array ipvt\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * n ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\n\tCoefficient Matrix (Real, Imaginary)\n\n");
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        double tmp_re;
        fscanf( fp, "%lf", &tmp_re );
        a[i+na*j] = tmp_re;
    }
}
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        double tmp_im;
        fscanf( fp, "%lf", &tmp_im );
        a[i+na*j] = a[i+na*j] + tmp_im * _Complex_I;
    }
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "                " );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", creal(a[i+na*j]),cimag(a[i+na*j]) );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector (Real, Imaginary)\n\n");
for( i=0 ; i<n ; i++ )
{
    double tmp_re;
    fscanf( fp, "%lf", &tmp_re );
    b[i] = tmp_re;
}
for( i=0 ; i<n ; i++ )
{

```

```

    double tmp_im;
    fscanf( fp, "%lf", &tmp_im );
    b[i] = b[i] + tmp_im * _Complex_I;
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t(%8.3g , %8.3g)\n", creal(b[i]),cimag(b[i]) );
}
fclose( fp );

ierr = ASL_hbhfsl(a, na, n, b, ipvt, w1, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tSolution (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = (%8.3g , %8.3g)\n", i,creal(b[i]),cimag(b[i]));
}

free( a );
free( b );
free( ipvt );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_hbhfsl ***

** Input **

n =      4

Coefficient Matrix (Real, Imaginary)
(      9 ,      0) (      7 ,      3) (      2 ,      5) (      1 ,      1)
                   (      10 ,      0) (      3 ,      2) (      2 ,      4)
                   (      8 ,      0) (      8 ,      0) (      5 ,      1)
                   (      6 ,      0)

Constant Vector (Real, Imaginary)
(      10 ,      6)
(      11 ,      2)
(      4 ,      6)
(      4 ,      6)

** Output **

ierr =      0

Solution (Real, Imaginary)
x[      0] = (      1 ,      0)
x[      1] = (      1 ,      0)
x[      2] = (-6.63e-17 ,      1)
x[      3] = (2.09e-17 ,      1)

```

### 3.9.2 ASL\_hbhfd, ASL\_gbhfd LDL\* Decomposition of a Hermitian Matrix

(1) **Function**

ASL\_hbhfd or ASL\_gbhfd uses the modified Cholesky method to perform an LDL\* decomposition of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type).

(2) **Usage**

Double precision:

`ierr = ASL_hbhfd (a, lna, n, ipvt, w1, nt);`

Single precision:

`ierr = ASL_gbhfd (a, lna, n, ipvt, w1, nt);`

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$lna \times n$	Input	Hermitian matrix $A$ (two-dimensional array type, upper triangular type)
				Output	Upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	ipvt	I*	n	Output	Pivoting information ipvt[i - 1]: Number of the row(column) exchanged with row(column) i in the i-th processing step. (See Note (b))
5	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of array a are not changed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000+i	A diagonal element became equal to 0.0 in the $i$ -th processing step. $A$ is nearly singular.	

(6) **Notes**

- (a) The upper triangular matrix  $L^*$  is stored in the upper triangular portion of array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in array a. This function uses only the upper triangular portion of array a. (See Fig. 3–9 in Section 3.9.1)
- (b) This function performs partial pivoting when obtaining the LDL\* decomposition of coefficient matrix  $A$ . The permutation of rows and columns is symmetrical for row and column. If the pivot row(column) in the  $i$ -th step is row(column)  $j$  ( $i < j$ ), then  $j$  is stored in `ipvt[i-1]`. In addition, among the column(row) elements corresponding to row(column)  $i$  and row(column)  $j$  of matrix  $A$ , elements from column(row)  $i$  to column(row)  $n$  actually are exchanged at this time.



### 3.10 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (COMPLEX ARGUMENT TYPE) (NO PIVOTING)

#### 3.10.1 ASL\_hbhesl, ASL\_gbhesl

##### Simultaneous Linear Equations (Hermitian Matrix) (No Pivoting)

(1) **Function**

ASL\_hbhesl or ASL\_gbhesl uses the modified Cholesky method to solve the simultaneous linear equations  $A\mathbf{x} = \mathbf{b}$  having the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) as coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_hbhesl (a, lna, n, b, nt);

Single precision:

ierr = ASL\_gbhesl (a, lna, n, b, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} Z^* \\ C^* \end{cases}$	$lna \times n$	Input	Coefficient matrix $A$ (Hermitian matrix, two-dimensional array type, upper triangular type)
				Output	Upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (b))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	b	$\begin{cases} Z^* \\ C^* \end{cases}$	n	Input	Constant vector $\mathbf{b}$
				Output	Solution $\mathbf{x}$
5	nt	I	1	Input	Number of tasks to be generated
6	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

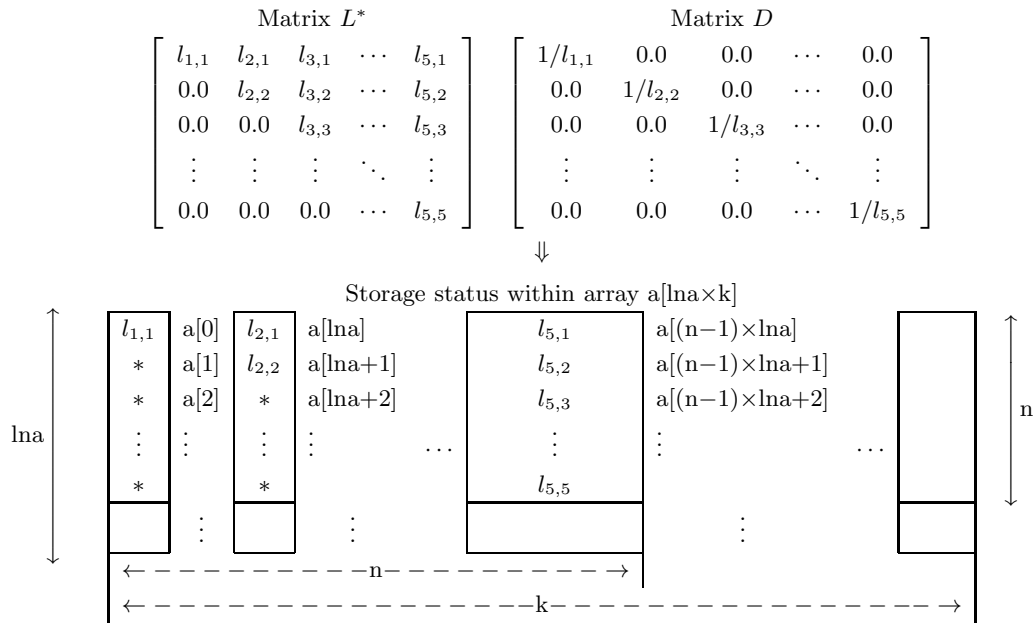
(b)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n = 1$ .	Contents of array a are not changed. $b[0] \leftarrow b[0]/a[0]$ is performed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$4000+i$	A diagonal element became equal to 0.0 in the $i$ -th processing step of the LDL* decomposition of coefficient matrix $A$ . $A$ is nearly singular.	

(6) Notes

- (a) To solve multiple sets of simultaneous linear equations where only the constant vector  $\mathbf{b}$  differs, call this function only once and then call function <Basic Functions Vol. 2> 2.12.4  $\left\{ \begin{matrix} \text{ASL\_zbhsl} \\ \text{ASL\_cbhsl} \end{matrix} \right\}$  the required number of times varying only the contents of  $\mathbf{b}$ . This enables you to eliminate unnecessary calculations by performing the LDL\* decomposition of matrix  $A$  only once.
- (b) The upper triangular matrix  $L^*$  is stored in the upper triangular portion of array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in array a. The matrix  $L$  is the adjoint matrix of the matrix  $L^*$ , and the matrix  $D$  is a diagonal matrix having the reciprocals of the diagonal elements of the matrix  $L^*$  as its components.



- Remarks**
- a. l<sub>na</sub> ≥ n and n ≤ k must hold.
  - b. Input time values of elements indicated by asterisks (\*) are not guaranteed.

Figure 3-10 Storage Status of Matrix  $L^*$  and Contents of Matrix  $D$

(7) Example

(a) Problem  
 Solve the following simultaneous linear equations.

$$\begin{bmatrix} 9 & 7+3i & 2+5i & 1+i \\ 7-3i & 10 & 3+2i & 2+4i \\ 2-5i & 3-2i & 8 & 5+i \\ 1-i & 2-4i & 5-i & 6 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} 10+6i \\ 11+2i \\ 4+6i \\ 4+6i \end{bmatrix}$$

(b) Input data  
 Coefficient matrix  $A$ , l<sub>na</sub> = 11, n = 4 and constant vector  $b$ .

(c) Main Program

```

/*      C interface example for ASL_hbhesl */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int na;
    int n;
    double _Complex *b;
    int nt = 2;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "hbhesl.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hbhesl ***\n" );
}
    
```

```

printf( "\n      ** Input **\n\n" );
fscanf( fp, "%d", &na );
fscanf( fp, "%d", &n );

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (na*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * n ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

printf( "\t n = %6d\n", n );
printf( "\n\tCoefficient Matrix (Real, Imaginary)\n\n");
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        double tmp_re;
        fscanf( fp, "%lf", &tmp_re );
        a[i+na*j] = tmp_re;
    }
}
for( i=0 ; i<n ; i++ )
{
    for( j=0 ; j<n ; j++ )
    {
        double tmp_im;
        fscanf( fp, "%lf", &tmp_im );
        a[i+na*j] = a[i+na*j] + tmp_im * _Complex_I;
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "          " );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", creal(a[i+na*j]),cimag(a[i+na*j]) );
    }
    printf( "\n" );
}

printf( "\n\tConstant Vector (Real, Imaginary)\n\n");
for( i=0 ; i<n ; i++ )
{
    double tmp_re;
    fscanf( fp, "%lf", &tmp_re );
    b[i] = tmp_re;
}
for( i=0 ; i<n ; i++ )
{
    double tmp_im;
    fscanf( fp, "%lf", &tmp_im );
    b[i] = b[i] + tmp_im * _Complex_I;
}
for( i=0 ; i<n ; i++ )
{
    printf( "\t(%8.3g , %8.3g)\n", creal(b[i]),cimag(b[i]) );
}
fclose( fp );

ierr = ASL_hbhesl(a, na, n, b, nt);

printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\tSolution (Real, Imaginary)\n\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t x[%6d] = (%8.3g , %8.3g)\n", i,creal(b[i]),cimag(b[i]));
}

```

```

    free( a );
    free( b );
    return 0;
}

```

(d) Output results

```

*** ASL_hbhesl ***
** Input **
n =      4
Coefficient Matrix (Real, Imaginary)
(      9 ,      0) (      7 ,      3) (      2 ,      5) (      1 ,      1)
(      0 ,      0) (     10 ,      0) (      3 ,      2) (      2 ,      4)
(      0 ,      0) (      0 ,      0) (      8 ,      0) (      5 ,      1)
(      0 ,      0) (      0 ,      0) (      0 ,      0) (      6 ,      0)

Constant Vector (Real, Imaginary)
(     10 ,      6)
(     11 ,      2)
(      4 ,      6)
(      4 ,      6)

** Output **
ierr =      0
Solution (Real, Imaginary)
x[  0] = (      1 , -9.87e-17)
x[  1] = (      1 , -6.25e-17)
x[  2] = (5.11e-17 ,      1)
x[  3] = (      0 ,      1)

```

### 3.10.2 ASL\_hbheud, ASL\_gbheud

#### LDL\* Decomposition of a Hermitian Matrix (No Pivoting)

(1) **Function**

ASL\_hbheud or ASL\_gbheud uses the modified Cholesky method to perform an LDL\* decomposition of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type).

(2) **Usage**

Double precision:

ierr = ASL\_hbheud (a, lna, n, nt);

Single precision:

ierr = ASL\_gbheud (a, lna, n, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\left\{ \begin{array}{l} Z^* \\ C^* \end{array} \right\}$	$lna \times n$	Input	Hermitian matrix $A$ (two-dimensional array type, upper triangular type)
				Output	Upper triangular matrix $L^*$ when $A$ is decomposed into $A = LDL^*$ (See Note (a))
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	nt	I	1	Input	Number of tasks to be generated
5	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	Contents of array a are not changed.
2100	There existed the diagonal element which was close to zero in the LDL* decomposition of the coefficient matrix $A$ . The result may not be obtained with a good accuracy.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000+i	A diagonal element became equal to 0.0 in the $i$ -th processing step. $A$ is nearly singular.	

(6) Notes

- (a) The upper triangular matrix  $L^*$  is stored in the upper triangular portion of array a. Since the diagonal matrix  $D$  and the lower triangular matrix  $L$  are calculated from  $L^*$ , they are not stored in array a. This function uses only the upper triangular portion of array a. (See Fig. 3–10 in Section 3.10.1)

## Chapter 4

---

# SIMULTANEOUS LINEAR EQUATIONS (ITERATIVE METHOD)

### 4.1 INTRODUCTION

This chapter describes functions that use iterative methods to solve simultaneous linear equations having a large dimensional sparse matrix as coefficient matrix.

**Function described in this chapter divides up and allocates internal processing among threads and executes allocated processing in parallel.**

This library having the functions listed below use either the nonstationary iterative methods as the basic iterative algorithms. In the nonstationary iterative methods are included the methods such as the CG iterative method or the GMRES method.

For the preconditioner, the Scaling method are supported.

Among these functions, one matrix storage format for and two matrix storage formats for random sparse matrices (ELLPACK format) are available.

These functions use the same user interface regardless of the iterative basic method. The users can only have to change the functions name which is called in the user's program to try different iterative methods under the the same preconditioning and the same storage format.

Table 4–1 Function for each basic iterative method

BASIC ITERATIVE METHOD	FUNCTION NAME
CG (For Symmetric Matrices only) Conjugate Gradient method	{ ASL_qxe010 } { ASL_pxe010 }
CGS Conjugate Gradient Squared method	{ ASL_qxe020 } { ASL_pxe020 }
BiCGSTAB Bi-Conjugate Gradient Stabilized method	{ ASL_qxe030 } { ASL_pxe030 }
GMRES General Minimal Residual method	{ ASL_qxe040 } { ASL_pxe040 }



### 4.1.1 Notes

(1) Proper use of iterative methods

The CG method functions is primarily used to solve positive definite symmetric large-scale sparse simultaneous linear equations that occur in finite difference approximations or finite element approximations of diffusion equations. The CGS, BiCGSTAB method functions are primarily used to solve asymmetric large-scale sparse simultaneous linear equations that occur in finite difference approximations or finite element approximations of advection diffusion equations. Since the memory area required when using an iterative method is smaller than that required when using a direct method, iterative methods are used for solving large-scale problems. However, an iterative method may not converge to the true solution. Therefore, before using these iterative method functions, you should test whether you can reduce the maximum number of iterations or the problem size.

In particular, it is generally difficult to solve large-scale asymmetric simultaneous linear equations, and there is no guarantee that the true solution will always be obtained by applying the iterative method algorithms used in this library.

(2) Matrix storage format

The nonzero elements of the coefficient matrix of simultaneous linear equations generated by using the finite difference method to discretize a two-dimensional rectangular area or three-dimensional rectangular parallelepiped area line along straight lines in the direction of the diagonal. This type of matrix is called a regular sparse matrix.

However, the coefficient matrix of simultaneous linear equations generated by using the finite element method to discretize a randomly shaped area generally has no regularity other than the fact that the diagonal elements are nonzero elements. This type of matrix is called an random (irregular) sparse matrix. The functions described in this chapter support a storage format for random (irregular) sparse matrices.

Each function has a common user interface, with a single exception that the GMRES method functions have an additional input argument (gmitr). Therefore, the users can only have to replace the function name in their programs if they want to try different iterative methods with the input values unchanged. (As for trying GMRES method, an input argument (gmitr) should also be added to the user interface.) However, attention must be paid to the fact that the CG method functions cannot be applied to problems for which the coefficient matrix is asymmetric, since the CG method is for solving positive definite symmetric simultaneous linear equations.

(3) Preconditioning methods

Preconditioning is a technique for accelerating convergence of a basic iterative method. The simultaneous linear equations  $A\mathbf{u} = \mathbf{b}$  are converted to equivalent equations that are easier to solve by using an easily invertible matrix  $M$ , which is an approximation matrix for matrix  $A$ . The procedure for the basic iterative method is then applied to these equations. In the functions described here, the scaling preconditioning is selected.

Scaling preconditioning is the most effective of these preconditioning methods for ordinary problems because the large vectors involved in the iteration calculations of scaling preconditioning enable the multiple parallel pipelines of the Vector Engine to operate efficiently, and also the required memory area is small.

(4) Conditions for terminating iterations

The iteration calculations are interrupted when any of the following conditions occurs:

- (a) The residual norm becomes less than or equal to the truncation residual norm.
- (b) The iteration count reaches the maximum number of iterations.

(c) An error in the algorithm is detected that makes further calculations impossible.

The truncation residual norm is briefly described below.

If  $\mathbf{u}^*$  is the approximate solution obtained by an iteration calculation for the simultaneous linear equations  $A\mathbf{u} = \mathbf{b}$ , then the residual vector at that time is defined by:

$$\mathbf{r} = \mathbf{b} - A\mathbf{u}^*.$$

Here, the term “residual norm” is used to mean the (ordinary) relative norm

$$\frac{\|\mathbf{r}\|}{\|\mathbf{b}\|}.$$

In addition, the user can determine the meaning of the norm function  $\|\cdot\|$  from among  $L^2$  norm, which are defined as follows:

$$\|\mathbf{r}\|_2 = \left( \sum_{i=1}^n r_i^2 \right)^{1/2}.$$

(5) Required memory areas

The functions (basic iterative functions plus matrix operation functions) described in this chapter require an area for storing the coefficient matrix of the equations, areas for storing the solution vector or right-hand side vector, and work area.

The sizes of these areas differ according to the basic iterative method, matrix storage format, and preconditioning method that are selected. For information about the sizes of these areas, see the notes in the explanations of the basic iterative functions and matrix operation functions.

### 4.1.2 Algorithms Used

#### 4.1.2.1 Nonstationary iterative method (for Symmetric Matrix only)

(1) **the CG method**

Consider the following simultaneous linear equations of order  $n$ :

$$A\mathbf{u} = \mathbf{b} \quad (1)$$

The CG method is an iterative method for solving the simultaneous linear equations shown in (1) for which the coefficient matrix  $A$  is a positive definite symmetric matrix. This method is represented as follows.

$$\begin{array}{l} \mathbf{u} = \mathbf{u}_1 \quad (\mathbf{u}_1 \text{ is the initial solution}), \\ \mathbf{r}_1 = \mathbf{b} - A\mathbf{u}_1, \\ \mathbf{p}_1 = \mathbf{r}_1. \\ \text{for } i = 1, 2, \dots, n \\ \left\{ \begin{array}{l} a_i = \frac{(\mathbf{r}_i, \mathbf{r}_i)}{(\mathbf{p}_i, A\mathbf{p}_i)}, \\ \mathbf{u}_{i+1} = \mathbf{u}_i + a_i\mathbf{p}_i, \\ \mathbf{r}_{i+1} = \mathbf{r}_i - a_iA\mathbf{p}_i, \\ b_i = \frac{(\mathbf{r}_{i+1}, \mathbf{r}_{i+1})}{(\mathbf{r}_i, \mathbf{r}_i)}, \\ \mathbf{p}_{i+1} = \mathbf{r}_{i+1} + b_i\mathbf{p}_i. \end{array} \right. \end{array}$$

A great advantage of the CG method is that, theoretically, a strict solution is obtained in  $n$  iterative calculations (where  $n$  is the order).

#### 4.1.2.2 Nonstationary iterative method (for Asymmetric Matrix)

(1) **the CGS and BiCGSTAB methods**

To explain the CGS and BiCGSTAB methods, we will first explain the BiCG method. The BiCG method considers the simultaneous linear equations of order  $2n$  formed by combining the equations shown in (1) with the dual equations as follows:

$$\begin{aligned} \tilde{A}\tilde{\mathbf{u}} &= \tilde{\mathbf{b}}, \\ \tilde{A} &= \begin{bmatrix} A & O \\ O & A^T \end{bmatrix}, \tilde{\mathbf{u}} = \begin{bmatrix} \mathbf{u} \\ \mathbf{u}^* \end{bmatrix}, \tilde{\mathbf{b}} = \begin{bmatrix} \mathbf{b} \\ \mathbf{b}^* \end{bmatrix} \end{aligned}$$

and uses the conjugate gradient method to obtain the stationary value of:

$$F(\tilde{\mathbf{u}}) = \langle (\tilde{\mathbf{u}} - \hat{\mathbf{u}}), A(\tilde{\mathbf{u}} - \hat{\mathbf{u}}) \rangle_H \quad (\hat{\mathbf{u}}; \text{Truesolution})$$

where,

$$\langle \mathbf{u}, \mathbf{v} \rangle_H \equiv (\mathbf{u}, H\mathbf{v}) = (H\mathbf{u}, \mathbf{v}), \quad H = \begin{bmatrix} O & I \\ I & O \end{bmatrix}.$$

Using the initial residual vector  $\mathbf{r}_0$  the residual vector after  $k$  iterations  $\mathbf{r}_k$  and direction vector after  $k$  iterations  $\mathbf{p}_k$  are represented by:

$$\begin{aligned} \mathbf{r}_k &= R_k(A)\mathbf{r}_0, \\ \mathbf{p}_k &= P_k(A)\mathbf{r}_0. \end{aligned}$$

Here,  $R_k(A)$ ,  $P_k(A)$  are polynomials generated from  $A$ . At this time, the CGS method uses this  $R_k(A)$  and  $P_k(A)$  to calculate new vectors  $\mathbf{r}'_k$  and  $\mathbf{p}'_k$  as follows.

$$\begin{aligned}\mathbf{r}'_k &= R_k^2(A)\mathbf{r}_0, \\ \mathbf{p}'_k &= P_k^2(A)\mathbf{r}_0\end{aligned}$$

Using this calculation method, as

$$\|\mathbf{r}_k\| = \|R_k(A)\mathbf{r}_0\|$$

gets smaller,

$$\|\mathbf{r}'_k\| = \|R_k^2(A)\mathbf{r}_0\|$$

is expected to get even smaller, and convergence is expected to be faster than with the BCG method.

The method in which the iterations proceed so that the residual becomes  $\mathbf{r}'_k$  is the CGS method.

However, the CGS method may exhibit extremely irregular convergence characteristics. In particular, the residual may be extremely large at the beginning of the iterations such as when the iteration initial value is close to the solution, and accurate iteration calculations may not be able to be performed due to the effect of rounding error. The BiCGSTAB method lets  $Q_k(A)$  represent the following:

$$Q_k(A) = (1 - \omega_1 A)(1 - \omega_2 A) \cdots (1 - \omega_k A)$$

and obtains the approximate solution  $\mathbf{u}_k$  so that the residual

$$\mathbf{r}_k = \mathbf{b} - A\mathbf{u}_k$$

in each iteration is given by:

$$\mathbf{r}_k = Q_k(A)R_k(A)\mathbf{r}_0$$

Here, the parameters  $\omega_k$  that minimize  $(\mathbf{r}_k, \mathbf{r}_k)$  are selected so that the residuals are reduced stably.

(2) **the GMRES method (GMRES(m))**

The GMRES method is an iterative method that for each  $j = 1, 2, \dots$  searches for the vector  $\mathbf{z}^{(j)}$  that minimizes the  $L^2$  norm of

$$\mathbf{b} - A(\mathbf{u}_0^{(j)} + \mathbf{z}^{(j)}) = \mathbf{r}^{(j)} - A\mathbf{z}^{(j)}$$

from the space spanned by

$$\mathbf{r}_0, \quad A\mathbf{r}_0, \quad A^2\mathbf{r}_0, \quad \dots, \quad A^i\mathbf{r}_0 \quad (i = 1, 2, \dots, m)$$

when the initial approximate solution vector of the  $j$ -th step is  $\mathbf{u}_0^{(j)}$  and the initial residual vector of the  $j$ -th step is

$$\mathbf{r}^{(j)} = \mathbf{b} - A\mathbf{u}_0^{(j)}$$

```

 $\mathbf{u}_0^{(1)} = \mathbf{u}_0.$ 
for  $j = 1, 2, \dots$ 
┌
   $\mathbf{r}^{(j)} = \mathbf{b} - A\mathbf{u}_0^{(j)},$ 
   $\mathbf{v}_1 = \mathbf{r}^{(j)} / \|\mathbf{r}^{(j)}\|_2.$ 
  for  $i = 1, 2, \dots, m$ 
  ┌
     $\mathbf{w} = A\mathbf{v}_i.$ 
    for  $k = 1, 2, \dots, i$ 
    ┌
       $\mathbf{w} = \mathbf{w} - (\mathbf{w}, \mathbf{v}_k)\mathbf{v}_k,$ 
    └
     $\mathbf{v}_{i+1} = \mathbf{w} / \|\mathbf{w}\|_2,$ 
    Determine the parameters  $\tilde{\mathbf{u}}^{(j)} = \mathbf{u}_0^{(j)} + y_1\mathbf{v}_1 + y_2\mathbf{v}_2 + \dots + y_i\mathbf{v}_i$ 
    so that  $\|\mathbf{b} - A\tilde{\mathbf{u}}^{(j)}\|_2$  takes the least value.
  └
 $\mathbf{u}_0^{(j+1)} = \tilde{\mathbf{u}}^{(j)}.$ 

```

As shown in this algorithm, the GMRES method consists of an external iterative calculation (loop related to  $j$ ) and an internal iterative calculation (loop related to  $i$ ). Since the GMRES method is an iterative method that depends on the parameter  $m$  representing the upper bound on the number of inner iterations, it is referred to, strictly speaking, as the GMRES( $m$ ) method.

In this library, the number of iterations in the GMRES( $m$ ) algorithm is defined by (number of internal iterations  $i$ ) + (number of external iterations  $j$ )  $\times m$ .

#### 4.1.2.3 Preconditioned Iterative Method

Since error is introduced in actual iterative method calculations, convergence characteristics become worse. In particular, if the difference mesh is fine in the difference scheme for partial differential equations, if the coefficient (dispersion or thermal conductivity) values severely fluctuate spatially or if the difference mesh widths or the coefficients described above are anisotropic, then the eigenvalues of matrix  $A$  are dispersed, and the iterative method has difficulty in converging under these circumstances. Conversely, convergence will be fast if the eigenvalues are clustered together.

To overcome the difficulties described above, preconditioned iterative methods were developed. A preconditioned iterative method is a technique for improving convergence by preconditioning the original matrix to convert it to a well-conditioned matrix before performing the iterative method calculations. This is represented as follows. Assume that the matrix  $M = M_1M_2$ , which approximates  $A = (a_{ij})$  in some way, is given.

$$A \sim M = M_1M_2$$

At this time, if the following equation, which is equivalent to equation (1), is created:

$$\begin{aligned}
 M_1^{-1}AM_2^{-1}\mathbf{u}' &= \mathbf{b}' & (2) \\
 (\mathbf{u}' &= M_2\mathbf{u} \ ) \\
 (\mathbf{b}' &= M_1^{-1}\mathbf{b} \ )
 \end{aligned}$$

the coefficient matrix of this equation is easier to solve since it is closer to the unit matrix than the original matrix was. Equation (2) is called the preconditioning equation, and matrix  $M$  is called the preconditioning matrix. The preconditioned iterative method for equation (1) is defined/derived as the (basic) iterative method for equation (1). The preconditioned iterative method must execute a matrix inversion on  $M_1$  and  $M_2$  or  $M$  (left operation on a column vector according to the inverse matrix). This is what is called the preconditioning operation (or simply preconditioning). Usually, a matrix for which the preconditioning operation can easily be performed is used as preconditioning matrix  $M = M_1M_2$ . The more closely the preconditioning matrix  $M$  approximates the original

matrix  $A$ , the more the convergence of the preconditioned iterative method will improve. However, there generally is a tendency for the amount of time required for a single preconditioning operation to increase in proportion to the quality of convergence of the preconditioning matrix.

See 4.1.2.4 for the detailed explanation of the preconditioning algorithms that ASL supports.

The algorithms for the preconditioned CG method (PCG method), preconditioned CGS method (PCGS method), preconditioned BiCGSTAB (PBiCGSTAB method), and preconditioned GMRES(m) method (PGMRES(m) method) are shown below.

(1) **the PCG method**

$$\mathbf{u} = \mathbf{u}_1 \quad (\mathbf{u}_1 \text{ is the initial solution}),$$

$$\mathbf{r}_1 = \mathbf{b} - A\mathbf{u}_1,$$

$$\mathbf{p}_1 = M^{-1}\mathbf{r}_1.$$

for  $i = 1, 2, \dots, n$

$$a_i = \frac{(\mathbf{r}_i, M^{-1}\mathbf{r}_i)}{(\mathbf{p}_i, A\mathbf{p}_i)},$$

$$\mathbf{u}_{i+1} = \mathbf{u}_i + a_i\mathbf{p}_i,$$

$$\mathbf{r}_{i+1} = \mathbf{r}_i - a_i A\mathbf{p}_i,$$

$$b_i = \frac{(\mathbf{r}_{i+1}, M^{-1}\mathbf{r}_{i+1})}{(\mathbf{r}_i, M^{-1}\mathbf{r}_i)},$$

$$\mathbf{p}_i = M^{-1}\mathbf{r}_{i+1} + b_i\mathbf{p}_i.$$

(2) **the PCGS method** (Bibliography(1))

$$\mathbf{r}_0 = \mathbf{b} - A\mathbf{u}_0,$$

$$\mathbf{p}_0 = M^{-1}\mathbf{r}_0,$$

$$\mathbf{e}_0 = \mathbf{r}_0.$$

for  $k = 0, 1, 2, \dots$

$$\alpha_{k+1} = \frac{(\mathbf{r}_0, \mathbf{r}_k)}{(\mathbf{r}_0, A\mathbf{p}_k)},$$

$$\mathbf{h}_{k+1} = \mathbf{e}_k - \alpha_{k+1}A\mathbf{p}_k,$$

$$\mathbf{u}_{k+1} = \mathbf{u}_k + \alpha_{k+1}M^{-1}(\mathbf{e}_k + \mathbf{h}_{k+1}),$$

$$\mathbf{r}_{k+1} = \mathbf{r}_k - \alpha_{k+1}AM^{-1}(\mathbf{e}_k + \mathbf{h}_{k+1}),$$

$$\beta_{k+1} = \frac{(\mathbf{r}_0, \mathbf{r}_{k+1})}{(\mathbf{r}_0, \mathbf{r}_k)},$$

$$\mathbf{e}_{k+1} = \mathbf{r}_{k+1} + \beta_{k+1}\mathbf{h}_{k+1},$$

$$\mathbf{p}_{k+1} = M^{-1}(\mathbf{e}_{k+1} + \beta_{k+1}\mathbf{h}_{k+1}) + \beta_{k+1}^2\mathbf{p}_k.$$

(3) the **PBiCGSTAB** method (Bibliography(2))

```

 $\mathbf{r}_0 = \mathbf{b} - A\mathbf{u}_0,$ 
for  $k = 0, 1, 2, \dots$ 
|
|  $\rho_{k+1} = (\mathbf{r}_0, \mathbf{r}_k),$ 
| if  $k = 0$  then
|    $\mathbf{p}_1 = \mathbf{r}_0.$ 
| else
|    $\beta = \left(\frac{\rho_{k+1}}{\rho_k}\right)\left(\frac{\alpha}{\omega_k}\right),$ 
|    $\mathbf{p}_{k+1} = \mathbf{r}_k + \beta(\mathbf{p}_k - \omega_k \mathbf{v}_k).$ 
| end if
|  $\mathbf{y} = M^{-1}\mathbf{p}_{k+1},$ 
|  $\mathbf{v}_{k+1} = A\mathbf{y},$ 
|  $\alpha = \frac{\rho_{k+1}}{(\mathbf{r}_0, \mathbf{v}_{k+1})},$ 
|  $\mathbf{s} = \mathbf{r}_k - \alpha\mathbf{v}_{k+1},$ 
|  $\mathbf{z} = M^{-1}\mathbf{s},$ 
|  $\mathbf{t} = A\mathbf{z},$ 
|  $\omega_{k+1} = \frac{(\mathbf{t}, \mathbf{s})}{(\mathbf{t}, \mathbf{t})},$ 
|  $\mathbf{u}_{k+1} = \mathbf{u}_k + \alpha\mathbf{y} + \omega_{k+1}\mathbf{z},$ 
|  $\mathbf{r}_{k+1} = \mathbf{s} - \omega_{k+1}\mathbf{t}.$ 

```

(4) the **GMRES(m)** method (Bibliography(3))

```

 $\mathbf{u}_0^{(1)} = \mathbf{u}_0.$ 
for  $j = 1, 2, \dots$ 
|
|  $\mathbf{r}^{(j)} = M^{-1}(\mathbf{b} - A\mathbf{u}_0^{(j)}),$ 
|  $\mathbf{v}_1 = \frac{\mathbf{r}^{(j)}}{\|\mathbf{r}^{(j)}\|_2}.$ 
| for  $i = 1, 2, \dots, m$ 
| |
| |  $\mathbf{w} = M^{-1}A\mathbf{v}_i.$ 
| | for  $k = 1, 2, \dots, i$ 
| | |  $\mathbf{w} = \mathbf{w} - (\mathbf{w}, \mathbf{v}_k)\mathbf{v}_k.$ 
| |  $\mathbf{v}_{i+1} = \frac{\mathbf{w}}{\|\mathbf{w}\|_2},$ 
| | Determine the parameters  $\tilde{\mathbf{u}}^{(j)} = \mathbf{u}_0^{(j)} + y_1\mathbf{v}_1 + y_2\mathbf{v}_2 + \dots + y_i\mathbf{v}_i$ 
| | so that  $\|\mathbf{b} - A\tilde{\mathbf{u}}^{(j)}\|_2$  takes the least value.
|  $\mathbf{u}_0^{(j+1)} = \tilde{\mathbf{u}}^{(j)}$ 

```

**Note:** the number of iterations in the GMRES(m) algorithm is defined by  
(number of internal iterations  $i$ ) +(number of external iterations  $j$ ) $\times m$  .

#### 4.1.2.4 Preconditioning Methods

In the following are described the details of those preconditioning algorithms that are implemented in functions in this chapter, synopses of which are explained 4.1.2.3.

(1) **Scaling preconditioning**

Scaling preconditioning is a technique that uses the following matrix consisting of the diagonal components of matrix  $A$  as the preconditioning matrix  $M$ .

$$D = \text{diag}(a_{11}, a_{22}, \dots, a_{nn}),$$

In particular, the scaling preconditioned CG method is known as the SCG method. On an Vector Engine, the scaling method often is faster than other preconditioning techniques.

#### 4.1.2.5 Advanced Techniques for Improving Performance

##### (1) Iterative Improvement Method

An iterative improvement method is a technique for constructing a higher precision solution based on an approximate solution obtained by some solution method. If a single-precision function was used, then a solution on the order of a double-precision solution will be obtained, and if a double-precision function was used, then a solution on the order of a quadruple-precision solution will be obtained. (However, the calculation time will be approximately two to four times longer.) Another benefit of an iterative improvement method is that the error of the obtained solution can be estimated. If an iterative improvement method is not used, then the error is difficult to estimate even if the residual is known. A problem with the PCG method is that the error cannot be improved beyond a certain degree no matter how close the residual gets to zero. Therefore, if the precision of the solution is particularly important or if a single-precision function has been used, then an iterative improvement of the solution should be performed.

The iterative improvement method can be intuitively derived from the following kind of transformation, where  $\mathbf{u}$  is the true solution of the equations  $A\mathbf{u} = \mathbf{b}$  and  $\mathbf{u}'$  is the approximation solution:

$$\begin{aligned} \mathbf{u} &= \mathbf{u}' + \mathbf{u} - \mathbf{u}' \\ &= \mathbf{u}' + A^{-1}\mathbf{b} - \mathbf{u}' \\ &= \mathbf{u}' + A^{-1}(\mathbf{b} - A\mathbf{u}') \end{aligned}$$

From this transformation the following algorithm is obtained:

- (a) Use a function from this chapter to solve the given equations  $A\mathbf{u} = \mathbf{b}$  and let  $\mathbf{u}'$  be the solution that was obtained.
- (b) Obtain the residual of the solution with high precision. That is, if a single-precision function was used, calculate  $\mathbf{r}' = \mathbf{b} - A\mathbf{u}'$  with double precision, and if a double-precision function was used, calculate  $\mathbf{r}'$  with quadruple precision.
- (c) Return  $\mathbf{r}'$  to the original precision and assign it to  $\mathbf{r}$ . That is, for single-precision function, assign  $\mathbf{r}'$  to a single-precision  $\mathbf{r}$ , and for a double-precision function, assign  $\mathbf{r}'$  to a double-precision  $\mathbf{r}$ .
- (d) Use a function from this chapter to solve the equations  $A\mathbf{v} = \mathbf{r}$  and let  $\mathbf{v}'$  be the solution that was obtained.
- (e) Calculate  $\mathbf{u}' + \mathbf{v}'$  and let it be called  $\mathbf{u}'$ .
- (f)  $\mathbf{v}'$  is the estimated error. If  $\mathbf{v}'$  is sufficiently small, then stop. If not, then return to (b).

The most important part of this algorithm is step (b). The error is improved by calculating the residual with high precision. The iterative improvement method generally can obtain a sufficient degree of precision in several iterations.

##### (2) Node numbering in the finite element method

Although it is known that odd-even ordering or multicolor ordering, for example, should be performed to obtain long vectors, this kind of tricky numbering may significantly degrade convergence in a preconditioned



iterative method. Therefore, care is required(Bibliography(4)). The most stable numbering is numbering that reduces the matrix band width. To automatically generate this kind of numbering, you should use a technique such as the Cuthill-McKee method(Bibliography(5)).

### 4.1.3 Reference Bibliography

- (1) Sonneveld, P. , “CGS, a Fast Lanczos–type Solver for Nonsymmetric Linear Systems”, Delft University of Technology, Report No. 84–16, Delft The Netherland (1984).
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## 4.2 SPARSE MATRIX—NONSTATIONARY ITERATIVE METHODS (BASIC ITERATION METHOD FUNCTIONS)

### 4.2.1 ASL\_qxe010, ASL\_pxe010

Positive Definite Symmetric Sparse Matrix (ELLPACK Format) (CG method)

(1) **Function**

ASL\_qxe010 or ASL\_pxe010 uses the scaling preconditioned CG method to solve the simultaneous linear equations  $A\mathbf{u} = \mathbf{b}$  having a sparse matrix  $A$  as the coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_qxe010 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

Single precision:

```
ierr = ASL_pxe010 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I: {int as for 32bit Integer}  
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$l_n \times m$	Input	Array of nonzero element values of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
2	lna	I	1	Input	Adjustable dimension of array a and ja
3	n	I	1	Input	Order of matrix $A$
4	m	I	1	Input	m denotes the column number of both array a and ja (See Note (d))
5	ja	$I^*$	$l_n \times m$	Input	Array storing the nonzero structure data of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
6	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Right-hand side vector $\mathbf{b}$ of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
7	u	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Solution vector $\mathbf{u}$ of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
8	itrmax	I	1	Input	Maximum number of iterations (Default value: n)
9	itr	$I^*$	1	Output	Actual number of iterations
10	epsmax	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Truncation residual norm Default value : $\left\{ \begin{array}{l} 10^{-12} \text{ (doubleprecision)} \\ 10^{-6} \text{ (singleprecision)} \end{array} \right\}$
11	eps	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	1	Output	Final residual norm
12	isw	I	1	Input	Processing switch (Default value: 0) (See Note (e)) isw= 0 : Check whether the values ja are all distinct. isw= 1 : Does not check whether the values ja are all distinct.
13	nt	I	1	Input	Number of tasks to be generated (Default value: 1)
14	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $1 \leq n \leq \text{lna}$
- (b)  $1 \leq m \leq n$
- (c) •  $\text{ja}[i-1] = i, \text{ja}[(i-1) + \text{lna} \times (j-1)] \neq i$  ( $i = 1, \dots, n; j = 2, \dots, j_i$ ),  
 $1 \leq \text{ja}[(i-1) + \text{lna} \times (j-1)] \leq n$  ( $j = 2, \dots, j_i$ ),  
 where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- $\text{ja}[(i-1) + \text{lna} \times j_i] = 0$ , if  $j_i < m$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (d)  $\text{ja}[(i-1) + \text{lna} \times (j-1)]$  ( $j = 1, \dots, j_i$ ) are all distinct for a fixed value  $i$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (e)  $\text{a}[i-1] \neq 0.0$  ( $i = 1, \dots, n$ )
- (f)  $\text{itrmax} \geq 1$
- (g)  $\text{epsmax} >$  underflow decision value
- (h)  $\text{isw} \in \{0, 1\}$
- (i)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	Restriction (f) was not satisfied.	Processing continues with $\text{itrmax} = n$ .
1200	Restriction (g) was not satisfied.	Processing continues with $\text{epsmax} = \begin{cases} 10^{-12} & \text{(doubleprecision)} \\ 10^{-6} & \text{(singleprecision)} \end{cases}$ .
1300	Restriction (h) was not satisfied.	Processing continues with $\text{isw}=0$ , on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
1400	Restriction (i) was not satisfied.	Processing continues with $\text{nt} = 1$ .
2000	The absolute norm of the right-hand side vector $\mathbf{b}$ is less than the underflow decision value.	$\text{u}[i-1] \leftarrow 0.0$ ( $i = 1, \dots, n$ ) is set for the solution.
2100	Restriction (e) was not satisfied.	Processing continues.
2200	$\text{isw}$ was equal to 0.	Processing continues on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	
3300	Restriction (d) was not satisfied.	
3500	The norm of the Maximum number of iterations has been reached.	The result obtained at that time is returned.

ierr value	Meaning	Processing
4000	a has a diagonal term whose absolute value is smaller than the underflow decision value.	Processing is aborted.
4100	The norm of the right-hand side vector $\mathbf{b}$ is greater than the overflow decision value.	
4210	The norm of the residual $\mathbf{r} = \mathbf{b} - \mathbf{A}\mathbf{u}$ is greater than the overflow decision value.	
4220	The relative norm of the residual $\mathbf{r}$ is greater than the overflow decision value.	
4310	$ (\mathbf{r}_i, M^{-1}\mathbf{r}_i) $ is less than the underflow decision value.	
4320	$ (\mathbf{r}_i, M^{-1}\mathbf{r}_i) $ is greater than the overflow decision value.	
4410	$ (\mathbf{p}_i, \mathbf{A}\mathbf{p}_i) $ is less than the underflow decision value.	
4420	$ (\mathbf{p}_i, \mathbf{A}\mathbf{p}_i) $ is greater than the overflow decision value.	
5000	Processing to reserve the work area required to execute the basic iterative algorithm failed (See Note (f)).	

(6) Notes

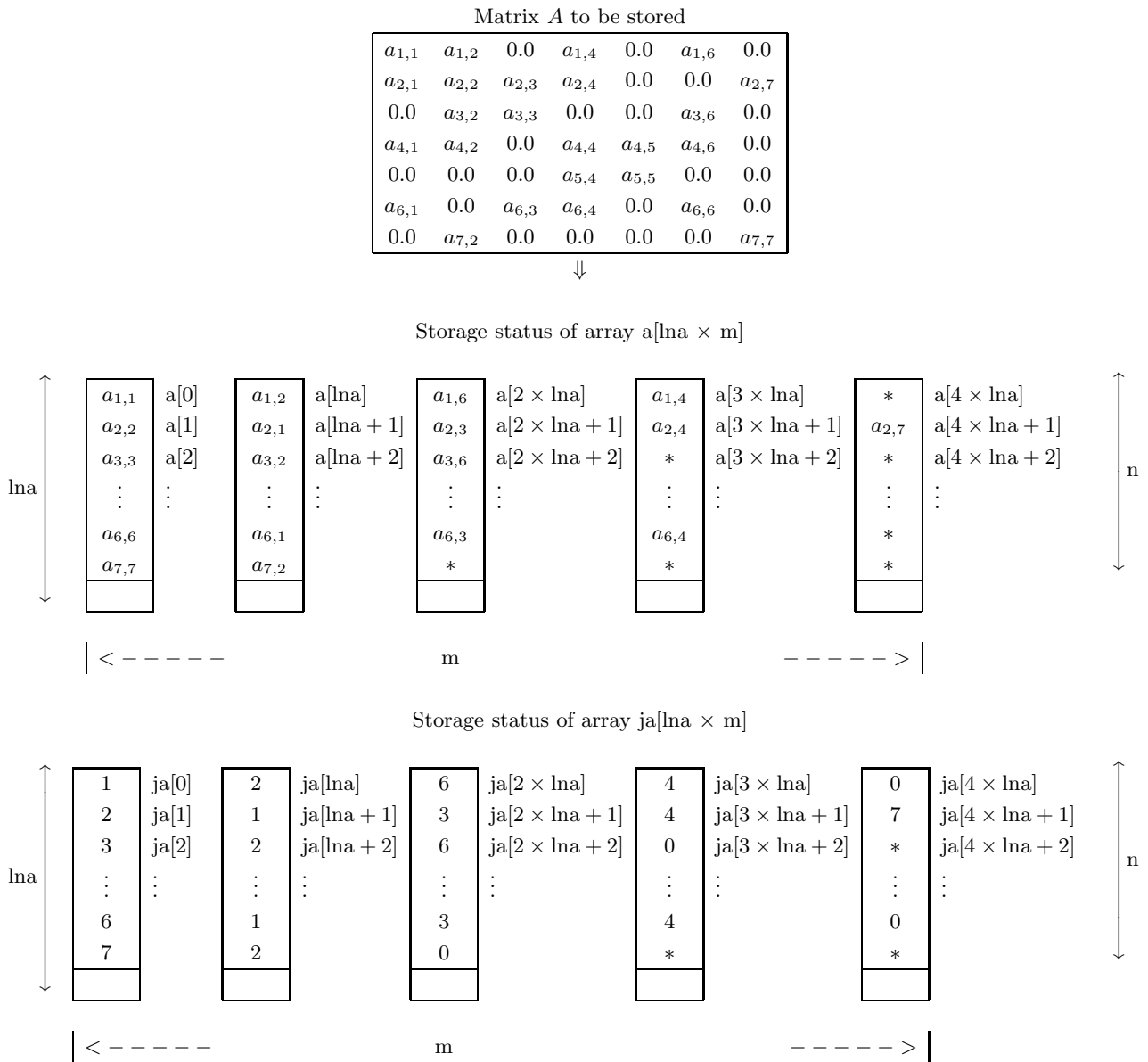
- (a) The table below shows maximum and minimum values of floating point data and so on, that are defined within ASL C interface.

Table 4-2 Constants used in ASL C interface

	Double-precision	Single-precision
Maximum value	$2^{1023}(2 - 2^{-52}) \simeq (1.80 \times 10^{308})$	$2^{127}(2 - 2^{-23}) \simeq (3.40 \times 10^{38})$
Positive minimum value	$2^{-1022} \simeq (2.23 \times 10^{-308})$	$2^{-126} \simeq (1.17 \times 10^{-38})$
Negative maximum value	$-2^{-1022} \simeq (-2.23 \times 10^{-308})$	$-2^{-126} \simeq (-1.17 \times 10^{-38})$
Minimum value	$-2^{1023}(2 - 2^{-52}) \simeq (-1.80 \times 10^{308})$	$-2^{127}(2 - 2^{-23}) \simeq (-3.40 \times 10^{38})$
Overflow decision value	Maximum value $\times 10^{-3}$	
Underflow decision value	Positive minimum value $\times 10^3$	

- (b) The storage method used for the arrays a and ja is as follows.

Figure 4-1 Storage format for Input Data



- (c) The input values of array a and array ja are partially changed in order to optimize performance.
- (d) m can take an arbitrary value as long as it is equal to or greater than the maximum of the numbers of nonzero elements in a row, but it is recommended from the viewpoint of calculation cost to make the padded area as little as possible by setting m as close as possible to the maximum of the numbers of nonzero elements in a row.
- (e) A better performance will be achieved by specifying isw = 0 (Performance will degrade remarkably when isw = 1 is specified). And so, isw = 1 should be specified when it is assured that ja satisfies Restriction (d). When isw = 1 is specified, the user should be very careful in giving the coefficient matrix data as input because the check for Restriction (d) is omitted. The validity of the result will not be guaranteed if isw = 0 was specified but ja doesn't satisfy Restriction (d).
- (f) The work area is automatically allocated within this function. If the work area could not be successfully allocated, processing will be aborted and ierr returns the value 5000. In this case, the problem cannot be solved by using this function unless the problem size is reduced or the machine environment is enhanced in memory size.

(7) **Example**

- (a) Problem

Solve  $Au = b$  for the following matrix  $A$  and vector  $b$ :

$$A = \begin{bmatrix} 3 & 1 & 0 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 5 & 2 & 0 & -2 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 7 & 3 & 0 & -3 & 0 & 0 & 0 & 0 \\ -1 & 0 & 3 & 9 & 4 & 0 & -4 & 0 & 0 & 0 \\ 0 & -2 & 0 & 4 & 11 & 5 & 0 & -5 & 0 & 0 \\ 0 & 0 & -3 & 0 & 5 & 13 & 6 & 0 & -6 & 0 \\ 0 & 0 & 0 & -4 & 0 & 6 & 15 & 7 & 0 & -7 \\ 0 & 0 & 0 & 0 & -5 & 0 & 7 & 17 & 8 & 0 \\ 0 & 0 & 0 & 0 & 0 & -6 & 0 & 8 & 19 & 9 \\ 0 & 0 & 0 & 0 & 0 & 0 & -7 & 0 & 9 & 21 \end{bmatrix}, b = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

- (b) Input data

Input Arrays a, ja and b,

lna = 11, n = 10, m = 11, itrmax = 100, epsmax =  $10^{-12}$ , isw = 0, nt = 2

- (c) Main program

```

/*      C interface example for ASL-qxe010 */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    double *a;
    int lna;
    int n;
    int m;
    int *ja;
    double *b;
    double *u;
    int itrmax;
    int itr;
    double epsmax;
    double eps;
    int isw;
    int nt;
    int ierr;
    int i,j;

```



```

printf( "    *** ASL_qxe010 ***\n" );
lna = 11;
n   = 10;
m   = 5;

a = ( double * )malloc((size_t)( sizeof(double) * (lna*m) ));
if ( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}
ja = ( int * )malloc((size_t)( sizeof(int) * (lna*m) ));
if ( ja == NULL )
{
    printf( "no enough memory for array ja\n" );
    return -1;
}
b = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}
u = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( u == NULL )
{
    printf( "no enough memory for array u\n" );
    return -1;
}

for( j=0 ; j < m ; j++ )
{
    for( i=0 ; i < n ; i++ )
    {
        a[i+lna*j] = 0.0;
        ja[i+lna*j] = 0;
    }
}

a[0 ] = 3.0;
a[ lna ] = 1.0;
a[ lna*2 ] = -1.0;

a[1 ] = 5.0;
a[1+lna ] = 1.0;
a[1+lna*2 ] = 2.0;
a[1+lna*3 ] = -2.0;

a[2 ] = 7.0;
a[2+lna ] = 2.0;
a[2+lna*2 ] = 3.0;
a[2+lna*3 ] = -3.0;

a[3 ] = 9.0;
a[3+lna ] = -1.0;
a[3+lna*2 ] = 3.0;
a[3+lna*3 ] = 4.0;
a[3+lna*4 ] = -4.0;

a[4 ] = 11.0;
a[4+lna ] = -2.0;
a[4+lna*2 ] = 4.0;
a[4+lna*3 ] = 5.0;
a[4+lna*4 ] = -5.0;

a[5 ] = 13.0;
a[5+lna ] = -3.0;
a[5+lna*2 ] = 5.0;
a[5+lna*3 ] = 6.0;
a[5+lna*4 ] = -6.0;

a[6 ] = 15.0;
a[6+lna ] = -4.0;
a[6+lna*2 ] = 6.0;
a[6+lna*3 ] = 7.0;
a[6+lna*4 ] = -7.0;

a[7 ] = 17.0;
a[7+lna ] = -5.0;
a[7+lna*2 ] = 7.0;
a[7+lna*3 ] = 8.0;

a[8 ] = 19.0;
a[8+lna ] = -6.0;
a[8+lna*2 ] = 8.0;
a[8+lna*3 ] = 9.0;

a[9 ] = 21.0;
a[9+lna ] = -7.0;
a[9+lna*2 ] = 9.0;

ja[0 ] = 1;

```

```

ja[ lna ] = 2;
ja[ lna*2 ] = 4;

ja[1 ] = 2;
ja[1+lna ] = 1;
ja[1+lna*2 ] = 3;
ja[1+lna*3 ] = 5;

ja[2 ] = 3;
ja[2+lna ] = 2;
ja[2+lna*2 ] = 4;
ja[2+lna*3 ] = 6;

ja[3 ] = 4;
ja[3+lna ] = 1;
ja[3+lna*2 ] = 3;
ja[3+lna*3 ] = 5;
ja[3+lna*4 ] = 7;

ja[4 ] = 5;
ja[4+lna ] = 2;
ja[4+lna*2 ] = 4;
ja[4+lna*3 ] = 6;
ja[4+lna*4 ] = 8;

ja[5 ] = 6;
ja[5+lna ] = 3;
ja[5+lna*2 ] = 5;
ja[5+lna*3 ] = 7;
ja[5+lna*4 ] = 9;

ja[6 ] = 7;
ja[6+lna ] = 4;
ja[6+lna*2 ] = 6;
ja[6+lna*3 ] = 8;
ja[6+lna*4 ] = 10;

ja[7 ] = 8;
ja[7+lna ] = 5;
ja[7+lna*2 ] = 7;
ja[7+lna*3 ] = 9;

ja[8 ] = 9;
ja[8+lna ] = 6;
ja[8+lna*2 ] = 8;
ja[8+lna*3 ] = 10;

ja[9 ] = 10;
ja[9+lna ] = 7;
ja[9+lna*2 ] = 9;

for( i=0 ; i < n ; i++ )
{
    b[i] = 1.0;
}

for( i=0 ; i < n ; i++ )
{
    u[i] = 0.0;
}

itrmax = 100;
epsmax = 1.0e-12;
isw = 1;
nt = 2;

printf( "\n *** Input ***\n\n" );
printf( "\titrmax = %6d\n", itrmax );
printf( "\tepsmax =%8.3g\n", epsmax );
printf( "\tnt = %6d\n", nt );

    ierr = ASL_qxe010
(a,lna,n,m,ja,b,u,itrmax,&itr,epsmax,&eps,isw,nt);

printf( "\n *** Output ***\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\titr = %6d\n", itr );
printf( "\teps =%8.3g\n", eps );

for( i=0 ; i < n ; i++ )
{
    printf( "\tu[%2d ] =%8.3g\n", i, u[i] );
}

free( a );

```

```
    free( ja );
    free( b );
    free( u );
}
return 0;
```

(d) Output results

```
*** ASL_qxe010 ***
***   Input   ***
itrmax =    100
epsmax = 1e-12
nt      =     2
***   Output  ***
ierr   =     0
itr    =     10
eps    =3.46e-16
u[ 0 ] =  0.335
u[ 1 ] =  0.22
u[ 2 ] = -0.0808
u[ 3 ] =  0.224
u[ 4 ] =  0.136
u[ 5 ] = -0.151
u[ 6 ] =  0.247
u[ 7 ] =  0.0408
u[ 8 ] = -0.0926
u[ 9 ] =  0.17
```

### 4.2.2 ASL\_qxe020, ASL\_pxe020

#### Asymmetric Sparse Matrix (ELLPACK Format) (CGS method)

(1) **Function**

ASL\_qxe020 or ASL\_pxe020 uses the scaling preconditioned CGS method to solve the simultaneous linear equations  $A\mathbf{u} = \mathbf{b}$  having an asymmetric sparse matrix  $A$  as the coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_qxe020 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

Single precision:

```
ierr = ASL_pxe020 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I: {int as for 32bit Integer}  
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$l_n \times m$	Input	Array of nonzero element values of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
2	$l_n$	I	1	Input	Adjustable dimension of array a and ja
3	n	I	1	Input	Order of matrix A
4	m	I	1	Input	m denotes the column number of both array a and ja (See Note (d))
5	ja	$I^*$	$l_n \times m$	Input	Array storing the nonzero structure data of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
6	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Right-hand side vector <b>b</b> of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
7	u	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Solution vector <b>u</b> of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
8	itrmax	I	1	Input	Maximum number of iterations (Default value: n)
9	itr	$I^*$	1	Output	Actual number of iterations
10	epsmax	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Truncation residual norm Default value: $\begin{Bmatrix} 10^{-12} & (\text{doubleprecision}) \\ 10^{-6} & (\text{singleprecision}) \end{Bmatrix}$
11	eps	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	1	Output	Final residual norm
12	isw	I	1	Input	Processing switch (Default value: 0) (See Note (e)) isw= 0 : Check whether the values ja are all distinct. isw= 1 : Does not check whether the values ja are all distinct.
13	nt	I	1	Input	Number of tasks to be generated (Default value: 1)
14	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $1 \leq n \leq \text{lna}$
- (b)  $1 \leq m \leq n$
- (c)
  - $\text{ja}[i-1] = i, \text{ja}[(i-1) + \text{lna} \times (j-1)] \neq i$  ( $i = 1, \dots, n; j = 2, \dots, j_i$ ),  
 $1 \leq \text{ja}[(i-1) + \text{lna} \times (j-1)] \leq n$  ( $j = 2, \dots, j_i$ ),  
 where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
  - $\text{ja}[(i-1) + \text{lna} \times j_i] = 0$ , if  $j_i < m$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (d)  $\text{ja}[(i-1) + \text{lna} \times (j-1)]$  ( $j = 1, \dots, j_i$ ) are all distinct for a fixed value  $i$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (e)  $\text{a}[i-1] \neq 0.0$  ( $i = 1, \dots, n$ )
- (f)  $\text{itrmax} \geq 1$
- (g)  $\text{epsmax} > \text{underflowdecisionvalue}$
- (h)  $\text{isw} \in \{0, 1\}$
- (i)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	Restriction (f) was not satisfied.	Processing continues with $\text{itrmax} = n$ .
1200	Restriction (g) was not satisfied.	Processing continues with $\text{epsmax} = \begin{cases} 10^{-12} & (\text{doubleprecision}) \\ 10^{-6} & (\text{singleprecision}) \end{cases}$ .
1300	Restriction (h) was not satisfied.	$\text{isw}$ has the input value 0. Processing continues on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
1400	Restriction (i) was not satisfied.	Processing continues with $\text{nt} = 1$ .
2000	The absolute norm of the right-hand side vector $\mathbf{b}$ is less than the underflow decision value.	$\text{u}[i-1] \leftarrow 0.0$ ( $i = 1, \dots, n$ ) is set for the solution.
2100	Restriction (e) was not satisfied.	Processing continues.
2200	$\text{isw}$ was equal to 0.	Processing continues on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	
3300	Restriction (d) was not satisfied.	

ierr value	Meaning	Processing
3500	The norm of the Maximum number of iterations has been reached.	The result obtained at that time is returned.
4000	a has a diagonal term whose absolute value is smaller than the underflow decision value.	Processing is aborted.
4100	The norm of the right-hand side vector $\mathbf{b}$ is greater than the overflow decision value.	
4210	The norm of the residual $\mathbf{r} = \mathbf{b} - A\mathbf{u}$ is greater than the overflow decision value.	
4220	The relative norm of the residual $\mathbf{r}$ is greater than the overflow decision value.	
4310	$ (\mathbf{r}_0, \mathbf{r}_i) $ is less than the underflow decision value.	
4320	$ (\mathbf{r}_0, \mathbf{r}_i) $ is greater than the overflow decision value.	
4410	$ (\mathbf{r}_0, A\mathbf{p}_i) $ is less than the underflow decision value.	
4420	$ (\mathbf{r}_0, A\mathbf{p}_i) $ is greater than the overflow decision value.	
5000	Processing to reserve the work area required to execute the basic iterative algorithm failed (See Note (f)).	

(6) Notes

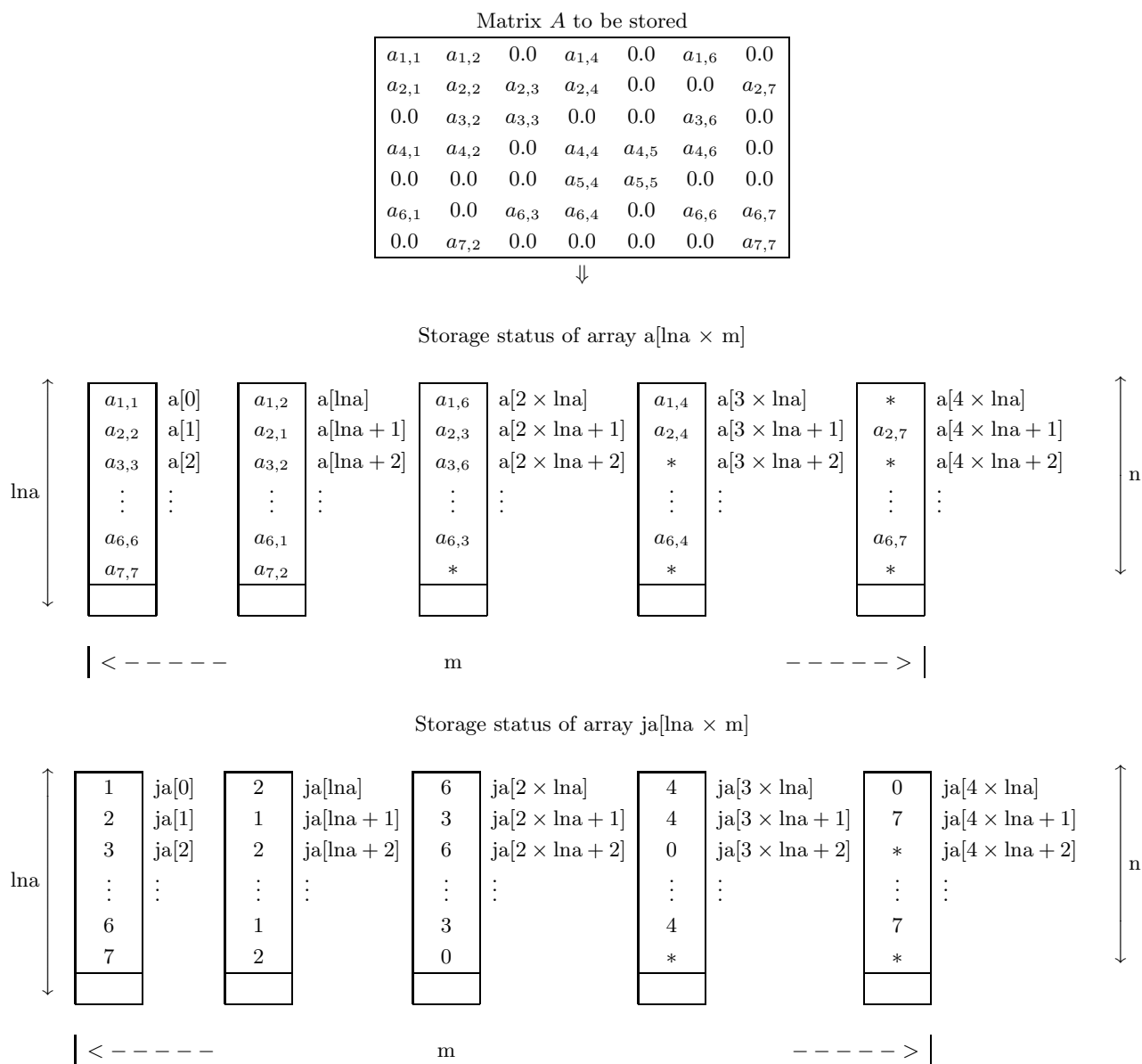
- (a) The table below shows maximum and minimum values of floating point data and so on, that are defined within ASL C interface.

Table 4-3 Constants used in ASL C interface

	Double-precision	Single-precision
Maximum value	$2^{1023}(2 - 2^{-52}) \simeq (1.80 \times 10^{308})$	$2^{127}(2 - 2^{-23}) \simeq (3.40 \times 10^{38})$
Positive minimum value	$2^{-1022} \simeq (2.23 \times 10^{-308})$	$2^{-126} \simeq (1.17 \times 10^{-38})$
Negative maximum value	$-2^{-1022} \simeq (-2.23 \times 10^{-308})$	$-2^{-126} \simeq (-1.17 \times 10^{-38})$
Minimum value	$-2^{1023}(2 - 2^{-52}) \simeq (-1.80 \times 10^{308})$	$-2^{127}(2 - 2^{-23}) \simeq (-3.40 \times 10^{38})$
Overflow decision value	Maximum value $\times 10^{-3}$	
Underflow decision value	Positive minimum value $\times 10^3$	

- (b) The storage method used for the arrays a and ja is as follows.

Figure 4-2 Storage format for Input Data





- (c) The input values of array a and array ja are partially changed in order to optimize performance.
- (d) m can take an arbitrary value as long as it is equal to or greater than the maximum of the numbers of nonzero elements in a row, but it is recommended from the viewpoint of calculation cost to make the padded area as little as possible by setting m as close as possible to the maximum of the numbers of nonzero elements in a row.
- (e) A better performance will be achieved by specifying isw = 0 (Performance will degrade remarkably when isw = 1 is specified). And so, isw = 1 should be specified when it is assured that ja satisfies Restriction (d). When isw = 1 is specified, the user should be very careful in giving the coefficient matrix data as input because the check for Restriction (d) is omitted. The validity of the result will not be guaranteed if isw = 0 was specified but ja doesn't satisfy Restriction (d).
- (f) The work area is automatically allocated within this function. If the work area could not be successfully allocated, processing will be aborted and ierr returns the value 5000. In this case, the problem cannot be solved by using this function unless the problem size is reduced or the machine environment is enhanced in memory size.

(7) **Example**

- (a) Problem

Solve  $Au = b$  for the following matrix  $A$  and vector  $b$ :

$$A = \begin{bmatrix} 3 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 5 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 7 & 3 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 3 & 9 & 4 & 0 & 0 & 0 & 0 & 0 \\ 0 & -2 & 0 & 4 & 11 & 5 & 0 & 0 & 0 & 0 \\ 0 & 0 & -3 & 0 & 5 & 13 & 6 & 0 & 0 & 0 \\ 0 & 0 & 0 & -4 & 0 & 6 & 15 & 7 & 0 & -7 \\ 0 & 0 & 0 & 0 & -5 & 0 & 7 & 17 & 8 & 0 \\ 0 & 0 & 0 & 0 & 0 & -6 & 0 & 8 & 19 & 9 \\ 0 & 0 & 0 & 0 & 0 & 0 & -7 & 0 & 9 & 21 \end{bmatrix}, b = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

- (b) Input data

Input Arrays a, ja and b,

lna = 11, n = 10, m = 11, itrmax = 100, epsmax =  $10^{-12}$ , isw = 0, nt = 2

- (c) Main program

```

/*      C interface example for ASL_qxe020 */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    double *a;
    int lna;
    int n;
    int m;
    int *ja;
    double *b;
    double *u;
    int itrmax;
    int itr;
    double epsmax;
    double eps;
    int isw;
    int nt;
    int ierr;
    int i,j;

```

```

printf( "    *** ASL_qxe020 ***\n" );

lna = 11;
n   = 10;
m   = 5;

a = ( double * )malloc((size_t)( sizeof(double) * (lna*m) ));
if ( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}
ja = ( int * )malloc((size_t)( sizeof(int) * (lna*m) ));
if ( ja == NULL )
{
    printf( "no enough memory for array ja\n" );
    return -1;
}
b = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}
u = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( u == NULL )
{
    printf( "no enough memory for array u\n" );
    return -1;
}

for( j=0 ; j < m ; j++ )
{
    for( i=0 ; i < n ; i++ )
    {
        a[i+lna*j] = 0.0;
        ja[i+lna*j] = 0;
    }
}

a[0 ] = 3.0;
a[lna ] = 1.0;

a[1 ] = 5.0;
a[1+lna ] = 1.0;
a[1+lna*2 ] = 2.0;

a[2 ] = 7.0;
a[2+lna ] = 2.0;
a[2+lna*2 ] = 3.0;

a[3 ] = 9.0;
a[3+lna ] = -1.0;
a[3+lna*2 ] = 3.0;
a[3+lna*3 ] = 4.0;

a[4 ] = 11.0;
a[4+lna ] = -2.0;
a[4+lna*2 ] = 4.0;
a[4+lna*3 ] = 5.0;

a[5 ] = 13.0;
a[5+lna ] = -3.0;
a[5+lna*2 ] = 5.0;
a[5+lna*3 ] = 6.0;

a[6 ] = 15.0;
a[6+lna ] = -4.0;
a[6+lna*2 ] = 6.0;
a[6+lna*3 ] = 7.0;
a[6+lna*4 ] = -7.0;

a[7 ] = 17.0;
a[7+lna ] = -5.0;
a[7+lna*2 ] = 7.0;
a[7+lna*3 ] = 8.0;

a[8 ] = 19.0;
a[8+lna ] = -6.0;
a[8+lna*2 ] = 8.0;
a[8+lna*3 ] = 9.0;

a[9 ] = 21.0;
a[9+lna ] = -7.0;
a[9+lna*2 ] = 9.0;

ja[0 ] = 1;
ja[lna ] = 2;

ja[1 ] = 2;
ja[1+lna ] = 1;
ja[1+lna*2 ] = 3;

```

```

ja[2      ] = 3;
ja[2+lna  ] = 2;
ja[2+lna*2] = 4;

ja[3      ] = 4;
ja[3+lna  ] = 1;
ja[3+lna*2] = 3;
ja[3+lna*3] = 5;

ja[4      ] = 5;
ja[4+lna  ] = 2;
ja[4+lna*2] = 4;
ja[4+lna*3] = 6;

ja[5      ] = 6;
ja[5+lna  ] = 3;
ja[5+lna*2] = 5;
ja[5+lna*3] = 7;

ja[6      ] = 7;
ja[6+lna  ] = 4;
ja[6+lna*2] = 6;
ja[6+lna*3] = 8;
ja[6+lna*4] = 10;

ja[7      ] = 8;
ja[7+lna  ] = 5;
ja[7+lna*2] = 7;
ja[7+lna*3] = 9;

ja[8      ] = 9;
ja[8+lna  ] = 6;
ja[8+lna*2] = 8;
ja[8+lna*3] = 10;

ja[9      ] = 10;
ja[9+lna  ] = 7;
ja[9+lna*2] = 9;

for( i=0 ; i < n ; i++ )
{
    b[i] = 1.0;
}

for( i=0 ; i < n ; i++ )
{
    u[i] = 0.0;
}

itrmax = 100;
epsmax = 1.0e-12;
isw    = 1;
nt     = 2;

printf( "\n *** Input ***\n\n" );
printf( "\titrmax = %6d\n", itrmax );
printf( "\tepsmax =%8.3g\n", epsmax );
printf( "\tnt     = %6d\n", nt );

    ierr = ASL_qxe020
(a,lna,n,m,ja,b,u,itrmax,&itr,epsmax,&eps,isw,nt);

printf( "\n *** Output ***\n\n" );
printf( "\tierr  = %6d\n", ierr );
printf( "\titr   = %6d\n", itr );
printf( "\teps   =%8.3g\n", eps );

for( i=0 ; i < n ; i++ )
{
    printf( "\tu[%2d ] =%8.3g\n", i, u[i] );
}

free( a );
free( ja );
free( b );
free( u );

return 0;
}

```

(d) Output results

```
*** ASL_qxe020 ***
***   Input   ***
itrmax =    100
epsmax = 1e-12
nt      =     2
***   Output  ***
ierr    =     0
itr     =    10
eps     =6.19e-16
u[ 0 ] =  0.296
u[ 1 ] =  0.111
u[ 2 ] =  0.0732
u[ 3 ] =  0.0882
u[ 4 ] =  0.0707
u[ 5 ] =  0.0185
u[ 6 ] =  0.104
u[ 7 ] =  0.0335
u[ 8 ] =  0.00671
u[ 9 ] =  0.0795
```

### 4.2.3 ASL\_qxe030, ASL\_pxe030

#### Asymmetric Sparse Matrix (ELLPACK Format) (BiCGSTAB method)

(1) **Function**

ASL\_qxe030 or ASL\_pxe030 uses the scaling preconditioned BiCGSTAB method to solve the simultaneous linear equations  $A\mathbf{u} = \mathbf{b}$  having an asymmetric sparse matrix  $A$  as the coefficient matrix.

(2) **Usage**

Double precision:

```
ierr = ASL_qxe030 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

Single precision:

```
ierr = ASL_pxe030 (a,lna,n,m,ja,b,u,itrmax, &itr,epsmax, &eps,isw,nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I: {int as for 32bit Integer}  
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$l_n \times m$	Input	Array of nonzero element values of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
2	lna	I	1	Input	Adjustable dimension of array a and ja
3	n	I	1	Input	Order of matrix $A$
4	m	I	1	Input	m denotes the column number of both array a and ja (See Note (d))
5	ja	$I^*$	$l_n \times m$	Input	Array storing the nonzero structure data of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
6	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Right-hand side vector $\mathbf{b}$ of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
7	u	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Solution vector $\mathbf{u}$ of the simultaneous linear equations $A\mathbf{u} = \mathbf{b}$
8	itrmax	I	1	Input	Maximum number of iterations (Default value: n)
9	itr	$I^*$	1	Output	Actual number of iterations
10	epsmax	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Truncation residual norm Default value: $\left\{ \begin{array}{l} 10^{-12} \text{ (doubleprecision)} \\ 10^{-6} \text{ (singleprecision)} \end{array} \right\}$
11	eps	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	1	Output	Final residual norm
12	isw	I	1	Input	Processing switch(Default value: 0) (See Note (e)) isw= 0 : Check whether the values ja are all distinct. isw= 1 : Does not check whether the values ja are all distinct.
13	nt	I	1	Input	Number of tasks to be generated (Default value: 1)
14	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $1 \leq n \leq \text{lna}$
- (b)  $1 \leq m \leq n$
- (c)
  - $\text{ja}[i-1] = i, \text{ja}[(i-1) + \text{lna} \times (j-1)] \neq i$  ( $i = 1, \dots, n; j = 2, \dots, j_i$ ),  
 $1 \leq \text{ja}[(i-1) + \text{lna} \times (j-1)] \leq n$  ( $j = 2, \dots, j_i$ ),  
 where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
  - $\text{ja}[(i-1) + \text{lna} \times j_i] = 0$ , if  $j_i < m$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (d)  $\text{ja}[(i-1) + \text{lna} \times (j-1)]$  ( $j = 1, \dots, j_i$ ) are all distinct for a fixed value  $i$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (e)  $\text{a}[i-1] \neq 0.0$  ( $i = 1, \dots, n$ )
- (f)  $\text{itrmax} \geq 1$
- (g)  $\text{epsmax} > \text{underflowdecisionvalue}$
- (h)  $\text{isw} \in \{0, 1\}$
- (i)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	Restriction (f) was not satisfied.	Processing continues with $\text{itrmax} = n$ .
1200	Restriction (g) was not satisfied.	Processing continues with $\text{epsmax} = \begin{cases} 10^{-12} & (\text{doubleprecision}) \\ 10^{-6} & (\text{singleprecision}) \end{cases}$ .
1300	Restriction (h) was not satisfied.	$\text{isw}$ has the input value 0. Processing continues on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
1400	Restriction (i) was not satisfied.	Processing continues with $\text{nt} = 1$ .
2000	The absolute norm of the right-hand side vector $\mathbf{b}$ is less than the underflow decision value.	$\text{u}[i-1] \leftarrow 0.0$ ( $i = 1, \dots, n$ ) is set for the solution.
2100	Restriction (e) was not satisfied.	Processing continues.
2200	$\text{isw}$ was equal to 0.	Processing continues on the assumption that $\text{ja}$ satisfies Restriction (d) (See Note (e)).
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	
3300	Restriction (d) was not satisfied.	
3500	The norm of the Maximum number of iterations has been reached.	

ierr value	Meaning	Processing
4000	a has a diagonal term whose absolute value is smaller than the underflow decision value.	Processing is aborted.
4100	The norm of the right-hand side vector $\mathbf{b}$ is greater than the overflow decision value.	
4210	The norm of the residual $\mathbf{r} = \mathbf{b} - A\mathbf{u}$ is greater than the overflow decision value.	
4220	The relative norm of the residual $\mathbf{r}$ is greater than the overflow decision value.	
4310	$ \rho_i $ is less than the underflow decision value.	
4320	$ \rho_i $ is greater than the overflow decision value.	
4410	$ (\mathbf{r}_0, \mathbf{v}_i) $ is less than the underflow decision value.	
4420	$ (\mathbf{r}_0, \mathbf{v}_i) $ is greater than the overflow decision value.	
4510	$ (\mathbf{t}, \mathbf{t}) $ is less than the underflow decision value.	
4520	$ (\mathbf{t}, \mathbf{t}) $ is greater than the overflow decision value.	
4610	$ \omega_i $ is less than the underflow decision value.	
4620	$ \omega_i $ is greater than the overflow decision value.	
4700	$ \beta $ is greater than the overflow decision value.	
4800	$ \alpha $ is greater than the overflow decision value.	
4900	$ (\mathbf{t}, \mathbf{s}) $ is greater than the overflow decision value.	
5000	Processing to reserve the work area required to execute the basic iterative algorithm failed (See Note (f)).	



(6) **Notes**

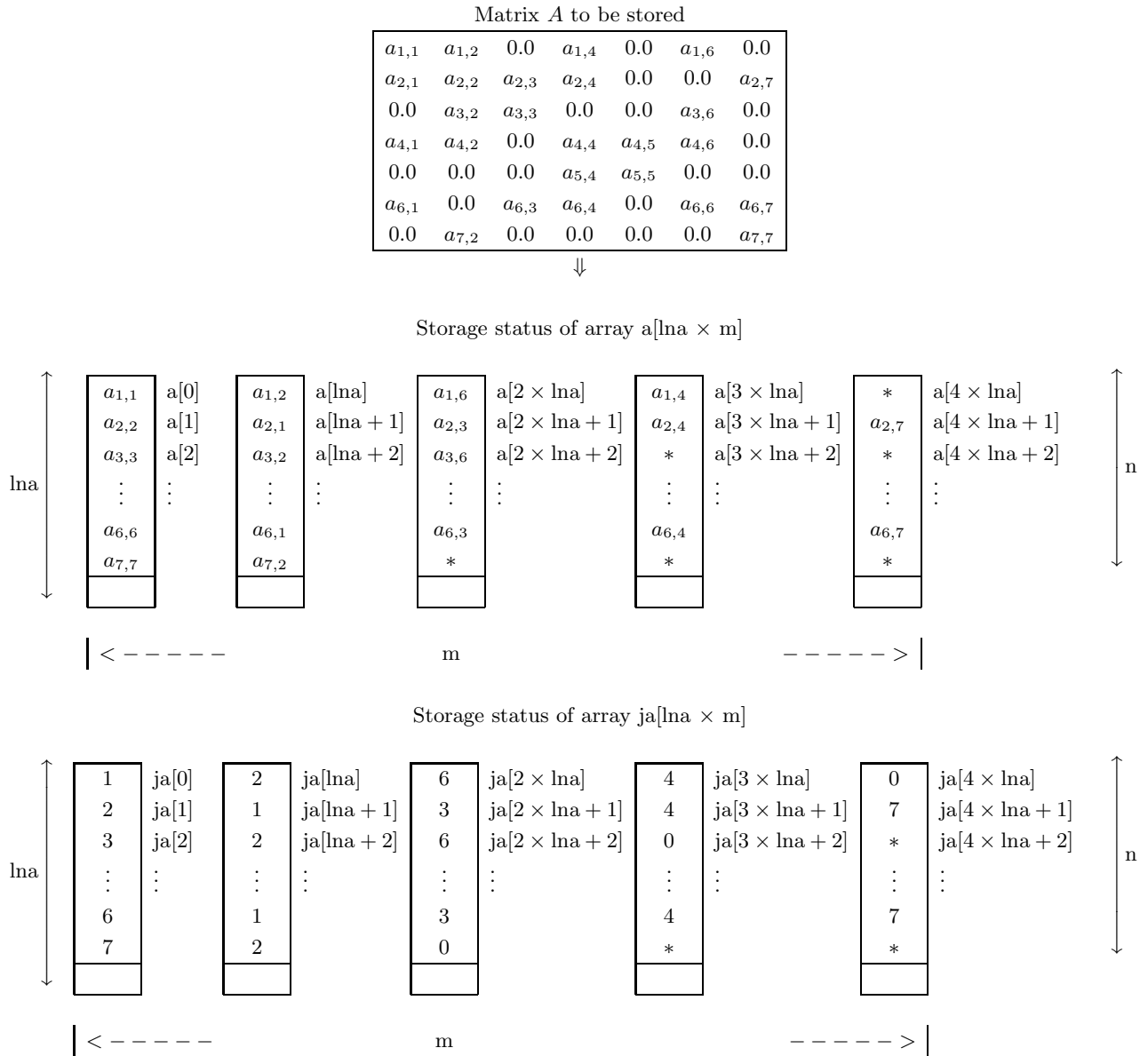
- (a) The table below shows maximum and minimum values of floating point data and so on, that are defined within ASL C interface.

Table 4-4 Constants used in ASL C interface

	Double-precision	Single-precision
Maximum value	$2^{1023}(2 - 2^{-52}) \simeq (1.80 \times 10^{308})$	$2^{127}(2 - 2^{-23}) \simeq (3.40 \times 10^{38})$
Positive minimum value	$2^{-1022} \simeq (2.23 \times 10^{-308})$	$2^{-126} \simeq (1.17 \times 10^{-38})$
Negative maximum value	$-2^{-1022} \simeq (-2.23 \times 10^{-308})$	$-2^{-126} \simeq (-1.17 \times 10^{-38})$
Minimum value	$-2^{1023}(2 - 2^{-52}) \simeq (-1.80 \times 10^{308})$	$-2^{127}(2 - 2^{-23}) \simeq (-3.40 \times 10^{38})$
Overflow decision value	Maximum value $\times 10^{-3}$	
Underflow decision value	Positive minimum value $\times 10^3$	

- (b) The storage method used for the arrays a and ja is as follows.

Figure 4–3 Storage format for Input Data



**Remarks**

- a.  $n$  is order of Matrix  $A$ .
- b.  $\text{lna} \geq n$  must hold.
- c.  $m$  is the column number of Array  $a$ , which contains the nonzero elements of Matrix  $A$ .
- d. Array  $a$  should contain nonzero elements of Matrix  $A$  so that:
  - Diagonal elements are stored in the first column.
  - Nonzero elements in the lower triangular part and the upper triangular part are stored in the second though  $m$ -th columns, with the first one in the second column, one adjacent to the next in each row. Here, it is unnecessary that nonzero elements in each row are stored sequentially.
  - Arbitrary values can be stored in the remaining positions that are marked with ‘\*’.
- e. Array  $ja$  should contain the column indices in Matrix  $A$  of those elements that correspond to the elements contained in Array  $a$ .  
 For those rows in which  $m - 1$  becomes greater than the number of nonzero elements in the lower and upper triangular part, value 0 should be stored in the right neighbor of the rightmost position of the region in  $ja$  in which the column indices of nonzero elements in Matrix  $A$  are stored. Arbitrary values can be stored in the remaining positions that are marked with ‘\*’.

- (c) The input values of array a and array ja are partially changed in order to optimize performance.
- (d) m can take an arbitrary value as long as it is equal to or greater than the maximum of the numbers of nonzero elements in a row, but it is recommended from the viewpoint of calculation cost to make the padded area as little as possible by setting m as close as possible to the maximum of the numbers of nonzero elements in a row.
- (e) A better performance will be achieved by specifying isw = 0 (Performance will degrade remarkably when isw = 1 is specified). And so, isw = 1 should be specified when it is assured that ja satisfies Restriction (d). When isw = 1 is specified, the user should be very careful in giving the coefficient matrix data as input because the check for Restriction (d) is omitted. The validity of the result will not be guaranteed if isw = 0 was specified but ja doesn't satisfy Restriction (d).
- (f) The work area is automatically allocated within this function. If the work area could not be successfully allocated, processing will be aborted and ierr returns the value 5000. In this case, the problem cannot be solved by using this function unless the problem size is reduced or the machine environment is enhanced in memory size.

(7) **Example**

- (a) Problem

Solve  $Au = b$  for the following matrix  $A$  and vector  $b$ :

$$A = \begin{bmatrix} 3 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 5 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 7 & 3 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 3 & 9 & 4 & 0 & 0 & 0 & 0 & 0 \\ 0 & -2 & 0 & 4 & 11 & 5 & 0 & 0 & 0 & 0 \\ 0 & 0 & -3 & 0 & 5 & 13 & 6 & 0 & 0 & 0 \\ 0 & 0 & 0 & -4 & 0 & 6 & 15 & 7 & 0 & -7 \\ 0 & 0 & 0 & 0 & -5 & 0 & 7 & 17 & 8 & 0 \\ 0 & 0 & 0 & 0 & 0 & -6 & 0 & 8 & 19 & 9 \\ 0 & 0 & 0 & 0 & 0 & 0 & -7 & 0 & 9 & 21 \end{bmatrix}, b = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

- (b) Input data

Input Arrays a, ja and b,

lna = 11, n = 10, m = 11, itrmax = 100, epsmax =  $10^{-12}$ , isw = 0, nt = 2

- (c) Main program

```

/*      C interface example for ASL_qxe030 */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    double *a;
    int lna;
    int n;
    int m;
    int *ja;
    double *b;
    double *u;
    int itrmax;
    int itr;
    double epsmax;
    double eps;
    int isw;
    int nt;
    int ierr;
    int i,j;

```

```

printf( "    *** ASL_qxe030 ***\n" );

lna = 11;
n   = 10;
m   = 5;

a = ( double * )malloc((size_t)( sizeof(double) * (lna*m) ));
if ( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}
ja = ( int * )malloc((size_t)( sizeof(int) * (lna*m) ));
if ( ja == NULL )
{
    printf( "no enough memory for array ja\n" );
    return -1;
}
b = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}
u = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( u == NULL )
{
    printf( "no enough memory for array u\n" );
    return -1;
}

for( j=0 ; j < m ; j++ )
{
    for( i=0 ; i < n ; i++ )
    {
        a[i+lna*j] = 0.0;
        ja[i+lna*j] = 0;
    }
}

a[0 ] = 3.0;
a[lna ] = 1.0;

a[1 ] = 5.0;
a[1+lna ] = 1.0;
a[1+lna*2 ] = 2.0;

a[2 ] = 7.0;
a[2+lna ] = 2.0;
a[2+lna*2 ] = 3.0;

a[3 ] = 9.0;
a[3+lna ] = -1.0;
a[3+lna*2 ] = 3.0;
a[3+lna*3 ] = 4.0;

a[4 ] = 11.0;
a[4+lna ] = -2.0;
a[4+lna*2 ] = 4.0;
a[4+lna*3 ] = 5.0;

a[5 ] = 13.0;
a[5+lna ] = -3.0;
a[5+lna*2 ] = 5.0;
a[5+lna*3 ] = 6.0;

a[6 ] = 15.0;
a[6+lna ] = -4.0;
a[6+lna*2 ] = 6.0;
a[6+lna*3 ] = 7.0;
a[6+lna*4 ] = -7.0;

a[7 ] = 17.0;
a[7+lna ] = -5.0;
a[7+lna*2 ] = 7.0;
a[7+lna*3 ] = 8.0;

a[8 ] = 19.0;
a[8+lna ] = -6.0;
a[8+lna*2 ] = 8.0;
a[8+lna*3 ] = 9.0;

a[9 ] = 21.0;
a[9+lna ] = -7.0;
a[9+lna*2 ] = 9.0;

ja[0 ] = 1;
ja[lna ] = 2;

ja[1 ] = 2;
ja[1+lna ] = 1;
ja[1+lna*2 ] = 3;

```

```

ja[2      ] = 3;
ja[2+lna  ] = 2;
ja[2+lna*2] = 4;

ja[3      ] = 4;
ja[3+lna  ] = 1;
ja[3+lna*2] = 3;
ja[3+lna*3] = 5;

ja[4      ] = 5;
ja[4+lna  ] = 2;
ja[4+lna*2] = 4;
ja[4+lna*3] = 6;

ja[5      ] = 6;
ja[5+lna  ] = 3;
ja[5+lna*2] = 5;
ja[5+lna*3] = 7;

ja[6      ] = 7;
ja[6+lna  ] = 4;
ja[6+lna*2] = 6;
ja[6+lna*3] = 8;
ja[6+lna*4] = 10;

ja[7      ] = 8;
ja[7+lna  ] = 5;
ja[7+lna*2] = 7;
ja[7+lna*3] = 9;

ja[8      ] = 9;
ja[8+lna  ] = 6;
ja[8+lna*2] = 8;
ja[8+lna*3] = 10;

ja[9      ] = 10;
ja[9+lna  ] = 7;
ja[9+lna*2] = 9;

for( i=0 ; i < n ; i++ )
{
    b[i] = 1.0;
}

for( i=0 ; i < n ; i++ )
{
    u[i] = 0.0;
}

itrmax = 100;
epsmax = 1.0e-12;
isw    = 1;
nt     = 2;

printf( "\n *** Input ***\n\n" );
printf( "\titrmax = %6d\n", itrmax );
printf( "\tepsmax =%8.3g\n", epsmax );
printf( "\tnt     = %6d\n", nt );

    ierr = ASL_qxe030
(a,lna,n,m,ja,b,u,itrmax,&itr,epsmax,&eps,isw,nt);

printf( "\n *** Output ***\n\n" );
printf( "\tierr  = %6d\n", ierr );
printf( "\titr   = %6d\n", itr );
printf( "\teps   =%8.3g\n", eps );

for( i=0 ; i < n ; i++ )
{
    printf( "\tu[%2d ] =%8.3g\n", i, u[i] );
}

free( a );
free( ja );
free( b );
free( u );

return 0;
}

```

(d) Output results

```
*** ASL_qxe030 ***
***   Input   ***
itrmax =    100
epsmax = 1e-12
nt      =     2
***   Output  ***
ierr   =     0
itr    =     10
eps    =3.88e-16
u[ 0 ] =  0.296
u[ 1 ] =  0.111
u[ 2 ] =  0.0732
u[ 3 ] =  0.0882
u[ 4 ] =  0.0707
u[ 5 ] =  0.0185
u[ 6 ] =  0.104
u[ 7 ] =  0.0335
u[ 8 ] =  0.00671
u[ 9 ] =  0.0795
```

#### 4.2.4 ASL\_qxe040, ASL\_pxe040 Asymmetric Sparse Matrix (ELLPACK Format) (GMRES(m) method)

(1) **Function**

ASL\_qxe040 or ASL\_pxe040 uses the scaling preconditioned GMRES(m) method to solve the simultaneous linear equations  $Au = b$  having an asymmetric sparse matrix  $A$  as the coefficient matrix.

(2) **Usage**

Double precision:

ierr = ASL\_qxe040 (a,lna,n,m,ja,b,u,itrmax, &itr,mgmrs,epsmax, &eps,isw,nt);

Single precision:

ierr = ASL\_pxe040 (a,lna,n,m,ja,b,u,itrmax, &itr,mgmrs,epsmax, &eps,isw,nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex  
 R:Single precision real    C:Single precision complex    I: {int as for 32bit Integer}

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lna \times m$	Input	Array of nonzero element values of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
2	lna	I	1	Input	Adjustable dimension of array a and ja
3	n	I	1	Input	Order of matrix $A$
4	m	I	1	Input	m denotes the column number of both array a and ja (See Note (d))
5	ja	$I^*$	$lna \times m$	Input	Array storing the nonzero structure data of the coefficient matrix (For the storage format(See Note (b)))
				Output	Values updated for optimization (See Note (c))
6	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Input	Right-hand side vector $b$ of the simultaneous linear equations $Au = b$
7	u	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Solution vector $u$ of the simultaneous linear equations $Au = b$
8	itrmax	I	1	Input	Maximum number of iterations (Default value: n)
9	itr	$I^*$	1	Output	Actual number of iterations
10	mgmrs	I	1	Input	GMRES(m) parameter m (Default value: 10)
11	epsmax	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Truncation residual norm Default value: $\begin{Bmatrix} 10^{-12} & (\text{doubleprecision}) \\ 10^{-6} & (\text{singleprecision}) \end{Bmatrix}$

No.	Argument and Return Value	Type	Size	Input/Output	Contents
12	eps	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	1	Output	Final residual norm
13	isw	I	1	Input	Processing switch(Default value: 0) (See Note (e)) isw= 0 : Check whether the values ja are all distinct. isw= 1 : Does not check whether the values ja are all distinct.
14	nt	I	1	Input	Number of tasks to be generated (Default value: 1)
15	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $1 \leq n \leq \text{lna}$
- (b)  $1 \leq m \leq n$
- (c)
  - $\text{ja}[i - 1] = i, \text{ja}[(i - 1) + \text{lna} \times (j - 1)] \neq i$  ( $i = 1, \dots, n; j = 2, \dots, j_i$ ),  
 $1 \leq \text{ja}[(i - 1) + \text{lna} \times (j - 1)] \leq n$  ( $j = 2, \dots, j_i$ ) ,  
where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
  - $\text{ja}[(i - 1) + \text{lna} \times j_i] = 0$ , if  $j_i < m$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (d)  $\text{ja}[(i - 1) + \text{lna} \times (j - 1)]$  ( $j = 1, \dots, j_i$ ) are all distinct for a fixed value  $i$ , where  $j_i$  ( $i = 1, \dots, n$ ) is the number of nonzero elements that are contained in the  $i$ -th column of Matrix  $A$ .
- (e)  $\text{a}[i - 1] \neq 0.0$  ( $i = 1, \dots, n$ )
- (f)  $\text{itrmax} \geq 1$
- (g)  $\text{mgmrs} \geq 1$
- (h)  $\text{epsmax} > \text{underflowdecisionvalue}$
- (i)  $\text{isw} \in \{0, 1\}$
- (j)  $\text{nt} \geq 1$



## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	Restriction (f) was not satisfied.	Processing continues with itrmax = n.
1100	Restriction (g) was not satisfied.	Processing continues with mgmrs = 10.
1200	Restriction (h) was not satisfied.	Processing continues with epsmax = $\left\{ \begin{array}{l} 10^{-12} \text{ (doubleprecision)} \\ 10^{-6} \text{ (singleprecision)} \end{array} \right\}$ .
1300	Restriction (i) was not satisfied.	isw has the input value 0. Processing continues on the assumption that ja satisfies Restriction (d) (See Note (e)).
1400	Restriction (j) was not satisfied.	Processing continues with nt = 1.
2000	The absolute norm of the right-hand side vector $\mathbf{b}$ is less than the underflow decision value.	$u[i - 1] \leftarrow 0.0$ ( $i = 1, \dots, n$ ) is set for the solution.
2100	Restriction (e) was not satisfied.	Processing continues.
2200	isw was equal to 0.	Processing continues on the assumption that ja satisfies Restriction (d) (See Note (e)).
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	
3300	Restriction (d) was not satisfied.	
3500	The norm of the Maximum number of iterations has been reached.	The result obtained at that time is returned.
4000	a has a diagonal term whose absolute value is smaller than the underflow decision value.	Processing is aborted.
4100	The norm of the right-hand side vector $\mathbf{b}$ is greater than the overflow decision value.	
4210	The norm of the initial residual $\mathbf{r} = \mathbf{b} - A\mathbf{u}$ is greater than the overflow decision value.	
4220	The norm of the residual $\mathbf{r}$ of the final solution is greater than the overflow decision value.	
4310	$\ AM^{-1}\mathbf{v}_i\ _2^2$ is less than the underflow decision value. In this case, it is very possible that $AM^{-1}$ is degenerate.	The solution obtained by the $(i - 1)$ iteration is returned, and processing is aborted.
4320	$\ AM^{-1}\mathbf{v}_i\ _2^2$ is greater than the overflow decision value.	Processing is aborted.

ierr value	Meaning	Processing
4410	$\ w\ _2^2$ after Gram-Schmidt orthogonalization became less than the underflow decision value, and processing could not be continued. This indicates that the iterative solution has already converged. It is very possible that the convergence decision conditions are too strict.	The solution obtained by that time is returned, and processing is aborted.
5000	Processing to reserve the work area required to execute the basic iterative algorithm failed (See Note (f)).	

(6) **Notes**

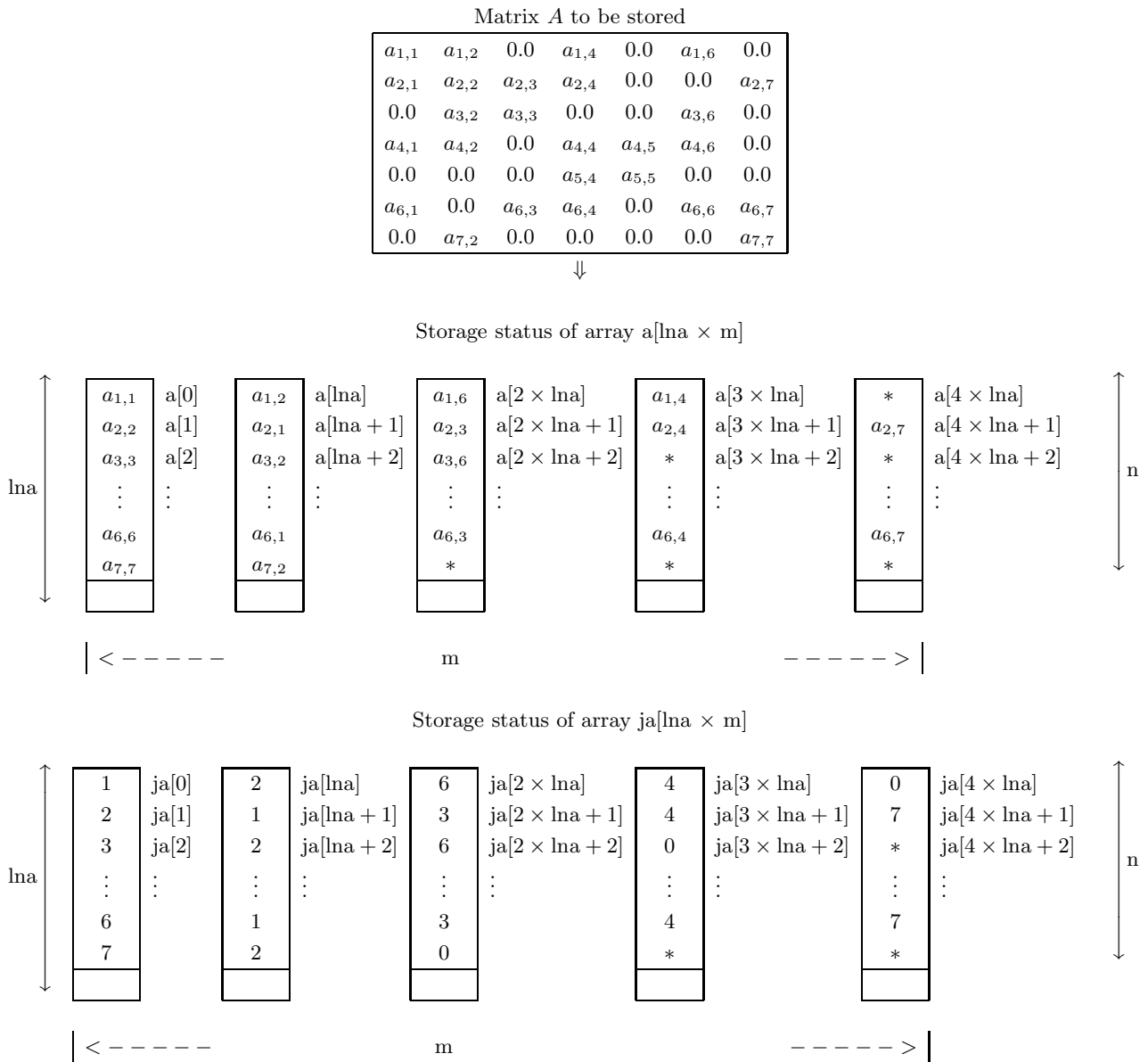
- (a) The table below shows maximum and minimum values of floating point data and so on, that are defined within ASL C interface.

Table 4-5 Constants used in ASL C interface

	Double-precision	Single-precision
Maximum value	$2^{1023}(2 - 2^{-52}) \simeq (1.80 \times 10^{308})$	$2^{127}(2 - 2^{-23}) \simeq (3.40 \times 10^{38})$
Positive minimum value	$2^{-1022} \simeq (2.23 \times 10^{-308})$	$2^{-126} \simeq (1.17 \times 10^{-38})$
Negative maximum value	$-2^{-1022} \simeq (-2.23 \times 10^{-308})$	$-2^{-126} \simeq (-1.17 \times 10^{-38})$
Minimum value	$-2^{1023}(2 - 2^{-52}) \simeq (-1.80 \times 10^{308})$	$-2^{127}(2 - 2^{-23}) \simeq (-3.40 \times 10^{38})$
Overflow decision value	Maximum value $\times 10^{-3}$	
Underflow decision value	Positive minimum value $\times 10^3$	

- (b) The storage method used for the arrays a and ja is as follows.

Figure 4-4 Storage format for Input Data



**Remarks**

- a.  $n$  is order of Matrix  $A$ .
- b.  $lna \geq n$  must hold.
- c.  $m$  is the column number of Array  $a$ , which contains the nonzero elements of Matrix  $A$ .
- d. Array  $a$  should contain nonzero elements of Matrix  $A$  so that:
  - Diagonal elements are stored in the first column.
  - Nonzero elements in the lower triangular part and the upper triangular part are stored in the second through  $m$ -th columns, with the first one in the second column, one adjacent to the next in each row. Here, it is unnecessary that nonzero elements in each row are stored sequentially.
  - Arbitrary values can be stored in the remaining positions that are marked with '\*'.
- e. Array  $ja$  should contain the column indices in Matrix  $A$  of those elements that correspond to the elements contained in Array  $a$ .  
 For those rows in which  $m - 1$  becomes greater than the number of nonzero elements in the lower and upper triangular part, value 0 should be stored in the right neighbor of the rightmost position of the region in  $ja$  in which the column indices of nonzero elements in Matrix  $A$  are stored. Arbitrary values can be stored in the remaining positions that are marked with '\*'.

- (c) The input values of array a and array ja are partially changed in order to optimize performance.
- (d) m can take an arbitrary value as long as it is equal to or greater than the maximum of the numbers of nonzero elements in a row, but it is recommended from the viewpoint of calculation cost to make the padded area as little as possible by setting m as close as possible to the maximum of the numbers of nonzero elements in a row.
- (e) A better performance will be achieved by specifying isw = 0 (Performance will degrade remarkably when isw = 1 is specified). And so, isw = 1 should be specified when it is assured that ja satisfies Restriction (d). When isw = 1 is specified, the user should be very careful in giving the coefficient matrix data as input because the check for Restriction (d) is omitted. The validity of the result will not be guaranteed if isw = 0 was specified but ja doesn't satisfy Restriction (d).
- (f) The work area is automatically allocated within this function. If the work area could not be successfully allocated, processing will be aborted and ierr returns the value 5000. In this case, the problem cannot be solved by using this function unless the problem size is reduced or the machine environment is enhanced in memory size.

(7) **Example**

- (a) Problem

Solve  $Au = b$  for the following matrix  $A$  and vector  $b$ :

$$A = \begin{bmatrix} 3 & 1 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 1 & 5 & 2 & 0 & 0 & 0 & 0 & 0 & 0 & 0 \\ 0 & 2 & 7 & 3 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 3 & 9 & 4 & 0 & 0 & 0 & 0 & 0 \\ 0 & -2 & 0 & 4 & 11 & 5 & 0 & 0 & 0 & 0 \\ 0 & 0 & -3 & 0 & 5 & 13 & 6 & 0 & 0 & 0 \\ 0 & 0 & 0 & -4 & 0 & 6 & 15 & 7 & 0 & -7 \\ 0 & 0 & 0 & 0 & -5 & 0 & 7 & 17 & 8 & 0 \\ 0 & 0 & 0 & 0 & 0 & -6 & 0 & 8 & 19 & 9 \\ 0 & 0 & 0 & 0 & 0 & 0 & -7 & 0 & 9 & 21 \end{bmatrix}, b = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}$$

- (b) Input data

Input Arrays a, ja and b,

lna = 11, n = 10, m = 11, itrmax = 100, mgmrs = 5, epsmax =  $10^{-12}$ , isw = 0, nt = 2

- (c) Main program

```

/*      C interface example for ASL_qxe040 */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    double *a;
    int lna;
    int n;
    int m;
    int *ja;
    double *b;
    double *u;
    int itrmax;
    int itr;
    int mgmrs;
    double epsmax;
    double eps;
    int isw;
    int nt;
    int ierr;

```

```

int i,j;
printf( "    *** ASL_qxe040 ***\n" );
lna = 11;
n    = 10;
m    = 5;

a = ( double * )malloc((size_t)( sizeof(double) * (lna*m) ));
if ( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}
ja = ( int * )malloc((size_t)( sizeof(int) * (lna*m) ));
if ( ja == NULL )
{
    printf( "no enough memory for array ja\n" );
    return -1;
}
b = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}
u = ( double * )malloc((size_t)( sizeof(double) * n ));
if ( u == NULL )
{
    printf( "no enough memory for array u\n" );
    return -1;
}

for( j=0 ; j < m ; j++ )
{
    for( i=0 ; i < n ; i++ )
    {
        a[i+lna*j] = 0.0;
        ja[i+lna*j] = 0;
    }
}

a[0 ] = 3.0;
a[ lna ] = 1.0;

a[1 ] = 5.0;
a[1+lna ] = 1.0;
a[1+lna*2] = 2.0;

a[2 ] = 7.0;
a[2+lna ] = 2.0;
a[2+lna*2] = 3.0;

a[3 ] = 9.0;
a[3+lna ] = -1.0;
a[3+lna*2] = 3.0;
a[3+lna*3] = 4.0;

a[4 ] = 11.0;
a[4+lna ] = -2.0;
a[4+lna*2] = 4.0;
a[4+lna*3] = 5.0;

a[5 ] = 13.0;
a[5+lna ] = -3.0;
a[5+lna*2] = 5.0;
a[5+lna*3] = 6.0;

a[6 ] = 15.0;
a[6+lna ] = -4.0;
a[6+lna*2] = 6.0;
a[6+lna*3] = 7.0;
a[6+lna*4] = -7.0;

a[7 ] = 17.0;
a[7+lna ] = -5.0;
a[7+lna*2] = 7.0;
a[7+lna*3] = 8.0;

a[8 ] = 19.0;
a[8+lna ] = -6.0;
a[8+lna*2] = 8.0;
a[8+lna*3] = 9.0;

a[9 ] = 21.0;
a[9+lna ] = -7.0;
a[9+lna*2] = 9.0;

ja[0 ] = 1;
ja[ lna ] = 2;

ja[1 ] = 2;
ja[1+lna ] = 1;

```

```

ja[1+lna*2] = 3;
ja[2      ] = 3;
ja[2+lna  ] = 2;
ja[2+lna*2] = 4;

ja[3      ] = 4;
ja[3+lna  ] = 1;
ja[3+lna*2] = 3;
ja[3+lna*3] = 5;

ja[4      ] = 5;
ja[4+lna  ] = 2;
ja[4+lna*2] = 4;
ja[4+lna*3] = 6;

ja[5      ] = 6;
ja[5+lna  ] = 3;
ja[5+lna*2] = 5;
ja[5+lna*3] = 7;

ja[6      ] = 7;
ja[6+lna  ] = 4;
ja[6+lna*2] = 6;
ja[6+lna*3] = 8;
ja[6+lna*4] = 10;

ja[7      ] = 8;
ja[7+lna  ] = 5;
ja[7+lna*2] = 7;
ja[7+lna*3] = 9;

ja[8      ] = 9;
ja[8+lna  ] = 6;
ja[8+lna*2] = 8;
ja[8+lna*3] = 10;

ja[9      ] = 10;
ja[9+lna  ] = 7;
ja[9+lna*2] = 9;

for( i=0 ; i < n ; i++ )
{
    b[i] = 1.0;
}

for( i=0 ; i < n ; i++ )
{
    u[i] = 0.0;
}

itrmax = 100;
mgmrs  = 10;
epsmax = 1.0e-12;
isw    = 1;
nt     = 2;

printf( "\n ***   Input   ***\n\n" );
printf( "\titrmax = %6d\n", itrmax );
printf( "\tepsmax =%8.3g\n", epsmax );
printf( "\tmgmrs  = %6d\n", mgmrs );
printf( "\tnt    = %6d\n", nt );

    ierr = ASL_qxe040
(a,lna,n,m,ja,b,u,itrmax,&itr,mgmrs,epsmax,&eps,isw,nt);

printf( "\n ***   Output  ***\n\n" );
printf( "\tierr  = %6d\n", ierr );
printf( "\titr   = %6d\n", itr );
printf( "\teps   =%8.3g\n\n", eps );

for( i=0 ; i < n ; i++ )
{
    printf( "\tu[%2d ] =%8.3g\n", i, u[i] );
}

free( a );
free( ja );
free( b );
free( u );

return 0;
}

```

(d) Output results

```
*** ASL_qxe040 ***  
***   Input   ***  
itrmax =    100  
epsmax = 1e-12  
mgmrs  =    10  
nt     =     2  
***   Output  ***  
ierr   =     0  
itr    =    10  
eps    =3.02e-16  
u[ 0 ] =  0.296  
u[ 1 ] =  0.111  
u[ 2 ] =  0.0732  
u[ 3 ] =  0.0882  
u[ 4 ] =  0.0707  
u[ 5 ] =  0.0185  
u[ 6 ] =  0.104  
u[ 7 ] =  0.0335  
u[ 8 ] =  0.00671  
u[ 9 ] =  0.0795
```

## Chapter 5

---

# EIGENVALUES AND EIGENVECTORS

## 5.1 INTRODUCTION

This chapter describes functions that obtain eigenvalues and eigenvectors of symmetric matrices.

**Function described in this chapter divides up and allocates internal processing among threads and executes allocated processing in parallel.**

In the standard eigenvalue problem, obtain the value  $\lambda$  and corresponding vector  $\mathbf{x}$  which satisfy the following equation for a given matrix  $A$ :

$$A\mathbf{x} = \lambda\mathbf{x}.$$

The value  $\lambda$  and the corresponding vector  $\mathbf{x}$  are called an eigenvalue and the corresponding eigenvector, respectively.

In generalized eigenvalue problem, obtain the value  $\lambda$  and corresponding vector  $\mathbf{x}$  which satisfy one of the following equations for given matrices  $A$  and  $B$ :

$$A\mathbf{x} = \lambda B\mathbf{x}$$

or

$$AB\mathbf{x} = \lambda\mathbf{x} \text{ ( both } A \text{ and } B \text{ are Hermitian, } B \text{ is positive definite )}$$

or

$$BA\mathbf{x} = \lambda\mathbf{x} \text{ ( both } A \text{ and } B \text{ are Hermitian, } B \text{ is positive definite )}.$$

These  $\lambda$  and  $\mathbf{x}$  are also called an eigenvalue and an eigenvector. If both  $A$  and  $B$  are Hermitian and  $B$  is positive definite, all the eigenvalues are real and the eigenvectors for different eigenvalues are orthogonal for each other.

The functions contained in this chapter provide functions corresponding to the following four categories.

**All eigenvalues and all eigenvectors:** Obtain all eigenvalues and the corresponding eigenvectors.

**All eigenvalues:** Obtain all eigenvalues only.

**Eigenvalues and eigenvectors:** Obtain a number of the largest or smallest eigenvalues and the corresponding eigenvectors.

**Eigenvalues:** Obtain a number of the largest or smallest eigenvalues.



### 5.1.1 Notes

- (1) In general, functions corresponding to ‘All eigenvalues and all eigenvectors’ or ‘Eigenvalues and eigenvectors’ require more processing time and memory than functions corresponding to ‘All eigenvalues’ or ‘Eigenvalues’ respectively.
- (2) In general, it is more efficient to use functions corresponding to ‘Eigenvalues and eigenvectors’ or ‘Eigenvalues’ if you want to obtain no more than 20% of the total number of eigenvalues. To obtain more than 20% of the total number of eigenvalues, functions corresponding to ‘All eigenvalues and all eigenvectors’ or ‘All eigenvalues’ require less processing time.
- (3) In this library, the functions of the generalized eigenvalue require the condition that  $B$  is positive definite. In the following cases, however, the eigenvalues and the eigenvectors can be obtained even if  $B$  is not positive definite.

- (a) Matrix  $B$  is not positive definite but matrix  $A$  is positive definite

$$Bv = \lambda^{-1}Av$$

gives non-zero eigenvalues.

- (b) Both of  $A$  and  $B$  are not positive definite but  $A + B$  is positive definite

$$Av = \frac{\lambda}{1 + \lambda}(A + B)v$$

gives the eigenvalues which are not  $-1$ .

- (4) If the input matrix is a symmetric matrix or Hermitian matrix, the use of the exclusive functions requires less processing time.
- (5) Since the parallel processing overhead significantly affects the computation cost if the order of the matrix is small, performance may be worse than when a non-parallel function is used.

## 5.1.2 Algorithms Used

### 5.1.2.1 Transforming a real symmetric matrix to a real symmetric tridiagonal matrix

The Householder method is used to transform an  $n \times n$  real symmetric matrix  $A$  to a real symmetric tridiagonal matrix  $T$ . That is,  $A_1 = A$  is assumed, and for  $k = 1, 2, \dots, n - 2$ , a vector  $u_k$  is taken so that:

$$H_k = \frac{1}{2} \mathbf{u}_k^T \mathbf{u}_k$$

$$P_k = I - \frac{\mathbf{u}_k \mathbf{u}_k^T}{H_k}$$

and all elements below the subdiagonal component in column  $k$  of:

$$A_{k+1} = P_k A_k P_k$$

become 0.  $A_{n-1}$  becomes the obtained real symmetric tridiagonal matrix. In addition, the transformation matrix  $P_k$  is an orthogonal symmetric matrix.

### 5.1.2.2 Transforming a Hermitian matrix to a real symmetric tridiagonal matrix

First, the Householder method is used to transform an  $n \times n$  Hermitian matrix  $A$  to a Hermitian tridiagonal matrix  $S$ .

$$S = P_{n-2} \cdots P_2 P_1 A P_1 P_2 \cdots P_{n-2}$$

Then a regular complex diagonal matrix  $D$  is used (similarity transformation) to transform matrix  $S$  to a real symmetric tridiagonal matrix  $T$ .

$$T = D^* S D$$

### 5.1.2.3 The Householder transformation by block algorithm

As for the Householder transformation, the block algorithm is used. This method simplifies the update of a matrix by applying the lump sum of a rank-one matrix update that transforms the original real symmetric matrix into a real symmetric tridiagonal matrix. Let  $A_{k+1}$  be the symmetric matrix that is generated after similarity transformations are performed for  $k$  times on the original symmetric matrix  $A$ . Then it holds that:

$$A_{k+1} = P_k A_k P_k = A_k - \mathbf{u}_k \mathbf{v}_k^T - \mathbf{v}_k \mathbf{u}_k^T$$

where  $P_k$  is an orthogonal matrix. Also the following relationship holds:

$$A_1 = A$$

$$P_k = I - \frac{\mathbf{u}_k \mathbf{u}_k^T}{H_k}$$

$$H_k = \frac{1}{2} \mathbf{u}_k^T \mathbf{u}_k$$

$$\mathbf{y}_k = A_k \mathbf{u}_k$$

$$\mathbf{v}_k = \frac{(\mathbf{y}_k - \frac{(\mathbf{u}_k^T \mathbf{y}_k) \mathbf{u}_k}{2H_k})}{H_k}$$

The Householder transformation updates the symmetric matrix twice for one similarity transformation. The  $A_{k+1}$  can be expressed without using  $A_k$  explicitly.

$$A_{k+1} = A_{k-1} - \mathbf{u}_{k-1}\mathbf{v}_{k-1}^T - \mathbf{v}_{k-1}\mathbf{u}_{k-1}^T - \mathbf{u}_k\mathbf{v}_k^T - \mathbf{v}_k\mathbf{u}_k^T$$

Similarly, matrix  $A_{p+1}$  after the  $p$  times of mirror transformation becomes:

$$A_{p+1} = A_1 - \sum_{i=1}^p (\mathbf{u}_i\mathbf{v}_i^T + \mathbf{v}_i\mathbf{u}_i^T)$$

Therefore, matrix  $A_{p+1}$  can be computed at a higher speed if matrix  $A_1$  is updated in the lump, once per  $2p$  matrices. If the original matrix is Hermitian, the transpose notation “ $T$ ” should be replaced with the Hermite conjugate notation “ $*$ ”. For details, refer to (3) and (4) in the reference bibliography.

#### 5.1.2.4 QR method

For a tridiagonal matrix  $T$ , there is a unitary matrix  $Q$  and an upper triangular matrix  $R$  (for which all diagonal components are positive) such that  $T$  is uniquely decomposed into  $T = QR$ .  $T_k = T$  is assumed,  $T_k$  is decomposed into  $T_k = Q_k R_k$ , and these are multiplied in the reverse order to form:

$$T_{k+1} \leftarrow R_k Q_k = Q_k^* T_k Q_k \text{ (} Q_k^* \text{ is the adjoint matrix of } Q_k \text{)} (k = 1, \dots)$$

If  $T_1, T_2, \dots, T_k, T_{k+1}$  are created, they are all tridiagonal matrices. As  $k \rightarrow \infty$ ,  $T_k$  converges to a diagonal matrix having the eigenvalues of  $T$  as its diagonal elements.

To accelerate convergence in the actual QR method, the values  $\mu_k$  are taken as approximations of the eigenvalues,  $T_k - \mu_k I$  are created in place of  $T_k$  by performing an origin shift, and these are decomposed into:

$$T_k - \mu_k I = Q_k R_k$$

To calculate the approximation of an eigenvalue, consider the case when there is an adjacent eigenvalue (or an eigenvalue having a close absolute value). Let the eigenvalue  $\mu_k$  of the lower-right corner submatrix obtained by the root-free QR method. If  $T_{k+1} = R_k Q_k + \mu_k I$  is created, then  $T_{k+1}$  becomes:

$$T_{k+1} = Q_k^* T_k Q_k$$

After this operation is iterated until the sequence of matrices converges, the values adjusted by the cumulative amount the origin was shifted become the eigenvalues.

The eigenvectors of the original matrix before tridiagonalization are obtained by the following procedures. First, sequentially accumulate the transformation matrices used when obtaining the trigonal matrix  $T$  according to the Householder transformation method. Next, accumulate the transformation matrices  $Q_1, Q_2, \dots, Q_k$  obtained according to the QR method.

#### 5.1.2.5 root-free QR method

The root-free QR method, which eliminates the square root calculations of the QR method, is faster when only seeking the eigenvalues of a real symmetric tridiagonal matrix. Let  $\alpha_1, \dots, \alpha_n$  be the diagonal elements and  $\beta_1, \dots, \beta_{n-1}$  be the subdiagonal elements. Let one of the components of the transformation matrix during the calculation be  $P^{(i)}$  and let  $S_i$  and  $C_i$  be  $\sin \theta$  and  $\cos \theta$  within  $P^{(i)}$ . Although square roots must be computed in the calculations:

$$P_i = \alpha_i C_{i-1} - \beta_{i-1} S_{i-1} C_{i-2}$$

$$S_i = \frac{\beta_i}{\sqrt{P_i^2 + \beta_i^2}}$$

$$C_i = \frac{P_i}{\sqrt{P_i^2 + \beta_i^2}}$$

$$\text{New } \alpha_{i-1} = \alpha_i + P_{i-1}C_{i-2} - P_iC_{i-1}$$

$$\text{New } \beta_{i-2} = S_{i-2}\sqrt{P_{i-1}^2 + \beta_{i-1}^2}$$

If these calculations are performed using the squared formats of each of  $P_i$ ,  $\beta_i$ ,  $S_i$ , and  $C_i$ , and if the following values are assumed:  $C_0 = 1, S_0 = 0, r_1 = \alpha_1, P_1^2 = \alpha_1^2, \alpha_{n+1} = \beta_{n+1} = 0$  then the equations can be expressed as follows:

$$t_i^2 = P_i^2 + \beta_{i+1}^2$$

$$\text{New } \beta_i^2 = S_{i-1}^2 t_i^2$$

$$S_i^2 = \frac{\beta_{i+1}^2}{t_i^2}, C_i^2 = \frac{P_i^2}{t_i^2}$$

$$P_{i+1}^2 = \alpha_{i+1}^2 C_i^2 - 2\alpha_i + S_i^2 \gamma_i + \beta_{i+1}^2 S_i^2 C_{i-1}^2$$

$$\gamma_{i+1} = \alpha_{i+1} C_i^2 = S_i^2 \gamma_i$$

$$\text{New } \alpha_i = \alpha_{i+1} + \gamma_i - \gamma_{i+1}$$

and the calculations can be performed without computing any square roots.

For details, refer to (9) in the reference bibliography.

### 5.1.2.6 Bisection method

The bisection method obtains several of the largest or smallest eigenvalues of a real symmetric tridiagonal matrix  $T$ . If we let  $d_1, d_2, \dots, d_n$  be the diagonal components of  $T$ , let  $s_1, s_2, \dots, s_{n-1}$  be the subdiagonal components, let  $\lambda$  be a variable, and create the sequence of functions:

$$f_0(\lambda) = 1$$

$$f_1(\lambda) = d_1 - \lambda$$

$$f_i(\lambda) = (d_i - \lambda)f_{i-1}(\lambda) - s_{i-1}^2 f_{i-2}(\lambda) \quad (i = 2, \dots, n)$$

then  $f_0(\lambda), f_1(\lambda), \dots, f_m(\lambda)$  is the Sturm sequence. That is, if we let  $L(\lambda)$  be the number of non-matching signs for the successive sequence of functions for a given  $\lambda$ , then this  $L(\lambda)$  is equal to the number of eigenvalues less than  $\lambda$ . To prevent overflow or underflow, if  $g_i(\lambda)$  is actually defined as:

$$g_i(\lambda) = \frac{f_i(\lambda)}{f_{i-1}(\lambda)} \quad (i = 1, 2, \dots, n)$$

$L(\lambda)$  becomes the number of  $g_i(\lambda)$  that are negative. Furthermore,  $g_i(\lambda)$  satisfies the following:

$$g_1(\lambda) = d_1 - \lambda$$

$$g_i = (d_i - \lambda) - \frac{s_{i-1}^2}{g_{i-1}(\lambda)} \quad (i = 2, \dots, n)$$

If  $g_{i-1}(\lambda)=0$ , then  $g_i(\lambda)$  is assumed to be:

$$g_i(\lambda) = (d_i - \lambda) - \frac{|s_{i-1}|}{\varepsilon} \quad (\varepsilon : \text{Units for determining error})$$

Assume that the eigenvalues of  $T$  satisfy:

$$\lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_m$$

From the Gerschgorin theorem, the lower limit ( $x_{\min}$ ) and upper limit ( $x_{\max}$ ) of all the eigenvalues are given by:

$$x_{\max} = \max(d_i + (|s_{i-1}| + |s_i|)) \quad (1 \leq i \leq n)$$

$$x_{\min} = \min(d_i - (|s_{i-1}| + |s_i|)) \quad (1 \leq i \leq n)$$

where,  $s_0 = s_n = 0$  is assumed.

We continue to make the eigenvalue existence range smaller by repeatedly subdividing the interval while counting the number of eigenvalues as described above, based on  $x_{\min}$  and  $x_{\max}$ . In this way, both ends of the infinitesimal interval can be made to converge to a given eigenvalue.

For information about the Sturm sequence of functions and the Gerschgorin theorem, refer to entries (2) and (5) of the bibliography.

### 5.1.2.7 Accumulation of similarity (unitary) transformation by block algorithm

In seeking the eigenvectors of a real symmetric matrix using the QR method or the inverse iteration method, it is necessary to calculate the accumulation of the similarity (unitary) matrices that had already been computed in the preceding Householder transformation. It is very effective to apply the block algorithm to this accumulating procedure for getting better performance.

Let  $P_k$  be a transformation matrix that is obtained in Householder transformation which transform the real symmetric matrix to a tridiagonal matrix.

$$P_k = \mathbf{I} - \frac{\mathbf{u}_k \mathbf{u}_k^T}{H_k}$$

$$H_k = \frac{1}{2} \mathbf{u}_k^T \mathbf{u}_k$$

The accumulation of the transformation matrix  $P_k$  becomes to:

$$P_1 P_2 \cdots P_{n-2} = \mathbf{I} - \sum_{i=1}^{n-2} \mathbf{u}_i \mathbf{w}_i^T$$

where  $\mathbf{w}_i^T$  is expressed by the following recurrence formula.

$$\mathbf{w}_{n-2}^T = \frac{\mathbf{u}_{n-2}^T}{H_{n-2}}$$

$$\mathbf{w}_i^T = \frac{\mathbf{u}_i^T - \sum_{j=i-1}^{n-2} (\mathbf{u}_i^T \mathbf{u}_j) \mathbf{w}_j^T}{H_i}$$

If we let  $\mathbf{V}$  be the eigenvectors of a real symmetric tridiagonal matrix which are obtained with the QR method or the inverse iteration method. The eigenvectors  $\mathbf{X}$  of the original matrix is obtained by the following.

$$\begin{aligned} \mathbf{X} &= P_1 \cdots P_{n-2} \mathbf{V} \\ &= \mathbf{V} - \sum_{i=1}^{n-2} \mathbf{u}_i \mathbf{w}_i^T \mathbf{V} \end{aligned}$$

A product of a similarity (unitary) matrix  $P_k$  and the eigenvectors  $\mathbf{V}$  is a rank-one update. Therefore, the accumulation of the transformation matrices can be obtained at a higher speed if the matrix updates are performed in the lump. If the original matrix is Hermitian, the transpose notation “ $T$ ” should be replaced with the Hermite conjugate notation “ $*$ ”.

### 5.1.2.8 Inverse iteration method

The inverse iteration method is used to obtain the eigenvector corresponding to the eigenvalues obtained by the root-free QR method or bisection method.

Assume the approximate value  $\mu_k$  of a given eigenvalue  $\lambda_k$  of the real symmetric tridiagonal matrix  $T$  has been obtained. If a suitable initial vector  $\mathbf{v}_0$  has been chosen at this time and the linear simultaneous equations:

$$(T - \mu_k I)\mathbf{v}_i = \mathbf{v}_{i-1} \quad (i = 1, 2, \dots)$$

are iteratively solved, then if  $\mathbf{v}_i$  satisfy the convergence conditions, they converge to the eigenvector.

To solve the simultaneous linear equations, partial pivoting is performed while using the Gauss method to perform an LU decomposition. Then forward elimination and back substitution are performed.

### 5.1.2.9 Generalized eigenvalue problem

A Cholesky decomposition of  $B$  is performed:

$$B = LL^*$$

in the generalized eigenvalue problem for a Hermitian matrix:

$$A\mathbf{x} = \lambda B\mathbf{x} \quad (A : \text{Hermitian}, B : \text{Positive Hermitian})$$

yielding:

$$(L^{-1}A(L^*)^{-1})(L^*\mathbf{x}) = \lambda(L^*\mathbf{x})$$

If we set:

$$P = L^{-1}A(L^*)^{-1}$$

$$L^*\mathbf{x} = \mathbf{y}$$

then the generalized eigenvalue problem is transformed to a standard eigenvalue problem for the Hermitian matrix  $P$ .

$$P\mathbf{y} = \lambda\mathbf{y}$$

The eigenvector of matrix  $A$  is given by:

$$\mathbf{x} = (L^*)^{-1}\mathbf{y}.$$

Generalized eigenvalue problems for Hermitian matrix other than  $A\mathbf{x} = \lambda B\mathbf{x}$  ( $B$  : Positive Hermitian) are classified into two cases by the position of positive Hermitian Matrix  $B$  as:

$$AB\mathbf{x} = \lambda\mathbf{x}$$

and

$$BA\mathbf{x} = \lambda\mathbf{x}$$

The eigenvalues  $\lambda$  and the eigenvectors  $x$  of these equations can be obtained by reducing them to standard eigenvalue problems using the following procedure:

- ① Apply the Cholesky decomposition to the positive matrix  $B$  as  $B = L^*L$ . (Where  $L$  is a lower triangle matrix.)
- ②  $ABx = \lambda x$  is reduced to the eigenvalue problem for  $C = LAL^*$ , and the eigenvectors are obtained by multiplying the inverse of  $L$ .
- ③  $BAx = \lambda x$  is reduced to the eigenvalue problem for  $C = LAL^*$ , and the eigenvectors are obtained by multiplying  $L^*$ .

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## 5.2 REAL SYMMETRIC MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE)

### 5.2.1 ASL\_qcsmaa, ASL\_pcsmaa

#### All Eigenvalues and All Eigenvectors of a Real Symmetric Matrix

(1) **Function**

ASL\_qcsmaa or ASL\_pcsmaa uses the Householder method and QR method to obtain all eigenvalues of the real symmetric matrix  $A$  (two-dimensional array type) (upper triangular type) and all corresponding eigenvectors.

(2) **Usage**

Double precision:

`ierr = ASL_qcsmaa (a, lna, n, e, w1, nt);`

Single precision:

`ierr = ASL_pcsmaa (a, lna, n, e, w1, nt);`

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	$lna \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Eigenvectors (column vectors) corresponding to each eigenvalue
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues
5	w1	$\begin{cases} D* \\ R* \end{cases}$	n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]$ and $a[0] \leftarrow 1.0$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
5000+i	The sequence did not converge in the step where the eigenvalues obtained. ( $1 \leq i \leq n$ )	Eigenvalues correctly obtained by this time are entered in $e[0], \dots, e[i-2]$ and eigenvectors corresponding to them are entered in a. (However, the order is irregular.)

(6) **Notes**

- (a) Data should be stored only in the upper triangular portion of array a.
- (b) Eigenvalues are stored in ascending order.
- (c) The eigenvectors are an orthonormal system.
- (d) If eigenvectors are not required, use 5.2.2  $\left\{ \begin{array}{l} \text{ASL\_qcsman} \\ \text{ASL\_pcsman} \end{array} \right\}$ .

(7) **Example**

(a) Problem

Obtain all eigenvalues of the matrix:

$$A = \begin{bmatrix} 6 & 4 & 4 & 1 \\ 4 & 6 & 1 & 4 \\ 4 & 1 & 6 & 4 \\ 1 & 4 & 4 & 6 \end{bmatrix}$$

and their corresponding eigenvectors.

(b) Input data

Matrix A, lda=11, n=4 and nt=2.

(c) Main program

```

/*      C interface example for ASL_qcsmaa */

#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lda=11;
    int n;
    double *e;
    double *w1;
    int nt = 2;
    int ierr;
    int i,j,k;
    FILE *fp;

    int mod;

    fp = fopen( "qcsmaa.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }
}

```

```

}

printf( "    *** ASL_qcsmaa ***\n" );
printf( "\n    ** Input **\n\n" );

fscanf( fp, "%d", &n );

mod = n % 4;

a = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

e = ( double * )malloc((size_t)( sizeof(double) * n ));
if( e == NULL )
{
    printf( "no enough memory for array e\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * n ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\tn = %6d\n", n );
printf( "\tnt= %6d\n\n", nt );
printf( "\tInput Matrix a\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=i ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &a[i+lda*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g ", a[j+lda*i] );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "%8.3g ", a[i+lda*j] );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_qcsmaa(a, lda, n, e, w1, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
for( k=0 ; k<n-3 ; k = k+4 )
{
    printf( "\n" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "\tEigenvalue " );
    }
    printf( "\n" );
    for( i=k ; i<k+4 ; i++ )
    {
        printf( "\t%8.3g      ", e[i] );
    }
    printf( "\n" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "\tEigenvector " );
    }
    printf( "\n" );
    for( j=0 ; j<n ; j++ )
    {
        for( i=k ; i<k+4 ; i++ )
        {
            printf( "\t%8.3g      ", a[j+lda*i] );
        }
        printf( "\n" );
    }
}

```

```

    }
}
if( mod != 0 )
{
    printf( "\n" );
    for( i=n-mod ; i<n ; i++ )
    {
        printf( "\tEigenvalue " );
    }
    printf( "\n" );
    for( i=n-mod ; i<n ; i++ )
    {
        printf( "\t%8.3g      ", e[i] );
    }
    printf( "\n" );
    for( i=n-mod ; i<n ; i++ )
    {
        printf( "\tEigenvector " );
    }
    printf( "\n" );
    for( j=1 ; j<n ; j++ )
    {
        for( i=n-mod ; i<n ; i++ )
        {
            printf( "\t%8.3g      ", a[j+lda*i] );
        }
        printf( "\n" );
    }
}

free( a );
free( e );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_qcsmaa ***

** Input **

n =      4
nt=      2

Input Matrix a

      6      4      4      1
      4      6      1      4
      4      1      6      4
      1      4      4      6

** Output **

ierr =      0

Eigenvalue  Eigenvalue  Eigenvalue  Eigenvalue
      -1           5           5           15
Eigenvector  Eigenvector  Eigenvector  Eigenvector
      0.5          0.707          0          0.5
      -0.5         2.78e-17         -0.707         0.5
      -0.5         -1.67e-16          0.707         0.5
      0.5          -0.707         8.33e-17         0.5

```

## 5.2.2 ASL\_qcsman, ASL\_pcsman All Eigenvalues of a Real Symmetric Matrix

### (1) Function

ASL\_qcsman or ASL\_pcsman uses the Householder method and root-free QR method to obtain all eigenvalues of the real symmetric matrix  $A$  (two-dimensional array type) (upper triangular type).

### (2) Usage

Double precision:

```
ierr = ASL_qcsman (a, lna, n, e, w1, nt);
```

Single precision:

```
ierr = ASL_pcsman (a, lna, n, e, w1, nt);
```

### (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type)(upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	e	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Output	Eigenvalues
5	w1	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

### (4) Restrictions

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$5000+i$	The sequence did not converge in the step where the eigenvalues obtained. ( $1 \leq i \leq n$ )	Eigenvalues correctly obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.)

(6) **Notes**

- (a) Data should be stored only in the upper triangular portion of array a.
- (b) Eigenvalues are stored in ascending order.

### 5.2.3 ASL\_qcsmss, ASL\_pcsms Eigenvalues and Eigenvectors of a Real Symmetric Matrix

(1) **Function**

ASL\_qcsmss or ASL\_pcsms uses the Householder method, root free QR method, or Bisection method to obtain the m largest or m smallest eigenvalues of the real symmetric matrix *A* (two-dimensional array type) (upper triangular type) and the inverse iteration method to obtain the corresponding eigenvectors.

(2) **Usage**

Double precision:

ierr = ASL\_qcsmss (a, lna, n, eps, e, m, ve, lnv, isw, iw1, w1, nt);

Single precision:

ierr = ASL\_pcsms (a, lna, n, eps, e, m, ve, lnv, isw, iw1, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lna×n	Input	Real symmetric matrix <i>A</i> (two-dimensional array type)(upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix <i>A</i>
4	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalues convergence test. (See Note (d))
5	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
6	m	I	1	Input	The number m of eigenvalues to be obtained.
7	ve	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lnv×m	Output	Eigenvectors (column vector) corresponding to each eigenvalue.
8	lnv	I	1	Input	Adjustable dimension of array ve
9	isw	I	1	Input	Processing switch isw ≥ 0: Obtain m eigenvalues from the largest one. isw < 0: Obtain m eigenvalues from the smallest one.
10	iw1	I*	m	Output	Eigenvectors flag (See Note (e))
11	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	8×n	Work	Work area
12	nt	I	1	Input	Number of tasks to be generated
13	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \text{lna}$ ,  $\text{lnv}$
- (b)  $0 < m \leq n$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$e[0] \leftarrow a[0]$ and $ve[0] \leftarrow 1.0$ are performed.
2000	The maximum number of iterations was exceeded by the inverse iterations for obtaining eigenvectors.	Some eigenvectors are obtained with low precision, and processing continues. (See Note (e))
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) Data should be stored only in the upper triangular portion of array  $a$ .
- (b) If  $\text{isw} \geq 0$ , the eigenvalues are stored in descending order. If  $\text{isw} < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $\text{eps} \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically.  $\text{eps}$  is used to obtain eigenvalues by using the Bisection method.
- (e) If the maximum number of iterations is exceeded when using the inverse iteration method ( $\text{ierr} = 2000$  is output), the following processing is performed.  
 If  $\text{iw1}[i - 1] = 0$ : The  $i$ -th eigenvector calculation is normally terminated.  
 If  $\text{iw1}[i - 1] \neq 0$ : The convergence condition is not satisfied for the  $i$ -th eigenvector calculation, and the eigenvector precision is low. In this case, the iteration count is set for  $\text{iw1}[i - 1]$ .  
 If processing is normally terminated ( $\text{ierr} = 0$  is output),  $\text{iw1}[i - 1] = 0$  is set.
- (f) The eigenvectors are an orthonormal system.
- (g) If eigenvectors are not required, use 5.2.4  $\left\{ \begin{array}{l} \text{ASL\_qcsmn} \\ \text{ASL\_pcsmn} \end{array} \right\}$ .



(7) **Example**

(a) Problem

Obtain the three smallest eigenvalues of the matrix:

$$A = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

and their corresponding eigenvectors.

(b) Input data

Matrix  $A$ , lna=11, n=6, eps=-1.0, m=3, lmv=11, isw=-1 and nt=2.

(c) Main program

```

/*      C interface example for ASL_qcsmss */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int lda=11;
    int n;
    double ceps = -1.0;
    double *e;
    int m;
    double *ve;
    int ldv=11;
    int isw = -1;
    int *iw1;
    double *w1;
    int nt = 2;
    int ierr;
    int i,j,k;
    FILE *fp;

    int mod;

    fp = fopen( "qcsmss.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_qcsmss ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &n );
    fscanf( fp, "%d", &m );

    mod = m % 4;

    a = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    e = ( double * )malloc((size_t)( sizeof(double) * m ));
    if( e == NULL )
    {
        printf( "no enough memory for array e\n" );
        return -1;
    }

    ve = ( double * )malloc((size_t)( sizeof(double) * (ldv*m) ));
    if( ve == NULL )
    {
        printf( "no enough memory for array ve\n" );
        return -1;
    }

    iw1 = ( int * )malloc((size_t)( sizeof(int) * m ));
    if( iw1 == NULL )

```

```

{
    printf( "no enough memory for array iw1\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * (8*n) ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\tn = %6d\n", n );
printf( "\tm = %6d\n", m );
printf( "\tnt= %6d\n", nt );

printf( "\n\tInput Matrix a\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=i ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &a[i+lda*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g ", a[j+lda*i] );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "%8.3g ", a[i+lda*j] );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_qcsmss(a, lda, n, cepts, e, m, ve, ldv, isw, iw1, w1, nt);

printf( "\n    ** Output **\n\n" );
printf( "\ttierr = %6d\n", ierr );

for( k=0 ; k<m-3 ; k = k+4 )
{
    printf( "\n" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "\tEigenvalue  " );
    }
    printf( "\n" );
    for( i=k ; i<k+4 ; i++ )
    {
        printf( "\t%8.3g      ", e[i] );
    }
    printf( "\n" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "\tEigenvector " );
    }
    printf( "\n" );
    for( j=0 ; j<n ; j++ )
    {
        for( i=k ; i<k+4 ; i++ )
        {
            printf( "\t%8.3g      ", ve[j+ldv*i] );
        }
        printf( "\n" );
    }
}

if( mod != 0 )
{
    printf( "\n" );
    for( i= m-mod ; i<m ; i++ )
    {
        printf( "\tEigenvalue  " );
    }
    printf( "\n" );
    for( i= m-mod ; i<m ; i++ )
    {
        printf( "\t%8.3g      ", e[i] );
    }
}

```

```

printf( "\n" );
for( i= m-mod ; i<m ; i++ )
{
    printf( "\tEigenvector " );
}
printf( "\n" );
for( j=0 ; j<n ; j++ )
{
    for( i= m-mod ; i<m ; i++ )
    {
        printf( "\t%8.3g      ", ve[j+ldv*i] );
    }
    printf( "\n" );
}
}

free( a );
free( e );
free( ve );
free( iw1 );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_qcsmss ***

** Input **

n =      4
m =      3
nt=      2

Input Matrix a

      4      1      0      0
      1      3      1      0
      0      1      3      1
      0      0      1      4

** Output **

ierr =      0

Eigenvalue      Eigenvalue      Eigenvalue
  1.59           3                4.41
Eigenvalue      Eigenvalue      Eigenvalue
-0.271          0.5              -0.653
 0.653          -0.5              -0.271
-0.653          -0.5              0.271
 0.271          0.5                0.653

```

### 5.2.4 ASL\_qcsmsn, ASL\_pcsmsn Eigenvalues of a Real Symmetric Matrix

(1) **Function**

ASL\_qcsmsn or ASL\_pcsmsn uses the Householder method, root-free QR method, or Bisection method to obtain the  $m$  largest or  $m$  smallest eigenvalues of the real symmetric matrix  $A$  (two-dimensional array type) (upper triangular type).

(2) **Usage**

Double precision:

ierr = ASL\_qcsmsn (a, lna, n, eps, e, m, isw, w1, nt);

Single precision:

ierr = ASL\_pcsmsn (a, lna, n, eps, e, m, isw, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type)(upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalues convergence test. (See Note (d))
5	e	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	m	Output	Eigenvalues
6	m	I	1	Input	The number m of eigenvalues to be obtained.
7	isw	I	1	Input	Processing switch isw $\geq$ 0: Obtain m eigenvalues from the largest one. isw $<$ 0: Obtain m eigenvalues from the smallest one.
8	w1	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$5 \times n$	Work	Work area
9	nt	I	1	Input	Number of tasks to be generated
10	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \ln a$
- (b)  $0 < m \leq n$
- (c)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]$ is performed.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) Data should be stored only in the upper triangular portion of array a.
- (b) If  $isw \geq 0$ , the eigenvalues are stored in descending order. If  $isw < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $eps \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically. eps is used to obtain eigenvalues by using the Bisection method.

---

## 5.3 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE)

### 5.3.1 ASL\_hchraa, ASL\_gchraa

#### All Eigenvalues and All Eigenvectors of a Hermitian Matrix

(1) **Function**

ASL\_hchraa or ASL\_gchraa uses the Householder method and QR method to obtain all eigenvalues of the Hermitian matrix  $A=(ar, ai)$  (two-dimensional array type) (upper triangular type) (real argument type) and all corresponding eigenvectors.

(2) **Usage**

Double precision:

```
ierr = ASL_hchraa (ar, ai, lna, n, e, vr, vi, lnv, w1, nt);
```

Single precision:

```
ierr = ASL_gchraa (ar, ai, lna, n, e, vr, vi, lnv, w1, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lna \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lna \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues
6	vr	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lnv \times n$	Output	Real part (column vector) of eigenvectors corresponding to each eigenvalue
7	vi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lnv \times n$	Output	Imaginary part (column vectors) of eigenvectors corresponding to each eigenvalue
8	lnv	I	1	Input	Adjustable dimension of arrays vr and vi
9	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$3 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $0 < n \leq lna, lnv$
- (b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]$ , $vr[0] \leftarrow 1.0$ and $vi[0] \leftarrow 0.0$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$5000+i$	The sequence did not converge in the step where the eigenvalue is obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.) Not eigenvector is obtained at this time.

(6) **Notes**

- (a) Real and imaginary parts of the Hermitian matrix are stored only in the upper triangular portions of arrays ar and ai respectively. (See Appendix A)
- (b) Eigenvalues are stored in ascending order.
- (c) The eigenvectors are an orthonormal set.
- (d) If eigenvectors are not required, use 5.3.2  $\left\{ \begin{array}{l} \text{ASL_hchran} \\ \text{ASL_gchran} \end{array} \right\}$ .

(7) **Example**

- (a) Problem

Obtain all eigenvalues of the matrix:

$$A = \begin{bmatrix} 7 & 3 & 1+2i & -1+2i \\ 3 & 7 & 1-2i & -1-2i \\ 1-2i & 1+2i & 7 & -3 \\ -1-2i & -1+2i & -3 & 7 \end{bmatrix}$$

and their corresponding eigenvectors.

- (b) Input data

Real part ar and imaginary part ai of matrix A, lna=11, n=4, lnv=10 and nt=2.

- (c) Main program

```

/*      C interface example for ASL_hchraa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lda=11;
    int n;
    double *e;
    double *vr;
    double *vi;
    int ldv=11;
    double *w1;
    int nt;
    int ierr;
    int i,j;
    FILE *fp;

    int mod;

```



```

fp = fopen( "hchraa.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_hchraa ***\n" );
printf( "\n    ** Input **\n\n" );

fscanf( fp, "%d", &n );
fscanf( fp, "%d", &nt );

mod = n % 2;

ar = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
if( ar == NULL )
{
    printf( "no enough memory for array ar\n" );
    return -1;
}

ai = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
if( ai == NULL )
{
    printf( "no enough memory for array ai\n" );
    return -1;
}

e = ( double * )malloc((size_t)( sizeof(double) * n ));
if( e == NULL )
{
    printf( "no enough memory for array e\n" );
    return -1;
}

vr = ( double * )malloc((size_t)( sizeof(double) * (ldv*n) ));
if( vr == NULL )
{
    printf( "no enough memory for array vr\n" );
    return -1;
}

vi = ( double * )malloc((size_t)( sizeof(double) * (ldv*n) ));
if( vi == NULL )
{
    printf( "no enough memory for array vi\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * (3*n) ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\tn = %6d\n", n );
printf( "\tnt= %6d\n\n", nt );

printf( "\n\tInput Matrix a ( Real , Imaginary )\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=i ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ar[i+lda*j] );
        fscanf( fp, "%lf", &ai[i+lda*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", ar[j+lda*i], -ai[j+lda*i] );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", ar[i+lda*j], ai[i+lda*j] );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_hchraa(ar, ai, lda, n, e, vr, vi, ldv, w1, nt);

printf( "\n    ** Output **\n\n" );

```

```

printf( "\tierr = %6d\n", ierr );
for( j=0 ; j<n-1 ; j = j+2 )
{
    printf( "\n" );
    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvalue          " );
    }
    printf( "\n" );
    printf( "\t\t%8.3g          \t%8.3g\n",
            e[j], e[j+1] );

    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvector          " );
    }
    printf( "\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t\t%8.3g , %8.3g          \t%8.3g , %8.3g\n",
                vr[i+ldv*j], vi[i+ldv*j], vr[i+ldv*(j+1)],vi[i+ldv*(j+1)] );
    }
}

if( mod != 0 )
{
    printf( "\n" );
    printf( "\tEigenvalue\n" );
    printf( "\t\t%8.3g\n", e[n-1] );
    printf( "\tEigenvector\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t\t%8.3g , %8.3g\n",
                vr[i+ldv*(n-1)], vi[i+ldv*(n-1)] );
    }
}

free( ar );
free( ai );
free( e );
free( vr );
free( vi );
free( wl );

return 0;
}

```

(d) Output results

```

*** ASL_hchraa ***

** Input **

n =      4
nt=      2

Input Matrix a ( Real , Imaginary )

(      7 ,      0) (      3 ,      0) (      1 ,      2) (      -1 ,      2)
(      3 ,      0) (      7 ,      0) (      1 ,     -2) (      -1 ,     -2)
(      1 ,     -2) (      1 ,      2) (      7 ,      0) (      -3 ,      0)
(     -1 ,     -2) (     -1 ,      2) (     -3 ,      0) (      7 ,      0)

** Output **

ierr =      0

Eigenvalue      0      Eigenvalue      8
Eigenvector
0.5 ,      0      Eigenvector
-0.707 ,      0
-0.5 ,      0      6.38e-16 ,      0
1.67e-16 ,      0.5      0.354 ,      0.354
-1.11e-16 ,      0.5      -0.354 ,      0.354

Eigenvalue      8      Eigenvalue     12
Eigenvector
0 ,      0      Eigenvector
0.5 ,      0
-0.0999 ,      0.7      0.5 ,      0
-0.3 ,     -0.4      0.5 , -5.55e-16
-0.4 ,      0.3      -0.5 , -4.44e-16

```

### 5.3.2 ASL\_hchran, ASL\_gchran All Eigenvalues of a Hermitian Matrix

(1) **Function**

ASL\_hchran or ASL\_gchran uses the Householder method and root-free QR method to obtain all eigenvalues of the Hermitian matrix  $A=(ar, ai)$  (two-dimensional array type) (upper triangular type) (real argument type).

(2) **Usage**

Double precision:

ierr = ASL\_hchran (ar, ai, lna, n, e, w1, nt);

Single precision:

ierr = ASL\_gchran (ar, ai, lna, n, e, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	ai	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	e	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Output	Eigenvalues
6	w1	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$3 \times n$	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow ar[0]$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$5000+i$	The sequence did not converge in the step where the eigenvalue is obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.)

(6) **Notes**

- (a) Real and imaginary parts of the Hermitian matrix are stored only in the upper triangular portions of arrays ar and ai respectively. (See Appendix A)
- (b) Eigenvalues are stored in ascending order.

### 5.3.3 ASL\_hchrss, ASL\_gchrss Eigenvalues and Eigenvectors of a Hermitian Matrix

(1) **Function**

ASL\_hchrss or ASL\_gchrss uses the Householder method, root-free QR method, or Bisection method to obtain the  $m$  largest or  $m$  smallest eigenvalues of the Hermitian matrix  $A=(ar, ai)$  (two-dimensional array type) (upper triangular type) (real argument type) and the inverse iteration method to obtain the corresponding eigenvectors.

(2) **Usage**

Double precision:

```
ierr = ASL_hchrss (ar, ai, lna, n, eps, e, m, vr, vi, lnv, isw, iw1, w1, nt);
```

Single precision:

```
ierr = ASL_gchrss (ar, ai, lna, n, eps, e, m, vr, vi, lnv, isw, iw1, w1, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
6	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
7	m	I	1	Input	The number of m of eigenvalues to be obtained.
8	vr	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnv} \times m$	Output	Real part (column vector) of eigenvectors corresponding to each eigenvalue
9	vi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnv} \times m$	Output	Imaginary parts (column vectors) of eigenvectors corresponding to each eigenvalue
10	lnv	I	1	Input	Adjustable dimension of arrays vr and vi
11	isw	I	1	Input	Processing switch isw $\geq$ 0: Obtain m eigenvalues from the largest one. isw $<$ 0: Obtain m eigenvalues from the smallest one.
12	iw1	I*	m	Output	Eigenvector flag (See Note (e))
13	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$10 \times n$	Work	Work area
14	nt	I	1	Input	Number of tasks to be generated
15	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \text{lna}$ ,  $\text{lnv}$
- (b)  $0 < m \leq n$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$e[0] \leftarrow ar[0]$ , $vr[0] \leftarrow 1.0$ and $vi[0] \leftarrow 0.0$ are performed.
2000	The maximum number of iterations was exceeded by the inverse iterations for obtaining eigenvectors.	Some eigenvectors are obtained with low precision, and processing continues. (See Note (e))
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) The real and imaginary parts of the Hermitian matrix should be stored only in the upper triangular portions of arrays  $ar$  and  $ai$  respectively (See Appendix A).
- (b) If  $\text{isw} \geq 0$ , the eigenvalues are stored in descending order. If  $\text{isw} < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $\text{eps} \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically.  $\text{eps}$  is used to obtain eigenvalues by using the Bisection method.
- (e) If the maximum number of iterations is exceeded when using the inverse iteration method ( $\text{ierr} = 2000$  is output), the following processing is performed.  
 If  $\text{iw1}[i - 1] = 0$ : The  $i$ -th eigenvector calculation is normally terminated.  
 If  $\text{iw1}[i - 1] \neq 0$ : The convergence condition is not satisfied for the  $i$ -th eigenvector calculation, and the eigenvector precision is low. In this case, the iteration count is set for  $\text{iw1}[i - 1]$ .  
 If processing is normally terminated ( $\text{ierr} = 0$  is output),  $\text{iw1}[i - 1] = 0$  is set.
- (f) The eigenvectors are an orthonormal set.
- (g) If eigenvectors are not required, use 5.3.4  $\left\{ \begin{array}{l} \text{ASL\_hchrns} \\ \text{ASL\_gchrns} \end{array} \right\}$ .

(7) **Example**

(a) Problem

Obtain the three largest eigenvalues of the Hermitian matrix  $A$ .

$$A = \begin{bmatrix} 7 & 3 & 1 + 2i & -1 + 2i \\ 3 & 7 & 1 - 2i & -1 - 2i \\ 1 - 2i & 1 + 2i & 7 & -3 \\ -1 - 2i & -1 + 2i & -3 & 7 \end{bmatrix}$$

and their corresponding eigenvectors.

(b) Input data

Real part ar and imaginary part ai of matrix  $A$ , lna=11, n=4, eps=-1.0, m=3, lnv=11, isw=1 and nt=2.

(c) Main program

```

/*      C interface example for ASL_hchrss */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar;
    double *ai;
    int lda=11;
    int n;
    double ceps= -1.0;
    double *e;
    int m;
    double *vr;
    double *vi;
    int ldv=11;
    int isw=1;
    int *iw1;
    double *w1;
    int nt;
    int ierr;
    int i,j;
    FILE *fp;

    int mod;

    fp = fopen( "hchrss.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hchrss ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &n );
    fscanf( fp, "%d", &m );
    fscanf( fp, "%d", &nt );

    mod = m % 2;

    ar = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
    if( ar == NULL )
    {
        printf( "no enough memory for array ar\n" );
        return -1;
    }

    ai = ( double * )malloc((size_t)( sizeof(double) * (lda*n) ));
    if( ai == NULL )
    {
        printf( "no enough memory for array ai\n" );
        return -1;
    }

    e = ( double * )malloc((size_t)( sizeof(double) * m ));
    if( e == NULL )
    {
        printf( "no enough memory for array e\n" );
        return -1;
    }

    vr = ( double * )malloc((size_t)( sizeof(double) * (ldv*m) ));

```



```

if( vr == NULL )
{
    printf( "no enough memory for array vr\n" );
    return -1;
}

vi = ( double * )malloc((size_t)( sizeof(double) * (ldv*m) ));
if( vi == NULL )
{
    printf( "no enough memory for array vi\n" );
    return -1;
}

iw1 = ( int * )malloc((size_t)( sizeof(int) * m ));
if( iw1 == NULL )
{
    printf( "no enough memory for array iw1\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * (10*n) ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

printf( "\tn = %6d\n", n );
printf( "\tm = %6d\n", m );
printf( "\tnt= %6d\n\n", nt );

printf( "\n\tInput Matrix a ( Real , Imaginary )\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=i ; j<n ; j++ )
    {
        fscanf( fp, "%lf", &ar[i+lda*j] );
        fscanf( fp, "%lf", &ai[i+lda*j] );
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g , %8.3g ", ar[j+lda*i], -ai[j+lda*i] );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "%8.3g , %8.3g ", ar[i+lda*j], ai[i+lda*j] );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_hchrss(ar, ai, lda, n, ceps, e, m, vr, vi, ldv, isw, iw1, w1, nt);

printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

for( j=0 ; j<m-1 ; j = j+2 )
{
    printf( "\n" );
    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvalue          " );
    }
    printf( "\n" );
    printf( "\t%8.3g          \t%8.3g\n",
        e[j], e[j+1] );

    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvector          " );
    }
    printf( "\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t%8.3g , %8.3g          \t%8.3g , %8.3g\n",
            vr[i+ldv*j], vi[i+ldv*j], vr[i+ldv*(j+1)], vi[i+ldv*(j+1)] );
    }
}

if( mod != 0 )
{
    printf( "\n" );
}

```

```

printf( "\tEigenvalue\n" );
printf( "\t%8.3g\n", e[m-1] );
printf( "\tEigenvector\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t%8.3g , %8.3g\n",
           vr[i+ldv*(m-1)], vi[i+ldv*(m-1)] );
}
}

free( ar );
free( ai );
free( e );
free( vr );
free( vi );
free( iw1 );
free( w1 );

return 0;
}

```

(d) Output results

```

*** ASL_hchrss ***

** Input **

n =      4
m =      3
nt=      2

Input Matrix a ( Real , Imaginary )

(      7 ,      0) (      3 ,      0) (      1 ,      2) (      -1 ,      2)
(      3 ,      0) (      7 ,      0) (      1 ,     -2) (      -1 ,     -2)
(      1 ,     -2) (      1 ,      2) (      7 ,      0) (      -3 ,      0)
(     -1 ,     -2) (     -1 ,      2) (     -3 ,      0) (      7 ,      0)

** Output **

ierr =      0

Eigenvalue      12      Eigenvalue      8
Eigenvector      Eigenvector
      0.5 ,      0      -0.0999 ,      0
      0.5 , 1.02e-16      -0.3 ,     -0.7
      0.5 , -5.41e-16      -0.4 ,     -0.4
     -0.5 , -4.58e-16      -0.4 ,      0.3

Eigenvalue      8
Eigenvector
      0.707 ,      0
     -6.66e-16 , -5.09e-17
     -0.354 ,     -0.354
      0.354 ,     -0.354

```

### 5.3.4 ASL\_hchrsn, ASL\_gchrsn Eigenvalues of a Hermitian Matrix

(1) **Function**

ASL\_hchrsn or ASL\_gchrsn uses the Householder method, root-free QR method, or Bisection method to obtain the  $m$  largest or  $m$  smallest eigenvalues of the Hermitian matrix  $A=(ar, ai)$  (two-dimensional array type) (upper triangular type) (real argument type).

(2) **Usage**

Double precision:

```
ierr = ASL_hchrsn (ar, ai, lna, n, eps, e, m, isw, w1, nt);
```

Single precision:

```
ierr = ASL_gchrsn (ar, ai, lna, n, eps, e, m, isw, w1, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lna \times n$	Input	Real part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$lna \times n$	Input	Imaginary part of Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrix $A$
5	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
6	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
7	m	I	1	Input	The number of m of eigenvalues to be obtained.
8	isw	I	1	Input	Processing switch isw $\geq$ 0: Obtain m eigenvalues from the largest one. isw $<$ 0: Obtain m eigenvalues from the smallest one.
9	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$5 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $0 < n \leq lna$
- (b)  $0 < m \leq n$
- (c)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow ar[0]$ , is performed.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) The real and imaginary parts of the Hermitian matrix should be stored only in the upper triangular portions of arrays ar and ai respectively (See Appendix A).
- (b) If  $isw \geq 0$ , the eigenvalues are stored in descending order. If  $isw < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $eps \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically. eps is used to obtain eigenvalues by using the Bisection method.

## 5.4 HERMITIAN MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (COMPLEX ARGUMENT TYPE)

### 5.4.1 ASL\_hcheaa, ASL\_gcheaa

#### All Eigenvalues and All Eigenvectors of a Hermitian Matrix

(1) **Function**

ASL\_hcheaa or ASL\_gcheaa uses the Householder method or QR method to obtain all eigenvalues of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) (complex argument type) and all corresponding eigenvectors.

(2) **Usage**

Double precision:

ierr = ASL\_hcheaa (a, lna, n, e, w1, w2, nt);

Single precision:

ierr = ASL\_gcheaa (a, lna, n, e, w1, w2, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\left\{ \begin{array}{l} Z^* \\ C^* \end{array} \right\}$	$lna \times n$	Input	Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Eigenvectors (column vector) corresponding to each eigenvalue
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	e	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	n	Output	Eigenvalues
5	w1	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	n	Work	Work area
6	w2	$\left\{ \begin{array}{l} Z^* \\ C^* \end{array} \right\}$	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $0 < n \leq \ln a$
- (b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \text{creal}(a[0])$ and $a[0] \leftarrow (1.0, 0.0)$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$5000+i$	The sequence did not converge in the step where the eigenvalue is obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ (However, the order is irregular). No eigenvector is obtained at this time.

(6) **Notes**

- (a) Only the upper triangular portion of the Hermitian matrix should be stored in array a. (See Appendix A)
- (b) Eigenvalues are stored in ascending order.
- (c) The eigenvectors are an orthonormal set.
- (d) If eigenvectors are not required, use 5.4.2  $\left\{ \begin{matrix} \text{ASL_hchean} \\ \text{ASL_gchean} \end{matrix} \right\}$ .

(7) **Example**

(a) Problem

Obtain all eigenvalues of the matrix:

$$A = \begin{bmatrix} 7 & 3 & 1+2i & -1+2i \\ 3 & 7 & 1-2i & -1-2i \\ 1-2i & 1+2i & 7 & -3 \\ -1-2i & -1+2i & -3 & 7 \end{bmatrix}$$

and their corresponding eigenvectors.

(b) Input data

Matrix A,  $\ln a=11$  and  $n=4$ .

(c) Main program

```

/*      C interface example for ASL_hcheaa */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lda=11;
    int n;
    double *e;
    double *w1;

```

```

double _Complex *w2;
int nt = 2;
int ierr;
int i,j;
FILE *fp;

int mod;

fp = fopen( "hcheaa.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_hcheaa ***\n" );
printf( "\n    ** Input **\n\n" );
fscanf( fp, "%d", &n );
mod = n % 2;

a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lda*n) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

e = ( double * )malloc((size_t)( sizeof(double) * n ));
if( e == NULL )
{
    printf( "no enough memory for array e\n" );
    return -1;
}

w1 = ( double * )malloc((size_t)( sizeof(double) * n ));
if( w1 == NULL )
{
    printf( "no enough memory for array w1\n" );
    return -1;
}

w2 = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * n ));
if( w2 == NULL )
{
    printf( "no enough memory for array w2\n" );
    return -1;
}

printf( "\tn = %6d\n", n );
printf( "\n\tInput Matrix a ( Real , Imaginary )\n\n" );
for( i=0 ; i<n ; i++ )
{
    for( j=i ; j<n ; j++ )
    {
        double tmp_re, tmp_im;
        fscanf( fp, "%lf", &tmp_re );
        fscanf( fp, "%lf", &tmp_im );
        a[i+lda*j] = tmp_re + tmp_im * _Complex_I;
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", creal(a[j+lda*i]), -cimag(a[j+lda*i]) );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "(%8.3g , %8.3g) ", creal(a[i+lda*j]), cimag(a[i+lda*j]) );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_hcheaa(a, lda, n, e, w1, w2, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
for( j=0 ; j<n-1 ; j = j+2 )
{

```



```

printf( "\n" );
for( i=0 ; i<2 ; i++ )
{
    printf( "\tEigenvalue          " );
}
printf( "\n" );
printf( "\t\t%8.3g          \t%8.3g\n",
        e[j], e[j+1] );

for( i=0 ; i<2 ; i++ )
{
    printf( "\tEigenvector          " );
}
printf( "\n" );
for( i=0 ; i<n ; i++ )
{
    printf( "\t\t%8.3g , %8.3g          \t%8.3g , %8.3g\n",
            creal(a[i+lda*j]), cimag(a[i+lda*j]), creal(a[i+lda*(j+1)]), cimag(a[i+lda*(j+1)]) );
}
}

if( mod != 0 )
{
    printf( "\n" );
    printf( "\tEigenvalue\n" );
    printf( "\t\t%8.3g\n", e[n-1] );
    printf( "\tEigenvector\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t\t%8.3g , %8.3g\n",
                creal(a[i+lda*(n-1)]), cimag(a[i+lda*(n-1)]) );
    }
}

free( a );
free( e );
free( w1 );
free( w2 );

return 0;
}

```

(d) Output results

```

*** ASL_hcheaa ***

** Input **

n =      4

Input Matrix a ( Real , Imaginary )

(      7 ,      0) (      3 ,      0) (      1 ,      2) (      -1 ,      2)
(      3 ,      0) (      7 ,      0) (      1 ,     -2) (      -1 ,     -2)
(      1 ,     -2) (      1 ,      2) (      7 ,      0) (      -3 ,      0)
(     -1 ,     -2) (     -1 ,      2) (     -3 ,      0) (      7 ,      0)

** Output **

ierr =      0

Eigenvalue      Eigenvalue
  0              8
Eigenvector
  0.5 ,          -0.707 ,
 -0.5 ,          6.11e-16 ,
2.22e-16 ,      0.5 ,      0.354 ,
-1.11e-16 ,    0.5 ,      -0.354 ,      0.354

Eigenvalue      Eigenvalue
  8             12
Eigenvector
  0 ,            0.5 ,
-0.0999 ,       0.5 ,
 -0.3 ,        -0.4 ,      0.5 ,
 -0.4 ,         0.3 ,     -0.5 ,

```

### 5.4.2 ASL\_hchean, ASL\_gchean All Eigenvalues of a Hermitian Matrix

(1) **Function**

ASL\_hchean or ASL\_gchean uses the Householder method or root-free QR method to obtain all eigenvalues of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) (complex argument type).

(2) **Usage**

Double precision:

ierr = ASL\_hchean (a, lna, n, e, w1, w2, nt);

Single precision:

ierr = ASL\_gchean (a, lna, n, e, w1, w2, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$lna \times n$	Input	Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$
4	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues
5	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Work	Work area
6	w2	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	n	Work	Work area
7	nt	I	1	Input	Number of tasks to be generated
8	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna$

(b)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \text{creal}(a[0])$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
$5000+i$	The sequence did not converge in the step where the eigenvalue is obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.)

(6) **Notes**

- (a) Only the upper triangular portion of the Hermitian matrix should be stored in array a. (See Appendix A)
- (b) Eigenvalues are stored in ascending order.

### 5.4.3 ASL\_hchess, ASL\_gchess Eigenvalues and Eigenvectors of a Hermitian Matrix

(1) **Function**

ASL\_hchess or ASL\_gchess uses the Householder method, root-free QR method, or Bisection method to obtain the  $m$  largest or  $m$  smallest eigenvalues of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) (complex argument type) and the inverse iterative method to obtain the corresponding eigenvectors.

(2) **Usage**

Double precision:

```
ierr = ASL_hchess (a, lna, n, eps, e, m, ve, lnv, isw, iw1, w1, w2, nt);
```

Single precision:

```
ierr = ASL_gchess (a, lna, n, eps, e, m, ve, lnv, isw, iw1, w1, w2, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a.
3	n	I	1	Input	Order of matrix $A$
4	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
5	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
6	m	I	1	Input	The number of m of eigenvalues to be obtained.
7	ve	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$\text{lnv} \times m$	Output	Eigenvectors (column vector) corresponding to each eigenvalue
8	lnv	I	1	Input	Adjustable dimension of array ve
9	isw	I	1	Input	Processing switch $\text{isw} \geq 0$ : Obtain m eigenvalues from the largest one. $\text{isw} < 0$ : Obtain m eigenvalues from the smallest one.
10	iw1	I*	m	Output	Eigenvector flag (See Note (e))
11	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$8 \times n$	Work	Work area
12	w2	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	n	Work	Work area
13	nt	I	1	Input	Number of tasks to be generated
14	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $0 < n \leq \text{lna}$ ,  $\text{lnv}$
- (b)  $0 < m \leq n$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \text{creal}(a[0])$ and $ve[0] \leftarrow (1.0, 0.0)$ are performed.
2000	The maximum number of iterations was exceeded by the inverse iterations for obtaining eigenvectors.	Some eigenvectors are obtained with low precision, and processing continues. (See Note (e))
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) Only the upper triangular portion of the Hermitian matrix should be stored in array  $A$ . (See Appendix A)
- (b) If  $isw \geq 0$ , the eigenvalues are stored in descending order. If  $isw < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $eps \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically.  $eps$  is used to obtain eigenvalues by using the Bisection method.
- (e) If the maximum number of iterations is exceeded when using the inverse iteration method ( $ierr = 2000$  is output), the following processing is performed.  
 If  $iw1[i - 1] = 0$ : The  $i$ -th eigenvector calculation is normally terminated.  
 If  $iw1[i - 1] \neq 0$ : The convergence condition is not satisfied for the  $i$ -th eigenvector calculation, and the eigenvector precision is low. In this case, the iteration count is set for  $iw1[i - 1]$ .  
 If processing is normally terminated ( $ierr = 0$  is output),  $iw1[i - 1] = 0$  is set.
- (f) The eigenvectors are an orthonormal set.
- (g) If eigenvectors are not required, use 5.4.4  $\left\{ \begin{array}{l} \text{ASL\_hchesn} \\ \text{ASL\_gchesn} \end{array} \right\}$ .

(7) **Example**

- (a) Problem

Obtain the three largest eigenvalues of the following Hermitian matrix  $A$ :

$$A = \begin{bmatrix} 7 & 3 & 1 + 2i & -1 + 2i \\ 3 & 7 & 1 - 2i & -1 - 2i \\ 1 - 2i & 1 + 2i & 7 & -3 \\ -1 - 2i & -1 + 2i & -3 & 7 \end{bmatrix}$$

and their corresponding eigenvectors.

- (b) Input data

Matrix  $A$ ,  $lma=11$ ,  $n=4$ ,  $eps=-1.0$ ,  $m=3$ ,  $lnv=11$  and  $isw=1$ .

(c) Main program

```

/*      C interface example for ASL_hchess */

#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    double _Complex *a;
    int lda=11;
    int n;
    double cepts= -1.0;
    double *e;
    int m;
    double _Complex *ve;
    int ldv=11;
    int isw=1;
    int *iw1;
    double *w1;
    double _Complex *w2;
    int nt=2;
    int ierr;
    int i,j;
    FILE *fp;

    int mod;

    fp = fopen( "hchess.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hchess ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &n );
    fscanf( fp, "%d", &m );

    mod = m % 2;

    a = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lda*n) ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    e = ( double * )malloc((size_t)( sizeof(double) * m ));
    if( e == NULL )
    {
        printf( "no enough memory for array e\n" );
        return -1;
    }

    ve = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (ldv*m) ));
    if( ve == NULL )
    {
        printf( "no enough memory for array ve\n" );
        return -1;
    }

    iw1 = ( int * )malloc((size_t)( sizeof(int) * m ));
    if( iw1 == NULL )
    {
        printf( "no enough memory for array iw1\n" );
        return -1;
    }

    w1 = ( double * )malloc((size_t)( sizeof(double) * (8*n) ));
    if( w1 == NULL )
    {
        printf( "no enough memory for array w1\n" );
        return -1;
    }

    w2 = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * n ));
    if( w2 == NULL )
    {
        printf( "no enough memory for array w2\n" );
        return -1;
    }

    printf( "\tn = %6d\n", n );
    printf( "\tm = %6d\n", m );

    printf( "\n\nInput Matrix a ( Real , Imaginary )\n\n" );
    for( i=0 ; i<n ; i++ )

```

```

{
    for( j=i ; j<n ; j++ )
    {
        double tmp_re, tmp_im;
        fscanf( fp, "%lf", &tmp_re );
        fscanf( fp, "%lf", &tmp_im );
        a[i+lda*j] = tmp_re + tmp_im * _Complex_I;
    }
}

for( i=0 ; i<n ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g , %8.3g ", creal(a[j+lda*i]), -cimag(a[j+lda*i]) );
    }
    for( j=i ; j<n ; j++ )
    {
        printf( "%8.3g , %8.3g ", creal(a[i+lda*j]), cimag(a[i+lda*j]) );
    }
    printf( "\n" );
}

fclose( fp );

ierr = ASL_hchess(a, lda, n, ceps, e, m, ve, ldv, isw, iw1, w1, w2, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

for( j=0 ; j<m-1 ; j = j+2 )
{
    printf( "\n" );
    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvalue          " );
    }
    printf( "\n" );
    printf( "\t%8.3g          \t%8.3g\n",
           e[j], e[j+1] );

    for( i=0 ; i<2 ; i++ )
    {
        printf( "\tEigenvector          " );
    }
    printf( "\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t%8.3g , %8.3g          \t%8.3g , %8.3g\n",
               creal(ve[i+ldv*j]), cimag(ve[i+ldv*j]), creal(ve[i+ldv*(j+1)]), cimag(ve[i+ldv*(j+1)]) );
    }
}

if( mod != 0 )
{
    printf( "\n" );
    printf( "\tEigenvalue\n" );
    printf( "\t%8.3g\n", e[m-1] );
    printf( "\tEigenvector\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t%8.3g , %8.3g\n",
               creal(ve[i+ldv*(m-1)]), cimag(ve[i+ldv*(m-1)]) );
    }
}

free( a );
free( e );
free( ve );
free( iw1 );
free( w1 );
free( w2 );

return 0;
}

```

(d) Output results

```

*** ASL_hchess ***

** Input **

n =      4
m =      3

```



```

Input Matrix a ( Real , Imaginary )
(      7 ,      0) (      3 ,      0) (      1 ,      2) (      -1 ,      2)
(      3 ,      0) (      7 ,      0) (      1 ,     -2) (      -1 ,     -2)
(      1 ,     -2) (      1 ,      2) (      7 ,      0) (      -3 ,      0)
(     -1 ,     -2) (     -1 ,      2) (     -3 ,      0) (      7 ,      0)

** Output **
ierr =      0

Eigenvalue      Eigenvalue
   12            8
Eigenvector      Eigenvector
  0.5 ,      0      0 ,      0
  0.5 , 1.02e-16  -0.0999 , 0.7
  0.5 ,  -5e-16   -0.3 ,   -0.4
 -0.5 , -4.44e-16 -0.4 ,    0.3

Eigenvalue      Eigenvector
      8
  0.707 ,      0
 -6.66e-16 , 0
 -0.354 ,   -0.354
  0.354 ,   -0.354
    
```

#### 5.4.4 ASL\_hchesn, ASL\_gchesn Eigenvalues of a Hermitian Matrix

(1) **Function**

ASL\_hchesn or ASL\_gchesn uses the Householder method, root-free QR method, or Bisection method to obtain the  $m$  largest or  $m$  smallest eigenvalues of the Hermitian matrix  $A$  (two-dimensional array type) (upper triangular type) (complex argument type).

(2) **Usage**

Double precision:

```
ierr = ASL_hchesn (a, lna, n, eps, e, m, isw, w1, w2, nt);
```

Single precision:

```
ierr = ASL_gchesn (a, lna, n, eps, e, m, isw, w1, w2, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	$\text{lna} \times \text{n}$	Input	Hermitian matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a.
3	n	I	1	Input	Order of matrix $A$
4	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
5	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
6	m	I	1	Input	The number of m of eigenvalues to be obtained.
7	isw	I	1	Input	Processing switch $\text{isw} \geq 0$ : Obtain m eigenvalues from the largest one. $\text{isw} < 0$ : Obtain m eigenvalues from the smallest one.
8	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times \text{n}$	Work	Work area
9	w2	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	n	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $0 < \text{n} \leq \text{lna}$
- (b)  $0 < \text{m} \leq \text{n}$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \text{creal}(a[0])$ is performed.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	

(6) **Notes**

- (a) Only the upper triangular portion of the Hermitian matrix should be stored in array a. (See Appendix A)
- (b) If  $\text{isw} \geq 0$ , the eigenvalues are stored in descending order. If  $\text{isw} < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $\text{eps} \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically.  $\text{eps}$  is used to obtain eigenvalues by using the Bisection method.

---

## 5.5 GENERALIZED EIGENVALUE PROBLEM FOR A REAL SYMMETRIC MATRIX (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) ( $Ax = \lambda Bx$ )

### 5.5.1 ASL\_qcgsaa, ASL\_pcgsaa

All Eigenvalues and All Eigenvectors of a Real Symmetric Matrix (Generalized Eigenvalue Problem  $Ax = \lambda Bx$ ,  $B$ : Positive)

(1) **Function**

ASL\_qcgsaa or ASL\_pcgsaa uses the Cholesky method to transform the real symmetric matrix (two-dimensional array type) (upper triangular type) generalized eigenvalue problem  $Ax = \lambda Bx$  ( $A$ : Real symmetric matrix,  $B$ : Positive real symmetric matrix) to a standard eigenvalue problem and uses the Householder method and QR method to obtain all eigenvalues and corresponding all eigenvectors.

(2) **Usage**

Double precision:

```
ierr = ASL_qcgsaa (a, lna, n, b, lnb, e, w1, nt);
```

Single precision:

```
ierr = ASL_pcgsaa (a, lna, n, b, lnb, e, w1, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	$\text{lna} \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Eigenvectors (column vector) corresponding to each eigenvalue
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$ and $B$
4	b	$\begin{cases} D* \\ R* \end{cases}$	$\text{lnb} \times n$	Input	Positive symmetric matrix $B$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues
7	w1	$\begin{cases} D* \\ R* \end{cases}$	$2 \times n$	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < n \leq \text{lna}, \text{lnb}$   
(b)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]/b[0]$ and $a[0] \leftarrow 1.0/\sqrt{b[0]}$ are performed.
2100	$B$ has a diagonal element very close to zero.	Some eigenvectors may be obtained with low precision, and processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000+i	The sequence did not converge in the step where the eigenvalue is obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.) No eigenvector is obtained at this time.

## (6) Notes

- (a) Data should be stored only in the upper triangular portions of arrays a and b.
- (b) Eigenvalues are stored in ascending order.
- (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^T B v_k = \delta_{j,k}$
- (d) If eigenvectors are not required, use 5.5.2  $\left\{ \begin{array}{l} \text{ASL\_qcgsan} \\ \text{ASL\_pcgsan} \end{array} \right\}$ .

## (7) Example

## (a) Problem

Obtain all eigenvalues of  $Ax = \lambda Bx$  and their corresponding eigenvectors, where matrices  $A$  and  $B$  are as follows:

$$A = \begin{bmatrix} 2 & 1 & 1 & 2 \\ 1 & 1 & 1 & 1 \\ 1 & 1 & 2 & 2 \\ 2 & 1 & 2 & 4 \end{bmatrix}$$

$$B = \begin{bmatrix} 153 & 31 & 58 & -58 \\ 31 & 153 & -53 & 58 \\ 58 & -58 & 153 & 31 \\ -58 & 58 & 31 & 153 \end{bmatrix}$$

## (b) Input data

Matrix  $A$ , lna=11, n=4, matrix  $B$ , lnb=11 and nt=2.

## (c) Main program

```

/*      C interface example for ASL_qcgsaa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int na=11;

```

```

int nn;
double *b;
int nb=11;
double *e;
double *wk;
int nt;
int ierr;
int i,j,k;
FILE *fp;

int mod;

fp = fopen( "qcgsaa.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_qcgsaa ***\n" );
printf( "\n    ** Input **\n\n" );

fscanf( fp, "%d", &nn );
fscanf( fp, "%d", &nt );

mod = nn % 4;

a = ( double * )malloc((size_t)( sizeof(double) * (na*nn) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double * )malloc((size_t)( sizeof(double) * (nb*nn) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

e = ( double * )malloc((size_t)( sizeof(double) * nn ));
if( e == NULL )
{
    printf( "no enough memory for array e\n" );
    return -1;
}

wk = ( double * )malloc((size_t)( sizeof(double) * (2*nn) ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

printf( "\tn = %6d\n", nn );
printf( "\tnt= %6d\n\n", nt );

printf( "\tInput Matrix a\n\n" );
for( i=0 ; i<nn ; i++ )
{
    for( j=i ; j<nn ; j++ )
    {
        fscanf( fp, "%lf", &a[i+na*j] );
    }
}

for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g", a[j+na*i] );
    }
    for( j=i ; j<nn ; j++ )
    {
        printf( "%8.3g", a[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tInput Matrix b\n\n" );
for( i=0 ; i<nn ; i++ )
{
    for( j=i ; j<nn ; j++ )
    {
        fscanf( fp, "%lf", &b[i+nb*j] );
    }
}

for( i=0 ; i<nn ; i++ )

```



```

{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g", b[j+nb*i] );
    }
    for( j=i ; j<nn ; j++ )
    {
        printf( "%8.3g", b[i+nb*j] );
    }
    printf( "\n" );
}
fclose( fp );

ierr = ASL_qcgsaa(a, na, nn, b, nb, e, wk, nt);

printf( "\n    ** Output **\n\n" );
printf( "\t(ierr = %6d\n", ierr );

for( k=0 ; k<nn-3 ; k = k+4 )
{
    printf( "\n\t" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "Eigenvalue  " );
    }
    printf( "\n\t" );
    for( i=k ; i<k+4 ; i++ )
    {
        printf( " %8.3g  ", e[i] );
    }
    printf( "\n" );

    printf( "\t" );
    for( i=0 ; i<4 ; i++ )
    {
        printf( "Eigenvector  " );
    }
    printf( "\n" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "\t" );
        for( i=k ; i<k+4 ; i++ )
        {
            printf( " %8.3g  ", a[j+na*i] );
        }
        printf( "\n" );
    }
}

if( mod != 0 )
{
    printf( "\n\t" );
    for( i= nn-mod ; i<nn ; i++ )
    {
        printf( "Eigenvalue  " );
    }
    printf( "\n\t" );
    for( i= nn-mod ; i<nn ; i++ )
    {
        printf( " %8.3g  ", e[i] );
    }
    printf( "\n" );

    printf( "\t" );
    for( i= nn-mod ; i<nn ; i++ )
    {
        printf( "Eigenvector  " );
    }
    printf( "\n" );
    for( j=0 ; j<nn ; j++ )
    {
        printf( "\t" );
        for( i= nn-mod ; i<nn ; i++ )
        {
            printf( " %8.3g  ", a[j+na*i] );
        }
        printf( "\n" );
    }
}

free( a );
free( b );
free( e );
free( wk );

```

```
    }
    return 0;
}
```

(d) Output results

```
*** ASL_qcgsaa ***
```

```
** Input **
```

```
n =      4
nt=      2
```

```
Input Matrix a
```

```
   2   1   1   2
   1   1   1   1
   1   1   2   2
   2   1   2   4
```

```
Input Matrix b
```

```
  153   31   58  -58
   31  153  -58   58
   58  -58  153   31
  -58   58   31  153
```

```
** Output **
```

```
ierr =      0
```

Eigenvalue	Eigenvalue	Eigenvalue	Eigenvalue
0.000648	0.00537	0.0274	0.217
Eigenvector	Eigenvector	Eigenvector	Eigenvector
0.0294	0.0498	-0.0161	0.205
-0.0469	0.0377	0.0686	-0.193
0.0311	-0.0194	0.086	-0.192
-0.0196	-0.0332	0.00145	0.21

### 5.5.2 ASL\_qcgsan, ASL\_pcgsan

#### All Eigenvalues of a Real Symmetric Matrix (Generalized Eigenvalue Problem $Ax = \lambda Bx$ , $B$ : Positive)

(1) **Function**

ASL\_qcgsan or ASL\_pcgsan uses the Cholesky method to transform the real symmetric matrix (two-dimensional array type) (upper triangular type) generalized eigenvalue problem  $Ax = \lambda Bx$  ( $A$ : Real symmetric matrix,  $B$ : Positive real symmetric matrix) to a standard eigenvalue problem and uses the Householder method and QR method to obtain all eigenvalues.

(2) **Usage**

Double precision:

ierr = ASL\_qcgsan (a, lna, n, b, lnb, e, w1, nt);

Single precision:

ierr = ASL\_pcgsan (a, lna, n, b, lnb, e, w1, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$ and $B$
4	b	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lnb \times n$	Input	Positive symmetric matrix $B$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Output	Eigenvalues
7	w1	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $0 < n \leq lna, lnb$

(b)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]/b[0]$ is performed.
2100	$B$ has a diagonal element very close to zero.	Processing continues.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000+i	The sequence did not converge in the step where the eigenvalue is obtained. ( $1 \leq i \leq n$ )	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.)

## (6) Notes

- (a) Data should be stored only in the upper triangular portions of arrays a and b.
- (b) Eigenvalues are stored in ascending order.

### 5.5.3 ASL\_qcgsss, ASL\_pcgsss

#### Eigenvalues and Eigenvectors of a Real Symmetric Matrix (Generalized Eigenvalue Problem $Ax = \lambda Bx$ , $B$ : Positive)

(1) **Function**

ASL\_qcgsss or ASL\_pcgsss uses the Cholesky method to transform the real symmetric matrix (two-dimensional array type) (upper triangular type) generalized eigenvalue problem  $A\mathbf{x} = \lambda B\mathbf{x}$  ( $A$ : Real symmetric matrix,  $B$ : Positive real symmetric matrix) to a standard eigenvalue problem and uses the Householder method and the root-free QR method or Bisection method to obtain the  $m$  largest eigenvalues or  $m$  smallest eigenvalues, and uses the reverse iterative method to obtain the eigenvectors.

(2) **Usage**

Double precision:

```
ierr = ASL_qcgsss (a, lna, n, b, lnb, eps, e, m, ve, lnv, isw, iw1, w1, nt);
```

Single precision:

```
ierr = ASL_pcgsss (a, lna, n, b, lnb, eps, e, m, ve, lnv, isw, iw1, w1, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$ and $B$
4	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Positive symmetric matrix $B$ (two-dimensional array type) (upper triangular type)
				Output	The strict upper triangular portion is not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
7	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
8	m	I	1	Input	The number of m of eigenvalues to be obtained.
9	ve	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnv} \times m$	Output	Eigenvectors (column vector) corresponding to each eigenvalue
10	lnv	I	1	Input	Adjustable dimension of array ve
11	isw	I	1	Input	Processing switch $\text{isw} \geq 0$ : Obtain m eigenvalues from the largest one. $\text{isw} < 0$ : Obtain m eigenvalues from the smallest one.
12	iw1	I*	m	Output	Eigenvector flag (See Note (a))
13	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$9 \times n$	Work	Work area
14	nt	I	1	Input	Number of tasks to be generated
15	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < n \leq \text{lna}$ ,  $\text{lmb}$ ,  $\text{lnv}$
- (b)  $0 < m \leq n$
- (c)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$e[0] \leftarrow a[0]/b[0]$ and $ve[0] \leftarrow 1.0/\sqrt{b[0]}$ are performed.
2000	The maximum number of iterations was exceeded by the inverse iterations for obtaining eigenvectors.	Some eigenvectors are obtained with low precision, and processing continues. (See Note (e))
2100	$B$ has a diagonal element very close to zero.	Some eigenvectors may be obtained with low precision, and processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	
4000	$B$ was not positive definite.	

## (6) Notes

- (a) Data should be stored only in the upper triangular portions of arrays  $a$  and  $b$ .
- (b) If  $\text{isw} \geq 0$ , the eigenvalues are stored in descending order. If  $\text{isw} < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $\text{eps} \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically.  $\text{eps}$  is used to obtain eigenvalues by using the Bisection method.
- (e) If the maximum number of iterations is exceeded when using the inverse iteration method ( $\text{ierr} = 2000$  is output), the following processing is performed.  
 If  $\text{iw1}[i - 1] = 0$ : The  $i$ -th eigenvector calculation is normally terminated.  
 If  $\text{iw1}[i - 1] \neq 0$ : The convergence condition is not satisfied for the  $i$ -th eigenvector calculation, and the eigenvector precision is low. In this case, the iteration count is set for  $\text{iw1}[i - 1]$ .  
 If processing is normally terminated ( $\text{ierr} = 0$  is output),  $\text{iw1}[i - 1] = 0$  is set.
- (f) Eigenvectors  $\mathbf{v}_i$  are an orthonormal set so that  $\mathbf{v}_j^T B \mathbf{v}_k = \delta_{j,k}$
- (g) If eigenvectors are not required, use 5.5.4  $\left\{ \begin{array}{l} \text{ASL\_qcgssn} \\ \text{ASL\_pcgssn} \end{array} \right\}$ .

## (7) Example

- (a) Problem Obtain all eigenvalues of
- $Ax = \lambda Bx$
- and their corresponding eigenvectors, where matrices
- $A$
- and
- $B$
- are as follows:

$$A = \begin{bmatrix} 611 & 196 & -192 & 407 & -8 & -52 & -49 & 29 \\ 196 & 899 & 113 & -192 & -71 & -43 & -8 & -44 \\ -192 & 113 & 899 & 196 & 61 & 49 & 8 & 52 \\ 407 & -192 & 196 & 611 & 8 & 44 & 59 & -23 \\ -8 & -71 & 61 & 8 & 411 & -599 & 208 & 208 \\ -52 & -43 & 49 & 44 & -599 & 411 & 208 & 208 \\ -49 & -8 & 8 & 59 & 208 & 208 & 99 & -911 \\ 29 & -44 & 52 & -23 & 208 & 208 & -911 & 99 \end{bmatrix}$$

$$B = \begin{bmatrix} 170 & 18 & 33 & -21 & -17 & 13 & 25 & -36 \\ 18 & 171 & -21 & 22 & 13 & -17 & -36 & 25 \\ 33 & -21 & 171 & 18 & 25 & -36 & -17 & 13 \\ -21 & 22 & 18 & 171 & -36 & 25 & 13 & -17 \\ -17 & 13 & 25 & -36 & 171 & 18 & 33 & -21 \\ 13 & -17 & -36 & 25 & 18 & 171 & -21 & -3 \\ 25 & -36 & -17 & 13 & 33 & -21 & 171 & 18 \\ -36 & 25 & 13 & -17 & -21 & -3 & 18 & 171 \end{bmatrix}$$

- (b) Input data

Matrix  $A$ , lna=11, n=8, matrix  $B$ , lnb=11, eps=-1.0, m=2, lnv=10 and isw=-1.

- (c) Main program

```

/*      C interface example for ASL_qcgsss */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a;
    int na=11;
    int nn;
    double *b;
    int nb=11;
    double cepts= -1.0;
    double *e;
    int mm;
    double *ve;
    int nv=10;
    int ksw= -1;
    int *kw1;
    double *wk;
    int nt;
    int ierr;
    int i,j,k;
    FILE *fp;

    int mod;

    fp = fopen( "qcgsss.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_qcgsss ***\n" );
    printf( "\n      ** Input **\n\n" );

    fscanf( fp, "%d", &nn );
    fscanf( fp, "%d", &mm );
    fscanf( fp, "%d", &nt );

    mod = mm % 4;

```



```

a = ( double * )malloc((size_t)( sizeof(double) * (na*nn) ));
if( a == NULL )
{
    printf( "no enough memory for array a\n" );
    return -1;
}

b = ( double * )malloc((size_t)( sizeof(double) * (nb*nn) ));
if( b == NULL )
{
    printf( "no enough memory for array b\n" );
    return -1;
}

e = ( double * )malloc((size_t)( sizeof(double) * mm ));
if( e == NULL )
{
    printf( "no enough memory for array e\n" );
    return -1;
}

ve = ( double * )malloc((size_t)( sizeof(double) * (nv*mm) ));
if( ve == NULL )
{
    printf( "no enough memory for array ve\n" );
    return -1;
}

kw1 = ( int * )malloc((size_t)( sizeof(int) * mm ));
if( kw1 == NULL )
{
    printf( "no enough memory for array kw1\n" );
    return -1;
}

wk = ( double * )malloc((size_t)( sizeof(double) * (9*nn) ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

printf( "\tn = %6d\n", nn );
printf( "\tm = %6d\n", mm );
printf( "\tnt = %6d\n\n", nt );

printf( "\tInput Matrix a\n\n" );
for( i=0 ; i<nn ; i++ )
{
    for( j=i ; j<nn ; j++ )
    {
        fscanf( fp, "%lf", &a[i+na*j] );
    }
}

for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g", a[j+na*i] );
    }
    for( j=i ; j<nn ; j++ )
    {
        printf( "%8.3g", a[i+na*j] );
    }
    printf( "\n" );
}

printf( "\n\tInput Matrix b\n\n" );
for( i=0 ; i<nn ; i++ )
{
    for( j=i ; j<nn ; j++ )
    {
        fscanf( fp, "%lf", &b[i+nb*j] );
    }
}

for( i=0 ; i<nn ; i++ )
{
    printf( "\t" );
    for( j=0 ; j<i ; j++ )
    {
        printf( "%8.3g", b[j+nb*i] );
    }
    for( j=i ; j<nn ; j++ )
    {
        printf( "%8.3g", b[i+nb*j] );
    }
    printf( "\n" );
}

```

```

    }
    fclose( fp );

    ierr = ASL_qcgsss(a, na, nn, b, nb, ceps, e, mm, ve, nv, ksw, kw1, wk, nt);

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    for( k=0 ; k<mm-3 ; k = k+4 )
    {
        printf( "\n" );
        printf( "\t" );
        for( i=0 ; i<4 ; i++ )
        {
            printf( "Eigenvalue  " );
        }
        printf( "\n" );
        printf( "\t" );
        for( i=k ; i<k+4 ; i++ )
        {
            printf( " %8.3g  ", e[i] );
        }
        printf( "\n" );
        printf( "\t" );
        for( i=0 ; i<4 ; i++ )
        {
            printf( "Eigenvector  " );
        }
        printf( "\n" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "\t" );
            for( i=k ; i<k+4 ; i++ )
            {
                printf( " %8.3g  ", ve[j+nv*i] );
            }
            printf( "\n" );
        }
    }

    if( mod != 0 )
    {
        printf( "\n" );
        printf( "\t" );
        for( i= mm-mod ; i<mm ; i++ )
        {
            printf( "Eigenvalue  " );
        }
        printf( "\n" );
        printf( "\t" );
        for( i= mm-mod ; i<mm ; i++ )
        {
            printf( " %8.3g  ", e[i] );
        }
        printf( "\n" );
        printf( "\t" );
        for( i= nn-mod ; i<nn ; i++ )
        {
            printf( "Eigenvector  " );
        }
        printf( "\n" );
        for( j=0 ; j<nn ; j++ )
        {
            printf( "\t" );
            for( i= mm-mod ; i<mm ; i++ )
            {
                printf( " %8.3g  ", ve[j+nv*i] );
            }
            printf( "\n" );
        }
    }

    free( a );
    free( b );
    free( e );
    free( ve );
    free( kw1 );
    free( wk );

    return 0;
}

```

(d) Output results

```

*** ASL_qcgsss ***
** Input **
n =      8
m =      2
nt=      2
Input Matrix a
   611   196  -192   407   -8   -52   -49   29
   196   899   113  -192   -71  -43   -8  -44
  -192   113   899   196    61   49    8   52
   407  -192   196   611    8    44   59  -23
    -8   -71    61    8   411  -599  208  208
   -52  -43    49    44  -599  411  208  208
   -49   -8    8    59  208  208   99  -911
    29  -44   52   -23  208  208  -911   99
Input Matrix b
   170   18   33   -21   -17   13   25  -36
    18   171  -21   22   13  -17  -36   25
    33  -21   171   18   25  -36  -17   13
   -21   22   18   171  -36   25   13  -17
   -17   13   25  -36   171   18   33  -21
    13  -17  -36   25   18   171  -21  -3
    25  -36  -17   13   33  -21   171   18
   -36   25   13  -17  -21   -3   18   171
** Output **
ierr =      0
Eigenvalue  Eigenvalue
   -5.3    -1.04e-15
Eigenvector Eigenvector
0.000789   -0.00329
 0.00146   -0.00658
0.000624   0.00658
-0.00168   0.00329
-0.0247    -0.0461
 -0.019    -0.0461
 0.0479    -0.023
 0.0445    -0.023

```

#### 5.5.4 ASL\_qcgssn, ASL\_pcgssn

##### Eigenvalues of a Real Symmetric Matrix (Generalized Eigenvalue Problem $Ax = \lambda Bx$ , $B$ : Positive)

(1) **Function**

ASL\_qcgssn or ASL\_pcgssn uses the Cholesky method to transform the real symmetric matrix (two-dimensional array type) (upper triangular type) generalized eigenvalue problem  $Ax = \lambda Bx$  ( $A$ : Real symmetric matrix,  $B$ : Positive real symmetric matrix) to a standard eigenvalue problem and uses the Householder method and the root-free QR method or Bisection method to obtain the  $m$  largest eigenvalues or  $m$  smallest eigenvalues.

(2) **Usage**

Double precision:

```
ierr = ASL_qcgssn (a, lna, n, b, lnb, eps, e, m, isw, w1, nt);
```

Single precision:

```
ierr = ASL_pcgssn (a, lna, n, b, lnb, eps, e, m, isw, w1, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real symmetric matrix $A$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrix $A$ and $B$
4	b	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Positive symmetric matrix $B$ (two-dimensional array type) (upper triangular type)
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	eps	$\begin{Bmatrix} D \\ R \end{Bmatrix}$	1	Input	Parameter that assigns an upper limit to the absolute error for use in the eigenvalue convergence test. (See Note (d))
7	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	m	Output	Eigenvalues
8	m	I	1	Input	The number of m of eigenvalues to be obtained.
9	isw	I	1	Input	Processing switch isw $\geq$ 0: Obtain m eigenvalues from the largest one. isw $<$ 0: Obtain m eigenvalues from the smallest one.
10	w1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$5 \times n$	Work	Work area
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $0 < n \leq \text{lna}, \text{lnb}$
- (b)  $0 < m \leq n$
- (c)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]/b[0]$ is performed.
2100	$B$ has a diagonal element very close to zero.	Processing continues.
3000	Restriction (a) or (b) was not satisfied.	Processing is aborted.
3010	Restriction (c) was not satisfied.	
4000	$B$ was not positive definite.	

(6) **Notes**

- (a) Data should be stored only in the upper triangular portions of arrays a and b.
- (b) If  $isw \geq 0$ , the eigenvalues are stored in descending order. If  $isw < 0$ , they are stored in ascending order.
- (c) Eigenvalue calculations are appropriately divided up between the root-free QR method and Bisection method internally.
- (d) If  $eps \leq 0$ , the optimum value is automatically set internally. Normally, a negative value should be set so that this value will be set automatically. eps is used to obtain eigenvalues by using the Bisection method.

## 5.6 GENERALIZED EIGENVALUE PROBLEM FOR REAL SYMMETRIC MATRICES (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) ( $ABx = \lambda x$ )

### 5.6.1 ASL\_qcgjaa, ASL\_pcgjaa

All Eigenvalues and All Eigenvectors of Real Symmetric Matrices (Generalized Eigenvalue Problem  $ABx = \lambda x$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$ABx = \lambda x$$

( $A$ : Real symmetric,  $B$ : Positive real symmetric) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  and corresponding all eigenvectors  $x$ .

(2) **Usage**

Double precision:

```
ierr = ASL_qcgjaa (a, lna, n, b, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_pcgjaa (a, lna, n, b, lnb, e, work, nt);
```

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	$lna \times n$	Input	Real symmetric matrix $A$
				Output	Eigenvectors $x$
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrices $A$ and $B$
4	b	$\begin{cases} D* \\ R* \end{cases}$	$lnb \times n$	Input	Real symmetric matrix $B$
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues $\lambda$
7	work	$\begin{cases} D* \\ R* \end{cases}$	$2 \times n$	Work	Work area
8	nt	I	1	Input	Number of tasks
9	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $1 \leq n \leq \text{lna}, \text{lnb}$   
 (b)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$e[0] \leftarrow a[0] * b[0]$ and $a[0] \leftarrow \frac{1.0}{\sqrt{b[0]}}$ are performed.
2100	$B$ has a diagonal element very close to zero.	Some results may be obtained with low precision
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
$5000+i$	The sequence did not converge in the step where the eigenvalue was obtained ( $1 \leq i \leq n$ )	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.) No eigenvector is obtained at this time.

## (6) Notes

- (a) Arrays  $a$  and  $b$  should be stored only in the upper triangular portions.  
 (b) Eigenvalues are stored in ascending order.  
 (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^T B v_k = \delta_{j,k}$   
 (d) 5.6.2  $\left\{ \begin{array}{l} \text{ASL\_qcgjan} \\ \text{ASL\_pcgjan} \end{array} \right\}$  should be used if the eigenvectors are not needed.  
 (e) 5.7.1  $\left\{ \begin{array}{l} \text{ASL\_qcgkaa} \\ \text{ASL\_pcgkaa} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## (7) Example

## (a) Problem

Obtain all eigenvalues and their corresponding eigenvectors non-symmetric matrix  $AB$  when  $A$  and  $B$  are positive symmetric matrices.

$$A = \begin{bmatrix} 1.07692 & 0.28571 & 0.09733 & 0.04887 \\ 0.28571 & 1.02041 & 0.26316 & 0.08610 \\ 0.09733 & 0.26316 & 1.00917 & 0.25676 \\ 0.04887 & 0.08610 & 0.25676 & 1.00518 \end{bmatrix}$$

$$B = \begin{bmatrix} 1.04762 & 0.18841 & 0.05996 & 0.02968 \\ 0.18841 & 1.01235 & 0.17460 & 0.05314 \\ 0.05996 & 0.17460 & 1.00552 & 0.17073 \\ 0.02968 & 0.05314 & 0.17073 & 1.00312 \end{bmatrix}$$



**Note** The input data  $A$  and  $B$  are defined as

$$A = P(3.0, 4), \quad B = P(5.0, 4)$$

where

$$P(a^2, n)_{i,j} = 2 \int_0^\infty \cos(ait) \cos(ajt) e^{-t} dt \quad (1 \leq i \leq n; 1 \leq j \leq n).$$

All eigenvalues of the each matrix exist within a finite interval which is independent of  $n$ .

(b) Input data

$n = 4$ ,  $lna = lnb = 4$ ,  $nt = 2$  and symmetric matrices  $A, B$ .

(c) Main program

```

/*      C interface example for ASL_qcgjaa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a, *b, *e, *work;
    double one,tre,fiv;
    int i,j,n,ierr;
    int nt;
    int lna, lnb;
    n=4;
    lna=4;lnb=4;nt=4;
    one=1.0;tre=3.0;fiv=5.0;
    printf( "      *** ASL_qcgjaa ***\n" );
    printf( "\n      ** Input **\n\n" );
    a = ( double * )malloc((size_t)(sizeof(double)*n*n));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n");
        return -1;
    }
    b = ( double * )malloc((size_t)(sizeof(double)*n*n));
    if( b == NULL )
    {
        printf( "no enough memory for array b\n");
        return -1;
    }
    e = ( double * )malloc((size_t)(sizeof(double)*n));
    if ( e == NULL )
    {
        printf( "no enough memory for array e\n");
        return -1;
    }
    work = ( double * )malloc((size_t)(sizeof(double)*(2*n)));
    if ( work == NULL )
    {
        printf( "no enough memory for array work\n");
        return -1;
    }
    printf( "\tn      = %6d\n" , n );
    printf( "\tlna   = %6d\n", lna );
    printf( "\tlnb   = %6d\n", lnb );
    printf( "\tnt    = %6d\n", nt );
    for(i=0; i<n; i++)
    {
        for(j=0; j<n; j++)
        {
            a[i+n*j]=one/(one+tre*(i+j+2)*(i+j+2))
                +one/(one+tre*(i-j)*(i-j));
            b[i+n*j]=one/(one+fiv*(i+j+2)*(i+j+2))
                +one/(one+fiv*(i-j)*(i-j));
        }
    }
    printf( "\n\tInput Matrix a\n\n" );
    for(i=0; i<n; i++)
    {
        printf( "\t" );
        for(j=0; j<n; j++)
        {
            printf( "%8.3g",a[i+n*j]);
        }
        printf( "\n" );
    }
    printf( "\n\tInput Matrix b\n\n" );
    for(i=0; i<n; i++)
    {
        printf( "\t" );
        for(j=0; j<n; j++)
        {

```

```

        printf( "%8.3g",b[i+n*j] );
    }
    printf( "\n" );
}
ierr = ASL_qcgjaa(a, lna, n, b, lnb, e, work, nt);
printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\t      Eigenvalue   " );
printf( "\n\t" );
for(j=0; j<n; j++)
{
    printf( "%8.3g",e[j] );
}
printf( "\n\t      Eigenvector   " );
for(i=0; i<n; i++)
{
    printf( "\n\t" );
    for(j=0; j<n; j++)
    {
        printf( "%8.3g",a[i+n*j] );
    }
}
printf( "\n" );
free(a);
free(b);
free(e);
free(work);
return 0;
}

```

(d) Output results

```

*** ASL_qcgjaa ***

** Input **

n   =   4
lna =   4
lnb =   4
nt  =   4

Input Matrix a

  1.08  0.286  0.0973  0.0489
  0.286  1.02  0.263  0.0861
  0.0973 0.263  1.01  0.257
  0.0489 0.0861 0.257  1.01

Input Matrix b

  1.05  0.188  0.06  0.0297
  0.188 1.01  0.175 0.0531
  0.06  0.175 1.01  0.171
  0.0297 0.0531 0.171  1

** Output **

ierr =   0

      Eigenvalue
0.503  0.706  1.15  2.13
      Eigenvector
-0.36 -0.573  0.6  0.414
 0.702 0.497  0.257 0.494
-0.718 0.42 -0.389 0.46
 0.406 -0.626 -0.604 0.324

```

### 5.6.2 ASL\_qcgjan, ASL\_pcgjan

#### All Eigenvalues of Real Symmetric Matrices (Generalized Eigenvalue Problem $ABx = \lambda x$ , $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$ABx = \lambda x$$

( $A$ : Real symmetric,  $B$ : Positive real symmetric) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$ .

(2) **Usage**

Double precision:

ierr = ASL\_qcgjan (a, lna, n, b, lnb, e, work, nt);

Single precision:

ierr = ASL\_pcgjan (a, lna, n, b, lnb, e, work, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	$lna \times n$	Input	Real symmetric matrix $A$
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrices $A$ and $B$
4	b	$\begin{cases} D* \\ R* \end{cases}$	$lnb \times n$	Input	Real symmetric matrix $B$
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues $\lambda$
7	work	$\begin{cases} D* \\ R* \end{cases}$	$2 \times n$	Work	Work area
8	nt	I	1	Input	Number of tasks
9	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $1 \leq n \leq lna, lnb$

(b)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ is performed.
2100	$B$ has a diagonal element very close to zero.	Some results may be obtained with low precision
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
$5000+i$	The sequence did not converge in the step where the eigenvalue was obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ (However, the order is irregular).

(6) Notes

- (a) Arrays a and b should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) 5.6.1  $\left\{ \begin{array}{l} \text{ASL\_qcgjaa} \\ \text{ASL\_pcgjaa} \end{array} \right\}$  should be used if the eigenvectors are needed.
- (d) 5.7.2  $\left\{ \begin{array}{l} \text{ASL\_qcgkan} \\ \text{ASL\_pcgkan} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## 5.7 GENERALIZED EIGENVALUE PROBLEM FOR REAL SYMMETRIC MATRICES (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) ( $BAx = \lambda x$ )

### 5.7.1 ASL\_qcgkaa, ASL\_pcgkaa

All Eigenvalues and All Eigenvectors of Real Symmetric Matrices (Generalized Eigenvalue Problem  $BAx = \lambda x$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$BAx = \lambda x$$

( $A$ : Real symmetric,  $B$ : Positive real symmetric) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  and corresponding all eigenvectors  $x$ .

(2) **Usage**

Double precision:

ierr = ASL\_qcgkaa (a, lna, n, b, lnb, e, work, nt);

Single precision:

ierr = ASL\_pcgkaa (a, lna, n, b, lnb, e, work, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lna \times n$	Input	Real symmetric matrix $A$
				Output	Eigenvectors $x$
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrices $A$ and $B$
4	b	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$lnb \times n$	Input	Real symmetric matrix $B$
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	n	Output	Eigenvalues $\lambda$
7	work	$\begin{Bmatrix} D* \\ R* \end{Bmatrix}$	$2 \times n$	Work	Work area
8	nt	I	1	Input	Number of tasks
9	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $1 \leq n \leq \text{lna}, \text{lnb}$   
 (b)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ and $a[0] \leftarrow \sqrt{b[0]}$ are performed.
2100	$B$ has a diagonal element very close to zero.	Some results may be obtained with low precision
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
$5000+i$	The sequence did not converge in the step where the eigenvalue was obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ (However, the order is irregular). No eigenvector is obtained at this time.

## (6) Notes

- (a) Arrays  $a$  and  $b$  should be stored only in the upper triangular portions.  
 (b) Eigenvalues are stored in ascending order.  
 (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^T B^{-1} v_k = \delta_{j,k}$   
 (d) 5.7.2  $\left\{ \begin{array}{l} \text{ASL-qcgkan} \\ \text{ASL-pcgkan} \end{array} \right\}$  should be used if the eigenvectors are not needed.  
 (e) 5.6.1  $\left\{ \begin{array}{l} \text{ASL-qcgjaa} \\ \text{ASL-pcgjaa} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## (7) Example

## (a) Problem

Obtain all eigenvalues and their corresponding eigenvectors non-symmetric matrix  $AB$  when  $A$  and  $B$  are positive symmetric matrices.

$$A = \begin{bmatrix} 1.07692 & 0.28571 & 0.09733 & 0.04887 \\ 0.28571 & 1.02041 & 0.26316 & 0.08610 \\ 0.09733 & 0.26316 & 1.00917 & 0.25676 \\ 0.04887 & 0.08610 & 0.25676 & 1.00518 \end{bmatrix}$$

$$B = \begin{bmatrix} 1.04762 & 0.18841 & 0.05996 & 0.02968 \\ 0.18841 & 1.01235 & 0.17460 & 0.05314 \\ 0.05996 & 0.17460 & 1.00552 & 0.17073 \\ 0.02968 & 0.05314 & 0.17073 & 1.00312 \end{bmatrix}$$

**Note** The input data  $A$  and  $B$  are defined as

$$A = P(3.0, 4), \quad B = P(5.0, 4)$$

where

$$P(a^2, n)_{i,j} = 2 \int_0^\infty \cos(ait) \cos(ajt) e^{-t} dt (1 \leq i \leq n; 1 \leq j \leq n).$$

All eigenvalues of the each matrix exist within a finite interval which is independent of  $n$ .

(b) Input data

$n=4$ ,  $lna=lnb=4$ ,  $nt=2$  and symmetric matrices  $A$ ,  $B$ .

(c) Main program

```

/*      C interface example for ASL_qcgkaa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a, *b, *e, *work;
    double one,tre,fiv;
    int i,j,n,ierr;
    int nt;
    int lna, lnb;
    n=4;
    lna=4;lnb=4;nt=4;
    one=1.0;tre=3.0;fiv=5.0;
    printf( "      *** ASL_qcgkaa ***\n" );
    printf( "\n      ** Input **\n\n" );
    a = ( double * )malloc((size_t)(sizeof(double)*n*n));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n");
        return -1;
    }
    b = ( double * )malloc((size_t)(sizeof(double)*n*n));
    if( b == NULL )
    {
        printf( "no enough memory for array b\n");
        return -1;
    }
    e = ( double * )malloc((size_t)(sizeof(double)*n));
    if ( e == NULL )
    {
        printf( "no enough memory for array e\n");
        return -1;
    }
    work = ( double * )malloc((size_t)(sizeof(double)*(2*n)));
    if ( work == NULL )
    {
        printf( "no enough memory for array work\n");
        return -1;
    }
    printf( "\tn      = %6d\n" , n );
    printf( "\tlna   = %6d\n", lna );
    printf( "\tlnb   = %6d\n", lnb );
    printf( "\tnt    = %6d\n", nt );
    for(i=0; i<n; i++)
    {
        for(j=0; j<n; j++)
        {
            a[i+n*j]=one/(one+tre*(i+j+2)*(i+j+2))
                +one/(one+tre*(i-j)*(i-j));
            b[i+n*j]=one/(one+fiv*(i+j+2)*(i+j+2))
                +one/(one+fiv*(i-j)*(i-j));
        }
    }
    printf( "\n\tInput Matrix a\n\n" );
    for(i=0; i<n; i++)
    {
        printf( "\t" );
        for(j=0; j<n; j++)
        {
            printf( "%8.3g",a[i+n*j]);
        }
        printf( "\n" );
    }
    printf( "\n\tInput Matrix b\n\n" );
    for(i=0; i<n; i++)
    {
        printf( "\t" );
        for(j=0; j<n; j++)
        {

```

```

        printf( "%8.3g",b[i+n*j] );
    }
    printf( "\n" );
}
ierr = ASL_qcgkaa(a, lna, n, b, lnb, e, work, nt);
printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\t      Eigenvalue   " );
printf( "\n\t" );
for(j=0; j<n; j++)
{
    printf( "%8.3g",e[j] );
}
printf( "\n\t      Eigenvector   " );
for(i=0; i<n; i++)
{
    printf( "\n\t" );
    for(j=0; j<n; j++)
    {
        printf( "%8.3g",a[i+n*j] );
    }
}
printf( "\n" );
free(a);
free(b);
free(e);
free(work);
return 0;
}

```

(d) Output results

```

*** ASL_qcgkaa ***
** Input **
n      =      4
lna    =      4
lnb    =      4
nt     =      4

Input Matrix a
      1.08  0.286  0.0973  0.0489
      0.286  1.02  0.263  0.0861
      0.0973 0.263  1.01  0.257
      0.0489 0.0861 0.257  1.01

Input Matrix b
      1.05  0.188  0.06  0.0297
      0.188 1.01  0.175 0.0531
      0.06  0.175 1.01  0.171
      0.0297 0.0531 0.171  1

** Output **
ierr =      0

      Eigenvalue
0.503  0.706  1.15  2.13
      Eigenvector
-0.276 -0.5  0.635  0.564
      0.54  0.436  0.273  0.676
-0.551  0.368 -0.414  0.629
      0.311 -0.547 -0.641  0.442

```



### 5.7.2 ASL\_qcgkan, ASL\_pcgkan

#### All Eigenvalues of Real Symmetric Matrices (Generalized Eigenvalue Problem $BAx = \lambda x$ , $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$BAx = \lambda x$$

( $A$  : Real symmetric,  $B$ : Positive real symmetric) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  .

(2) **Usage**

Double precision:

ierr = ASL\_qcgkan (a, lna, n, b, lnb, e, work, nt);

Single precision:

ierr = ASL\_pcgkan (a, lna, n, b, lnb, e, work, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	$lna \times n$	Input	Real symmetric matrix $A$
				Output	Input-time contents are not retained.
2	lna	I	1	Input	Adjustable dimension of array a
3	n	I	1	Input	Order of matrices $A$ and $B$
4	b	$\begin{cases} D* \\ R* \end{cases}$	$lnb \times n$	Input	Real symmetric matrix $B$
				Output	Input-time contents are not retained.
5	lnb	I	1	Input	Adjustable dimension of array b
6	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues $\lambda$
7	work	$\begin{cases} D* \\ R* \end{cases}$	$2 \times n$	Work	Work area
8	nt	I	1	Input	Number of tasks
9	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $1 \leq n \leq lna, lnb$

(b)  $nt \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ is performed.
2100	$B$ has a diagonal element very close to zero.	Some results may be obtained with low precision
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
$5000+i$	The sequence did not converge in the step where the eigenvalue was obtained ( $1 \leq i \leq n$ ).	Eigenvalues obtained by this time are entered in $e[0], \dots, e[i-2]$ . (However, the order is irregular.)

## (6) Notes

- (a) Arrays a and b should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) 5.7.1  $\left\{ \begin{array}{l} \text{ASL\_qcgkaa} \\ \text{ASL\_pcgkaa} \end{array} \right\}$  should be used if the eigenvectors are needed.
- (d) 5.6.2  $\left\{ \begin{array}{l} \text{ASL\_qcgjan} \\ \text{ASL\_pcgjan} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

---

## 5.8 GENERALIZED EIGENVALUE PROBLEM ( $Az = \lambda Bz$ ) FOR HERMITIAN MATRICES (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE)

### 5.8.1 ASL\_hcgraa, ASL\_gcgraa

All Eigenvalues and All Eigenvectors of Hermitian Matrices (Generalized Eigenvalue Problem  $Az = \lambda Bz$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$Az = \lambda Bz$$

( $A$ : Hermitian,  $B$ : Positive Hermitian ) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  and corresponding all eigenvectors  $z$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hcgraa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_gcgraa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$
				Output	Real part of eigenvector $x$
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Imaginary part of eigenvector $x$
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues $\lambda$
9	work	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

- (a)  $1 \leq n \leq \text{lna}$  ,  $\text{lnb}$   
(b)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \frac{a[0]}{b[0]}$ and $a[0] \leftarrow \frac{1.0}{\sqrt{b[0]}}$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

## (6) Notes

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^* B v_k = \delta_{j,k}$
- (d) 5.8.2  $\left\{ \begin{array}{l} \text{ASL\_hcgran} \\ \text{ASL\_gcgran} \end{array} \right\}$  should be used if the eigenvectors are not needed.

## (7) Example

## (a) Problem

For Hermitian matrix of the degree 4

$$A = \begin{bmatrix} 8 & 3 & 1-2i & -1-2i \\ 3 & 9 & 1+2i & -1+2i \\ 1+2i & 1-2i & 10 & -3 \\ -1+2i & -1-2i & -3 & 11 \end{bmatrix}$$

and its conjugate Hermitian matrix

$$B = \begin{bmatrix} 8 & 3 & 1+2i & -1+2i \\ 3 & 9 & 1-2i & -1-2i \\ 1-2i & 1+2i & 10 & -3 \\ -1-2i & -1+2i & -3 & 11 \end{bmatrix}$$

obtain eigenvector of generalized eigenvalue problem.

## (b) Input data

$n=4$ ,  $lna=4$ , matrix  $A$ ,  $lnb=4$ ,  $nt=2$  and matrix  $B$ .

## (c) Main program

```

/*      C interface example for ASL_hcgraa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar, *br, *ai, *bi, *e, *work;
    int    i,j,n,ierr;
    int    nt;
    int    lna, lnb;
    n=4;
    lna=4;lnb=4,nt=4;

```

```

printf( "    *** ASL_hcgraa ***\n" );
printf( "\n    ** Input **\n\n" );
ar = ( double * )malloc(sizeof(double)*n*n);
if( ar == NULL )
{
    printf( "no enough memory for array ar\n");
    return -1;
}
br = ( double * )malloc(sizeof(double)*n*n);
if( br == NULL )
{
    printf( "no enough memory for array br\n");
    return -1;
}
ai = ( double * )malloc(sizeof(double)*n*n);
if( ai == NULL )
{
    printf( "no enough memory for array ai\n");
    return -1;
}
bi = ( double * )malloc(sizeof(double)*n*n);
if( bi == NULL )
{
    printf( "no enough memory for array bi\n");
    return -1;
}
e = ( double * )malloc(sizeof(double)*n);
if ( e == NULL )
{
    printf( "no enough memory for array e\n");
    return -1;
}
work = ( double * )malloc(sizeof(double)*4*n);
if ( work == NULL )
{
    printf( "no enough memory for array work\n");
    return -1;
}
printf( "\tn    = %6d\n" , n );
printf( "\tlna = %6d\n",lna );
printf( "\tlnb = %6d\n",lnb );
printf( "\tnt  = %6d\n",nt );
br[0+n*0]=8.0;bi[0+n*0]=0.0;
br[1+n*1]=9.0;bi[1+n*1]=0.0;
br[2+n*2]=10.0;bi[2+n*2]=0.0;
br[3+n*3]=11.0;bi[3+n*3]=0.0;
br[0+n*1]=3.0;bi[0+n*1]=0.0;
br[0+n*2]=1.0;bi[0+n*2]=2.0;
br[0+n*3]=-1.0;bi[0+n*3]=2.0;
br[1+n*2]=1.0;bi[1+n*2]=-2.0;
br[1+n*3]=-1.0;bi[1+n*3]=-2.0;
br[2+n*3]=-3.0;bi[2+n*3]=0.0;
for(i=1;i<n;i++)
{
    for(j=0;j<i;j++)
    {
        br[i+n*j]= br[j+n*i];
        bi[i+n*j]=-bi[j+n*i];
    }
}
for(i=0;i<n;i++)
{
    for(j=0;j<n;j++)
    {
        ar[i+n*j]= br[i+n*j];
        ai[i+n*j]=-bi[i+n*j];
    }
}
printf( "\n\tInput Matrix a\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g",ar[i+n*j]);
        printf( "%8.3g" ,ai[i+n*j]);
    }
    printf( "\n" );
}
printf( "\n\tInput Matrix b\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g",br[i+n*j]);
        printf( "%8.3g" ,bi[i+n*j]);
    }
    printf( "\n" );
}

```

```

ierr = ASL_hcgrra(ar,ai, lna, n, br,bi, lnb, e, work, nt);
printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\t      Eigenvalue  " );
printf( "\n\t" );
for(j=0; j<n; j++)
{
    printf( "      %8.3g      ",e[j]);
}
printf( "\n\t      Eigenvector  " );
for(i=0; i<n; i++)
{
    printf( "\n\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g,",ar[i+n*j]);
        printf( "%8.3g) ",ai[i+n*j]);
    }
}
printf( "\n" );
free(ar);
free(br);
free(ai);
free(bi);
free(e);
free(work);
return 0;
}

```

(d) Output results

```

*** ASL_hcgrra ***
** Input **
n      =      4
lna    =      4
lnb    =      4
nt     =      4

Input Matrix a
(      8,      0) (      3,      0) (      1,     -2) (     -1,     -2)
(      3,      0) (      9,      0) (      1,      2) (     -1,      2)
(      1,      2) (      1,     -2) (     10,      0) (     -3,      0)
(     -1,      2) (     -1,     -2) (     -3,      0) (     11,      0)

Input Matrix b
(      8,      0) (      3,      0) (      1,      2) (     -1,      2)
(      3,      0) (      9,      0) (      1,     -2) (     -1,     -2)
(      1,     -2) (      1,      2) (     10,      0) (     -3,      0)
(     -1,     -2) (     -1,      2) (     -3,      0) (     11,      0)

** Output **
ierr =      0

      Eigenvalue
0.231              1              1              4.33
      Eigenvector
(  0.175,      0) (      0,      0) (  0.211,      0) (  0.364,      0)
( -0.16, 0.00182) (-6.89e-19,-2.03e-16) (  0.211,-1.71e-18) ( -0.333,-0.00378)
(-0.00199, -0.149) ( -0.192, 0.00302) ( -0.0313,-4.35e-17) (-0.00414,  0.31)
( 0.00029, -0.138) (  0.192,-0.00302) (  0.0313,-1.66e-16) (0.000603,  0.287)

```

### 5.8.2 ASL\_hcgran, ASL\_gcgran

**All Eigenvalues of Hermitian Matrices (Generalized Eigenvalue Problem  $Az = \lambda Bz$ ,  $B$ : Positive)**

(1) **Function**

Generalized eigenvalue problem

$$Az = \lambda Bz$$

( $A$ : Hermitian,  $B$ : Positive Hermitian) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hcgran ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_gcgran ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```



## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues $\lambda$
9	work	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

(a)  $1 \leq n \leq \text{lna}, \text{lnb}$ (b)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow \frac{a[0]}{b[0]}$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

(6) **Notes**

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) 5.8.1  $\left\{ \begin{array}{l} \text{ASL\_hcgraa} \\ \text{ASL\_gcgraa} \end{array} \right\}$  should be used if the eigenvectors are needed.

---

## 5.9 GENERALIZED EIGENVALUE PROBLEM ( $ABz = \lambda z$ ) FOR HERMITIAN MATRICES (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE)

### 5.9.1 ASL\_hcgjaa, ASL\_gcgjaa

All Eigenvalues and All Eigenvectors of Hermitian Matrices (Generalized Eigenvalue Problem  $ABz = \lambda z$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$ABz = \lambda z$$

( $A$ : Hermitian,  $B$ : Positive Hermitian) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  and corresponding all eigenvectors  $z$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hcgjaa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_gcgjaa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{cases} D* \\ R* \end{cases}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$
				Output	Real part of eigenvector $x$
2	ai	$\begin{cases} D* \\ R* \end{cases}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Imaginary part of eigenvector $x$
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\begin{cases} D* \\ R* \end{cases}$	$\text{lnb} \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\begin{cases} D* \\ R* \end{cases}$	$\text{lnb} \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\begin{cases} D* \\ R* \end{cases}$	n	Output	Eigenvalues $\lambda$
9	work	$\begin{cases} D* \\ R* \end{cases}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

(a)  $1 \leq n \leq \text{lna}, \text{lnb}$ (b)  $\text{nt} \geq 1$ 

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ and $a[0] \leftarrow \frac{1.0}{\sqrt{b[0]}}$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

## (6) Notes

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^* B v_k = \delta_{j,k}$
- (d) 5.9.2  $\left\{ \begin{array}{l} \text{ASL\_hcgjan} \\ \text{ASL\_gcgjan} \end{array} \right\}$  should be used if the eigenvectors are not needed.
- (e) 5.10.1  $\left\{ \begin{array}{l} \text{ASL\_hcgkaa} \\ \text{ASL\_gcgkaa} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## (7) Example

## (a) Problem

For Hermitian matrix of the degree 4

$$A = \begin{bmatrix} 8 & 3 & 1-2i & -1-2i \\ 3 & 9 & 1+2i & -1+2i \\ 1+2i & 1-2i & 10 & -3 \\ -1+2i & -1-2i & -3 & 11 \end{bmatrix}$$

and its conjugate Hermitian matrix

$$B = \begin{bmatrix} 8 & 3 & 1+2i & -1+2i \\ 3 & 9 & 1-2i & -1-2i \\ 1-2i & 1+2i & 10 & -3 \\ -1-2i & -1+2i & -3 & 11 \end{bmatrix}$$

obtain eigenvector of generalized eigenvalue problem  $ABx = \lambda x$ .

## (b) Input data

n=4, lna=4, matrix  $A$ , lnb=4, nt=2, matrix  $B$ .

## (c) Main program

```

/*      C interface example for ASL_hcgjaa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar, *br, *ai, *bi, *e, *work;
    int i,j,n,ierr;
    int nt;
    int lna, lnb;
    n=4;
    lna=4;lnb=4,nt=4;
    printf( "      *** ASL_hcgjaa ***\n" );
    printf( "\n      ** Input **\n\n" );
    ar = ( double * )malloc(sizeof(double)*n*n);
    if( ar == NULL )
    {
        printf( "no enough memory for array ar\n");
        return -1;
    }
    br = ( double * )malloc(sizeof(double)*n*n);
    if( br == NULL )
    {
        printf( "no enough memory for array br\n");
        return -1;
    }
    ai = ( double * )malloc(sizeof(double)*n*n);
    if( ai == NULL )
    {
        printf( "no enough memory for array ai\n");
        return -1;
    }
    bi = ( double * )malloc(sizeof(double)*n*n);
    if( bi == NULL )

```

```

{
    printf( "no enough memory for array bi\n");
    return -1;
}
e = ( double * )malloc(sizeof(double)*n);
if ( e == NULL )
{
    printf( "no enough memory for array e\n");
    return -1;
}
work = ( double * )malloc(sizeof(double)*4*n);
if ( work == NULL )
{
    printf( "no enough memory for array work\n");
    return -1;
}
printf( "\tn = %6d\n" , n );
printf( "\tlna = %6d\n", lna );
printf( "\tlnb = %6d\n", lnb );
printf( "\tnt = %6d\n", nt );
br[0+n*0]=8.0;bi[0+n*0]=0.0;
br[1+n*1]=9.0;bi[1+n*1]=0.0;
br[2+n*2]=10.0;bi[2+n*2]=0.0;
br[3+n*3]=11.0;bi[3+n*3]=0.0;
br[0+n*1]=3.0;bi[0+n*1]=0.0;
br[0+n*2]=1.0;bi[0+n*2]=2.0;
br[0+n*3]=-1.0;bi[0+n*3]=2.0;
br[1+n*2]=1.0;bi[1+n*2]=-2.0;
br[1+n*3]=-1.0;bi[1+n*3]=-2.0;
br[2+n*3]=-3.0;bi[2+n*3]=0.0;
for(i=1;i<n;i++)
{
    for(j=0;j<i;j++)
    {
        br[i+n*j]= br[j+n*i];
        bi[i+n*j]=-bi[j+n*i];
    }
}
for(i=0;i<n;i++)
{
    for(j=0;j<n;j++)
    {
        ar[i+n*j]= br[i+n*j];
        ai[i+n*j]=-bi[i+n*j];
    }
}
printf( "\n\tInput Matrix a\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g," ,ar[i+n*j]);
        printf( "%8.3g) ",ai[i+n*j]);
    }
    printf( "\n" );
}
printf( "\n\tInput Matrix b\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g," ,br[i+n*j]);
        printf( "%8.3g) ",bi[i+n*j]);
    }
    printf( "\n" );
}

ierr = ASL_hcgjaa(ar,ai, lna, n, br,bi, lnb, e, work, nt);

printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );
printf( "\n\t      Eigenvalue   " );
printf( "\n\t" );
for(j=0; j<n; j++)
{
    printf( "      %8.3g      ",e[j]);
}
printf( "\n\t      Eigenvector   " );
for(i=0; i<n; i++)
{
    printf( "\n\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g," ,ar[i+n*j]);
        printf( "%8.3g) ",ai[i+n*j]);
    }
}

```

```

}
printf( "\n" );
free(ar);
free(br);
free(ai);
free(bi);
free(e);
free(work);
return 0;
}

```

(d) Output results

```
*** ASL_hcgjaa ***
```

```
** Input **
```

```
n   =    4
lna =    4
lnb =    4
nt  =    4
```

```
Input Matrix a
```

```
(      8,      0) (      3,      0) (      1,     -2) (     -1,     -2)
(      3,      0) (      9,      0) (      1,      2) (     -1,      2)
(      1,      2) (      1,     -2) (     10,      0) (     -3,      0)
(     -1,     -2) (     -1,     -2) (     -3,      0) (     11,      0)
```

```
Input Matrix b
```

```
(      8,      0) (      3,      0) (      1,      2) (     -1,      2)
(      3,      0) (      9,      0) (      1,     -2) (     -1,     -2)
(      1,     -2) (      1,      2) (     10,      0) (     -3,      0)
(     -1,     -2) (     -1,      2) (     -3,      0) (     11,      0)
```

```
** Output **
```

```
ierr =    0
```

```

Eigenvalue
16.7
Eigenvector
( -0.399,      0) ( -0.109,      0) (  0.17,      0) (  0.0917,      0)
(  0.345, 0.000936) (  0.102,  0.016) (  0.203,  0.00834) (  0.101, -0.00404)
(  0.00726, -0.132) (  0.0187, -0.327) ( -0.0997, -0.00601) (  0.147, -0.00401)
(-0.00421, -0.125) (  0.0156, -0.281) (  0.13, -0.000644) ( -0.166,  0.00269)

```

### 5.9.2 ASL\_hcgjan, ASL\_gcgjan All Eigenvalues of Hermitian Matrices (Generalized Eigenvalue Problem $ABz = \lambda z$ , $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$ABz = \lambda z$$

( $A$ : Hermitian,  $B$ : Positive Hermitian) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$ .

(2) **Usage**

Double precision:

ierr = ASL\_hcgjan ( ar, ai, lna, n, br, bi, lnb, e, work, nt);

Single precision:

ierr = ASL\_gcgjan ( ar, ai, lna, n, br, bi, lnb, e, work, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	$lna \times n$	Input	Real part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
2	ai	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	$lna \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	$lnb \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	$lnb \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	n	Output	Eigenvalues $\lambda$
9	work	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $1 \leq n \leq lna, lnb$

(b)  $nt \geq 1$



(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

(6) Notes

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) 5.9.1  $\left\{ \begin{array}{l} \text{ASL\_hcgjaa} \\ \text{ASL\_gcgjaa} \end{array} \right\}$  should be used if the eigenvectors are needed.
- (d) 5.10.2  $\left\{ \begin{array}{l} \text{ASL\_hcgkan} \\ \text{ASL\_gcgkan} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

---

## 5.10 GENERALIZED EIGENVALUE PROBLEM ( $BAz = \lambda z$ ) FOR HERMITIAN MATRICES (TWO-DIMENSIONAL ARRAY TYPE) (UPPER TRIANGULAR TYPE) (REAL ARGUMENT TYPE)

### 5.10.1 ASL\_hcgkaa, ASL\_gcgkaa

All Eigenvalues and All Eigenvectors of Hermitian Matrices (Generalized Eigenvalue Problem  $BAz = \lambda z$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$BAz = \lambda z$$

( $A$ : Hermitian,  $B$ : Positive Hermitian) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$  and corresponding all eigenvectors  $z$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hcgkaa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_gcgkaa ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

## (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$
				Output	Real part of eigenvector $x$
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Imaginary part of eigenvector $x$
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues $\lambda$
9	work	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

## (4) Restrictions

(a)  $1 \leq n \leq \text{lna}, \text{lnb}$ (b)  $\text{nt} \geq 1$

## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ and $a[0] \leftarrow \sqrt{b[0]}$ are performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

## (6) Notes

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) Eigenvectors  $v_i$  are an orthonormal set so that  $v_j^* B^{-1} v_k = \delta_{j,k}$
- (d) 5.10.2  $\left\{ \begin{array}{l} \text{ASL\_hcgkan} \\ \text{ASL\_gcgkan} \end{array} \right\}$  should be used if the eigenvectors are not needed.
- (e) 5.9.1  $\left\{ \begin{array}{l} \text{ASL\_hcgjaa} \\ \text{ASL\_gcgjaa} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## (7) Example

## (a) Problem

For Hermitian matrix of the degree 4

$$A = \begin{bmatrix} 8 & 3 & 1-2i & -1-2i \\ 3 & 9 & 1+2i & -1+2i \\ 1+2i & 1-2i & 10 & -3 \\ -1+2i & -1-2i & -3 & 11 \end{bmatrix}$$

and its conjugate Hermitian matrix

$$B = \begin{bmatrix} 8 & 3 & 1+2i & -1+2i \\ 3 & 9 & 1-2i & -1-2i \\ 1-2i & 1+2i & 10 & -3 \\ -1-2i & -1+2i & -3 & 11 \end{bmatrix}$$

obtain eigenvector of generalized eigenvalue problem  $BAx = \lambda x$ .

## (b) Input data

$n=4$ ,  $lna=4$ , matrix  $A$ ,  $lnb=4$ ,  $nt=2$  and matrix  $B$ .

## (c) Main program

```

/*      C interface example for ASL_hcgkaa */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *ar, *br, *ai, *bi, *e, *work;
    int    i,j,n,ierr;
    int    nt;
    int    lna, lnb;

```

```

n=4;
lna=4;lnb=4,nt=4;
printf( "      *** ASL_hcgkaa ***\n" );
printf( "\n      ** Input **\n\n" );
ar = ( double * )malloc(sizeof(double)*n*n);
if( ar == NULL )
{
    printf( "no enough memory for array ar\n");
    return -1;
}
br = ( double * )malloc(sizeof(double)*n*n);
if( br == NULL )
{
    printf( "no enough memory for array br\n");
    return -1;
}
ai = ( double * )malloc(sizeof(double)*n*n);
if( ai == NULL )
{
    printf( "no enough memory for array ai\n");
    return -1;
}
bi = ( double * )malloc(sizeof(double)*n*n);
if( bi == NULL )
{
    printf( "no enough memory for array bi\n");
    return -1;
}
e = ( double * )malloc(sizeof(double)*n);
if ( e == NULL )
{
    printf( "no enough memory for array e\n");
    return -1;
}
work = ( double * )malloc(sizeof(double)*4*n);
if ( work == NULL )
{
    printf( "no enough memory for array work\n");
    return -1;
}
printf( "\tn      = %6d\n" , n );
printf( "\tlna    = %6d\n", lna );
printf( "\tlnb    = %6d\n", lnb );
printf( "\tnt     = %6d\n", nt );
br[0+n*0]=8.0;bi[0+n*0]=0.0;
br[1+n*1]=9.0;bi[1+n*1]=0.0;
br[2+n*2]=10.0;bi[2+n*2]=0.0;
br[3+n*3]=11.0;bi[3+n*3]=0.0;
br[0+n*1]=3.0;bi[0+n*1]=0.0;
br[0+n*2]=1.0;bi[0+n*2]=2.0;
br[0+n*3]=-1.0;bi[0+n*3]=2.0;
br[1+n*2]=1.0;bi[1+n*2]=-2.0;
br[1+n*3]=-1.0;bi[1+n*3]=-2.0;
br[2+n*3]=-3.0;bi[2+n*3]=0.0;
for(i=1;i<n;i++)
{
    for(j=0;j<i;j++)
    {
        br[i+n*j]= br[j+n*i];
        bi[i+n*j]=-bi[j+n*i];
    }
}
for(i=0;i<n;i++)
{
    for(j=0;j<n;j++)
    {
        ar[i+n*j]= br[i+n*j];
        ai[i+n*j]=-bi[i+n*j];
    }
}
printf( "\n\tInput Matrix a\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g,",ar[i+n*j]);
        printf( "%8.3g) ",ai[i+n*j]);
    }
    printf( "\n" );
}
printf( "\n\tInput Matrix b\n\n" );
for(i=0; i<n; i++)
{
    printf( "\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g,",br[i+n*j]);
        printf( "%8.3g) ",bi[i+n*j]);
    }
}

```

```

    printf( "\n" );
}
ierr = ASL_hcgkaa(ar,ai, lna, n, br,bi, lnb, e, work, nt);
printf( "\n  ** Output **\n\n" );
printf( "\t ierr = %6d\n", ierr );
printf( "\n\t      Eigenvalue  " );
printf( "\n\t" );
for(j=0; j<n; j++)
{
    printf( "      %8.3g      ",e[j]);
}
printf( "\n\t      Eigenvector  " );
for(i=0; i<n; i++)
{
    printf( "\n\t" );
    for(j=0; j<n; j++)
    {
        printf( "(%8.3g,",ar[i+n*j]);
        printf( "%8.3g) ",ai[i+n*j]);
    }
}
printf( "\n" );
free(ar);
free(br);
free(ai);
free(bi);
free(e);
free(work);
return 0;
}

```

(d) Output results

```

*** ASL_hcgkaa ***
** Input **
n =      4
lna =     4
lnb =     4
nt =     4

Input Matrix a
(      8,      0) (      3,      0) (      1,     -2) (     -1,     -2)
(      3,      0) (      9,      0) (      1,      2) (     -1,      2)
(      1,      2) (      1,     -2) (     10,      0) (     -3,      0)
(     -1,     -2) (     -1,     -2) (     -3,      0) (     11,      0)

Input Matrix b
(      8,      0) (      3,      0) (      1,      2) (     -1,      2)
(      3,      0) (      9,      0) (      1,     -2) (     -1,     -2)
(      1,     -2) (      1,      2) (     10,      0) (     -3,      0)
(     -1,     -2) (     -1,      2) (     -3,      0) (     11,      0)

** Output **
ierr =      0

      Eigenvalue
      16.7
      Eigenvector
(  -1.63,  0.00149) (  0.656,  0.0713) (  1.76,  0.08) (  1.35, -0.0577)
(  1.41, -0.00511) ( -0.625,  0.0299) (  2.09,  0.00933) (  1.5, -0.00416)
(  0.0301,  0.54) (  0.102, -1.99) ( -1.03,  0.0151) (  2.16, -0.0331)
( -0.0167,  0.51) (  0.0905, -1.71) (  1.34,  0.0676) ( -2.45,  0.0648)

```

### 5.10.2 ASL\_hcgkan, ASL\_gcgkan

All Eigenvalues of Hermitian Matrices (Generalized Eigenvalue Problem  
 $BAz = \lambda z$ ,  $B$ : Positive)

(1) **Function**

Generalized eigenvalue problem

$$BAz = \lambda z$$

( $A$ : Hermitian,  $B$ : Positive Hermitian) is solved by using the Cholesky method, the Householder method and QR method to obtain all eigenvalues  $\lambda$ .

(2) **Usage**

Double precision:

```
ierr = ASL_hcgkan ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

Single precision:

```
ierr = ASL_gcgkan ( ar, ai, lna, n, br, bi, lnb, e, work, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	ar	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Real part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
2	ai	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lna} \times n$	Input	Imaginary part of Hermitian matrix $A$
				Output	Input-time contents are not retained.
3	lna	I	1	Input	Adjustable dimension of arrays ar and ai
4	n	I	1	Input	Order of matrices $A$ and $B$
5	br	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Real part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
6	bi	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$\text{lnb} \times n$	Input	Imaginary part of Hermitian matrix $B$
				Output	Input-time contents are not retained.
7	lnb	I	1	Input	Adjustable dimension of arrays br and bi
8	e	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	n	Output	Eigenvalues $\lambda$
9	work	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$4 \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $1 \leq n \leq \text{lna}, \text{lnb}$
- (b)  $\text{nt} \geq 1$



## (5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
1000	n was equal to 1.	$e[0] \leftarrow a[0]*b[0]$ is performed.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
4000	$B$ was not positive definite.	
5000	The sequence did not converge in the step where the eigenvalue was obtained.	

## (6) Notes

- (a) Arrays ar, ai, br and bi should be stored only in the upper triangular portions.
- (b) Eigenvalues are stored in ascending order.
- (c) 5.10.1  $\left\{ \begin{array}{l} \text{ASL\_hcgkaa} \\ \text{ASL\_gcgkaa} \end{array} \right\}$  should be used if the eigenvectors are needed.
- (d) 5.9.2  $\left\{ \begin{array}{l} \text{ASL\_hcgjan} \\ \text{ASL\_gcgjan} \end{array} \right\}$  should be used if matrix  $A$  is only positive.

## Chapter 6

---

# FOURIER TRANSFORMS AND THEIR APPLICATIONS

## 6.1 INTRODUCTION

This chapter describes functions that perform fast Fourier transforms, convolutions, correlations, power spectrum analysis.

**Function described in this chapter divides up and allocates internal processing among threads and executes allocated processing in parallel.**

The following functions are provided for computing the discrete Fourier transform according to the data characteristics. You can perform processing efficiently if your data satisfies any of the characteristics described below.

- (1) Multiple one-dimensional complex Fourier transform  
Data values are complex numbers and are multiple one-dimensional data.
- (2) Multiple one-dimensional real Fourier transform  
Data values are real numbers and are multiple one-dimensional data.
- (3) Two-dimensional complex Fourier transform  
Data values are complex numbers and are two-dimensional data.
- (4) Two-dimensional real Fourier transform  
Data values are real numbers and are two-dimensional data.
- (5) Three-dimensional complex Fourier transform  
Data values are complex numbers and are three-dimensional data.
- (6) Three-dimensional real Fourier transform  
Data values are real numbers and are three-dimensional data.

In addition, although the number of equal divisions  $n$  of the input data can be transformed by the fast Fourier transforms handled in this chapter no matter what prime number is used as radix, calculation efficiency decreases for a sequence formed from a large prime number. Therefore, the number of equal divisions  $n$  should be able to be factored into small radices (2, 3, 5, 7).

The functions, which handles the following data, are provided for performing the convolutions, correlations, power spectrum analysis.

- (1) Two-dimensional data
- (2) Three-dimensional data

### 6.1.1 Notes

- (1) For a two-dimensional complex Fourier transform or two-dimensional real Fourier transform, the number of sample points  $n_x$ ,  $n_y$  ( $n_x$ ,  $n_y$ ,  $n_z$  for a three-dimensional complex Fourier transform or three-dimensional real Fourier transform) corresponds to a sample on a single period  $(0, 2\pi)$ .
- (2) You first must use a transform function that includes initialization. This function performs such tasks as creating a trigonometric function table. If you plan to continue computing Fourier transforms for the same number of data  $N$ , processing will be more efficient if you use the post-initialization transform function thereafter since the contents of work areas are retained.

## 6.1.2 Algorithms Used

### 6.1.2.1 Two-dimensional complex Fourier transform

The complex Fourier forward transform  $C_{j_x, j_y}$  for one period  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) of complex multiperiodic data  $\hat{c}_{k_x, k_y}$  satisfying  $\hat{c}_{k_x, k_y} = \hat{c}_{k_x + n_x, k_y + n_y}$  for arbitrary integers  $k_x$  and  $k_y$  is defined by the following equation.

$$C_{j_x, j_y} = \frac{1}{\alpha} \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)}$$

$$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

$\alpha$  is an arbitrary constant for which 1 or  $n_x n_y$  normally is selected. At this time, the complex data after the transform  $C_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) also corresponds to one period of complex multiperiodic data  $\hat{C}_{j_x, j_y}$  satisfying  $\hat{C}_{j_x, j_y} = \hat{C}_{j_x + n_x, j_y + n_y}$  for an arbitrary integers  $j_x$  and  $j_y$ . The corresponding backward transform is as follows:

$$c_{k_x, k_y} = \frac{1}{n_x n_y} \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} (\alpha C_{j_x, j_y}) e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

The functions in this manual obtain  $\alpha C_{j_x, j_y}$  and  $n_x n_y c_{k_x, k_y}$  except for a constant factor, from the normal Fourier transform definition.

### 6.1.2.2 Two-dimensional real Fourier transform

When the data for which the two-dimensional Fourier forward transform is to be calculated is real, the following relationship is satisfied.

$$\begin{aligned} \alpha C_{n_x - j_x, n_y - j_y}^* &= \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y}^* e^{2\pi\sqrt{-1}\left\{\frac{(n_x - j_x)k_x}{n_x} + \frac{(n_y - j_y)k_y}{n_y}\right\}} \\ &= \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left\{\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right\}} \\ &= \alpha C_{j_x, j_y} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ . In particular,  $C_{0,0}$  is real and when  $n_x$  and  $n_y$  are even, then  $C_{\frac{n_x}{2}, \frac{n_y}{2}}$  is also real.

Similarly, the following relationship is satisfied.

$$\begin{aligned} \alpha C_{n_x - j_x, j_y}^* &= \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y}^* e^{2\pi\sqrt{-1}\left\{\frac{(n_x - j_x)k_x}{n_x} + \frac{j_y k_y}{n_y}\right\}} \\ &= \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left\{\frac{j_x k_x}{n_x} + \frac{(n_y - j_y)k_y}{n_y}\right\}} \\ &= \alpha C_{j_x, n_y - j_y} \end{aligned}$$

Therefore, the calculations for a Fourier transform can be executed using half the data for the case of general complex data (for  $c_{k_x, k_y}$ , only the real part, and for  $C_{j_x, j_y}$ , half of its period for either  $j_x$  or  $j_y$ ).

### 6.1.2.3 Three-dimensional complex Fourier transform

The complex Fourier forward transform  $C_{j_x, j_y, j_z}$  for one period  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) of complex multiperiodic data  $\hat{c}_{k_x, k_y, k_z}$  satisfying  $\hat{c}_{k_x, k_y, k_z} = \hat{c}_{k_x+n_x, k_y+n_y, k_z+n_z}$  for arbitrary integers  $k_x, k_y$  and  $k_z$  is defined by the following equation.

$$C_{j_x, j_y, j_z} = \frac{1}{\alpha} \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$$

$\alpha$  is an arbitrary constant for which 1 or  $n_x n_y n_z$  normally is selected. At this time, the complex data after the transform  $C_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) also corresponds to one period of complex multiperiodic data  $\hat{C}_{j_x, j_y, j_z}$  satisfying  $\hat{C}_{j_x, j_y, j_z} = \hat{C}_{j_x+n_x, j_y+n_y, j_z+n_z}$  for an arbitrary integers  $j_x, j_y$  and  $j_z$ . The corresponding backward transform is as follows:

$$c_{k_x, k_y, k_z} = \frac{1}{n_x n_y n_z} \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \sum_{j_z=0}^{n_z-1} (\alpha C_{j_x, j_y, j_z}) e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

The functions in this manual obtain  $\alpha C_{j_x, j_y, j_z}$  and  $n_x n_y n_z c_{k_x, k_y, k_z}$  except for a constant factor, from the normal Fourier transform definition.

### 6.1.2.4 Three-dimensional real Fourier transform

When the data for which the three-dimensional Fourier forward transform is to be calculated is real, the following relationship is satisfied.

$$\alpha C_{n_x-j_x, n_y-j_y, n_z-j_z}^* = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} C_{k_x, k_y, k_z}^* e^{2\pi\sqrt{-1}\left\{\frac{(n_x-j_x)k_x}{n_x} + \frac{(n_y-j_y)k_y}{n_y} + \frac{(n_z-j_z)k_z}{n_z}\right\}}$$

$$= \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} C_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left\{\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right\}}$$

$$= \alpha C_{j_x, j_y, j_z}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ . In particular,  $C_{0,0,0}$  is real and when  $n_x, n_y$  and  $n_z$  are all even, then  $C_{\frac{n_x}{2}, \frac{n_y}{2}, \frac{n_z}{2}}$  is also real.

Similarly, the following kinds of relationships are satisfied.

$$C_{n_x-j_x, j_y, j_z}^* = C_{j_x, n_y-j_y, n_z-j_z}$$

That is, the complex number after substitutions are made in the subscripts using the following correspondence relationships for all of  $j_x, j_y$  and  $j_z$  is mutually related as a complex conjugate to the complex number before the substitutions are made.

$$j_x \leftrightarrow n_x - j_x$$

$$j_y \leftrightarrow n_y - j_y$$

$$j_z \leftrightarrow n_z - j_z$$

Therefore, the calculations for a Fourier transform can be executed using half the data for the case of general complex data (for  $c_{k_x, k_y, k_z}$ , only the real part, and for  $C_{j_x, j_y, j_z}$ , half of its period for either  $j_x, j_y$ , or  $j_z$ ).

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## 6.2 MULTIPLE ONE-DIMENSIONAL COMPLEX FOURIER TRANSFORM (REAL ARGUMENT TYPE)

### 6.2.1 [DEPRECATED]ASL\_qfcmfb, ASL\_pfcmfb Multiple One-Dimensional Complex Fourier Transforms (Include Initialization)

#### (1) Function

##### Forward transform

ASL\_qfcmfb or ASL\_pfcmfb computes the  $m$ -fold one-dimensional complex Fourier forward transform (arbitrary radix) for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ).

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

##### Backward transform

ASL\_qfcmfb or ASL\_pfcmfb computes the  $m$ -fold one-dimensional complex Fourier backward transform (arbitrary radix) for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ).

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

#### (2) Usage

Double precision:

ierr = ASL\_qfcmfb (n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfcmfb (n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	cr	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Input	Real part of input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Real part of output data $d_{j,l}$ (See Notes (b) and (c))
4	ci	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Input	Imaginary part of input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Imaginary part of output data $d_{j,l}$ (See Notes (b) and (c))
5	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
6	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
7	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
8	ifax	I*	20	Output	Factorization results and number of factors (See Note (d))
9	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times n$	Output	Trigonometric function table (See Note (d))
10	wk	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times m \times n$	Work	Work area
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)



(4) **Restrictions**

- (a)  $n \geq 1, m \geq 1$
- (b)  $incn \geq 1, incm \geq 1$
- (c)  $incn \geq m \times gcm(incn, incm)$  or  
 $incm \geq n \times gcm(incn, incm)$   
 where,  $gcm(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n=1$ .	Input data is output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) When the number of transformed data  $n$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n = 289 (= 17^2)$ , it is usually more efficient to set  $n = 300 (= 2^2 \times 3 \times 5^2)$ ,  $n = 320 (= 2^6 \times 5)$ ,  $n = 384 (= 2^7 \times 3)$  or the like.
- (b) If we let the real and imaginary parts of the complex data  $c_{k,l}$  ( $k = 0, \dots, n-1; l = 1, \dots, m$ ) be  $\Re\{c_{k,l}\}$  and  $\Im\{c_{k,l}\}$ , respectively, the  $c_{k,l}$  and elements of arrays  $cr$  and  $ci$  are associated as follows.

$$\begin{aligned} \Re\{c_{k,l}\} &\leftrightarrow cr[incn * k + incm * (l - 1)] \\ \Im\{c_{k,l}\} &\leftrightarrow ci[incn * k + incm * (l - 1)] \end{aligned}$$

For example, if we let  $incn=1$  and  $incm=n$ , then the associations are as follows:

$$\Re\{c_{k,l}\} \leftrightarrow cr[k + n * (l - 1)], \quad \Im\{c_{k,l}\} \leftrightarrow ci[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $incn=m$  and  $incm=1$ , then the associations are as follows:

$$\Re\{c_{k,l}\} \leftrightarrow cr[(l - 1) + m * k], \quad \Im\{c_{k,l}\} \leftrightarrow ci[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . Similarly, for the complex data  $d_{j,l}$  ( $j = 0, \dots, n-1; l = 1, \dots, m$ ). Values in areas where the data of arrays  $cr$  and  $ci$  is not stored do not change when this function is called.

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform

immediately following the forward transform for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ) be represented by  $\hat{c}_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ), then the following relationship holds.

$$\hat{c}_{k,l} = nc_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data  $n$ , you should call this function once, and then use the after-initialization transform 6.2.2  $\left\{ \begin{array}{l} \text{ASL\_qfcmfb} \\ \text{ASL\_pfcmbf} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays `ifax` and `trigs` so they can be used as input to the function 6.2.2  $\left\{ \begin{array}{l} \text{ASL\_qfcmfb} \\ \text{ASL\_pfcmbf} \end{array} \right\}$ .

To perform initialization only by setting `isw=0`, you need not set input data for arrays `cr` and `ci`.

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.2.2 (7).

## 6.2.2 [DEPRECATED]ASL\_qfcmbf, ASL\_pfcmbf Multiple One-Dimensional Complex Fourier Transforms (After Initialization)

### (1) Function

#### Forward transform

ASL\_qfcmbf or ASL\_pfcmbf computes the  $m$ -fold one-dimensional complex Fourier forward transform (arbitrary radix) for the complex data  $c_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ .

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

#### Backward transform

ASL\_qfcmbf or ASL\_pfcmbf computes the  $m$ -fold one-dimensional complex Fourier backward transform (arbitrary radix) for the complex data  $c_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ .

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

### (2) Usage

Double precision:

ierr = ASL\_qfcmbf (n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfcmbf (n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	cr	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	See Contents	Input	Real part of input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Real part of output data $d_{j,l}$ (See Notes (b) and (c))
4	ci	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	See Contents	Input	Imaginary part of input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Imaginary part of output data $d_{j,l}$ (See Notes (b) and (c))
5	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
6	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
7	isw	I	1	Input	Processing switch isw= 1:Forward transform isw=-1:Backward transform
8	ifax	I*	20	Input	Factorization results and number of factors (See Note (a))
9	trigs	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times n$	Input	Trigonometric function table (See Note (a))
10	wk	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times m \times n$	Work	Work area
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $n \geq 1, m \geq 1$
- (b)  $\text{incn} \geq 1, \text{incm} \geq 1$
- (c)  $\text{incn} \geq m \times \text{gcm}(\text{incn}, \text{incm})$  or  
 $\text{incm} \geq n \times \text{gcm}(\text{incn}, \text{incm})$   
 where,  $\text{gcm}(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $\text{isw} \in \{1, -1\}$
- (e)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n=1$ .	Input data is output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) This function can be used to repeatedly compute the transform for the same number of transformed data  $n$  after the including-initialization function 6.2.1  $\left\{ \begin{matrix} \text{ASL\_qfcmfb} \\ \text{ASL\_pfcmbf} \end{matrix} \right\}$  has been used. In this case, you must retain the contents of arrays  $\text{ifax}$  and  $\text{trigs}$  so they can be used as input in this function.
- (b) If we let the real and imaginary parts of the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ) be  $\Re\{c_{k,l}\}$  and  $\Im\{c_{k,l}\}$ , respectively, the  $c_{k,l}$  and elements of arrays  $\text{cr}$  and  $\text{ci}$  are associated as follows.

$$\begin{aligned} \Re\{c_{k,l}\} &\leftrightarrow \text{cr}[\text{incn} * k + \text{incm} * (l - 1)] \\ \Im\{c_{k,l}\} &\leftrightarrow \text{ci}[\text{incn} * k + \text{incm} * (l - 1)] \end{aligned}$$

For example, if we let  $\text{incn}=1$  and  $\text{incm}=n$ , then the associations are as follows:

$$\Re\{c_{k,l}\} \leftrightarrow \text{cr}[k + n * (l - 1)], \quad \Im\{c_{k,l}\} \leftrightarrow \text{ci}[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $\text{incn}=m$  and  $\text{incm}=1$ , then the associations are as follows:

$$\Re\{c_{k,l}\} \leftrightarrow \text{cr}[(l - 1) + m * k], \quad \Im\{c_{k,l}\} \leftrightarrow \text{ci}[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . Similarly, for the complex data  $d_{j,l}$  ( $j = 0, \dots, n - 1; l = 1, \dots, m$ ). Values in areas where the data of arrays  $\text{cr}$  and  $\text{ci}$  is not stored do not change when this function is called.

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform

immediately following the forward transform for the complex data  $c_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$  be represented by  $\hat{c}_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ , then the following relationship holds.

$$\hat{c}_{k,l} = nc_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

- (a) Problem

Compute the multiple one-dimensional complex Fourier transform using the following sequence of numbers as input data.

```

cr[ 0]= 1.000   ci[ 0]=4.000
cr[ 1]= 2.000   ci[ 1]=3.000
cr[ 2]= 3.000   ci[ 2]=2.000
cr[ 3]= 4.000   ci[ 3]=1.000
cr[ 4]= 4.000   ci[ 4]=1.000
cr[ 5]= 3.000   ci[ 5]=2.000
cr[ 6]= 2.000   ci[ 6]=3.000
cr[ 7]= 1.000   ci[ 7]=4.000
cr[ 9]= 1.000   ci[ 9]=2.000
cr[10]= 1.000   ci[10]=2.000
cr[11]= 2.000   ci[11]=1.000
cr[12]= 2.000   ci[12]=1.000
cr[13]= 2.000   ci[13]=1.000
cr[14]= 2.000   ci[14]=1.000
cr[15]= 1.000   ci[15]=2.000
cr[16]= 1.000   ci[16]=2.000
cr[18]= 1.000   ci[18]=2.000
cr[19]= 1.000   ci[19]=2.000
cr[20]= 1.000   ci[20]=2.000
cr[21]= 1.000   ci[21]=2.000
cr[22]= 2.000   ci[22]=1.000
cr[23]= 2.000   ci[23]=1.000
    
```

```

cr[24]= 2.000   ci[24]=1.000
cr[25]= 2.000   ci[25]=1.000
cr[27]= 1.000   ci[27]=1.000
cr[28]= 1.000   ci[28]=1.000
cr[29]= 1.000   ci[29]=1.000
cr[30]= 1.000   ci[30]=1.000
cr[31]= 1.000   ci[31]=1.000
cr[32]= 1.000   ci[32]=1.000
cr[33]= 1.000   ci[33]=1.000
cr[34]= 1.000   ci[34]=1.000
    
```

(b) Input data

Array cr and ci, n=8, m=4, incn=1, incm=9, isw=1 (forward transform), isw=-1 (backward transform) and nt=2.

(c) Main program

```

/*      C interface example for ASL_qfcmbf , ASL_qfcmbf */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int ld=35;
    int n;      int m;
    double *cr; double *ci;
    int incn;   int incm;
    int isw;
    int ifax[20];   double *trigs;
    double *wk;
    int nt;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "qfcmbf.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_qfcmbf , ASL_qfcmbf ***\n" );
    printf( "\n      ** Input **\n\n" );

    cr = ( double * )malloc((size_t)( sizeof(double) * ld ));
    if( cr == NULL )
    {
        printf( "no enough memory for array cr\n" );
        return -1;
    }

    ci = ( double * )malloc((size_t)( sizeof(double) * ld ));
    if( ci == NULL )
    {
        printf( "no enough memory for array ci\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*ld) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double * )malloc((size_t)( sizeof(double) * (2*ld) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    fscanf( fp, "%d,%d,%d,%d,%d", &n, &m, &incn, &incm, &nt );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
    
```

```

        {
            fscanf( fp, "%lf,%lf", &cr[i*incn+j*incm], &ci[i*incn+j*incm] );
        }
    }

    printf("\t  n   = %d \n", n );
    printf("\t  m   = %d \n", m );
    printf("\t  incn = %d \n", incn );
    printf("\t  incm = %d \n", incm );
    printf("\t  nt   = %d \n\n", nt );

    printf( "\t Real Part                Imaginary Part\n" );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            printf( "\t cr[%3d] = %8.3g          ci[%3d] = %8.3g\n",
                i*incn+j*incm, cr[i*incn+j*incm],
                i*incn+j*incm, ci[i*incn+j*incm] );
        }
    }

    fclose( fp );

    printf( "\n    ** Output **\n" );

    isw = 1;
    ierr = ASL_qfcmbf(n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            cr[i*incn+j*incm] /= n ;
            ci[i*incn+j*incm] /= n ;
        }
    }

    printf( "\n\t< Forward Transform >\n" );
    printf( "\t\tierr = %6d\n", ierr );

    printf( "\n\tSolution\n\n" );
    printf( "\t Real Part                Imaginary Part\n" );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            printf( "\t cr[%3d] = %8.3g          ci[%3d] = %8.3g\n",
                i*incn+j*incm, cr[i*incn+j*incm],
                i*incn+j*incm, ci[i*incn+j*incm] );
        }
    }

    isw = -1;
    ierr = ASL_pfcmbf(n, m, cr, ci, incn, incm, isw, ifax, trigs, wk, nt);

    printf( "\n\t< Backward Transform >\n" );
    printf( "\t\tierr = %6d\n", ierr );

    printf( "\n\tSolution\n\n" );
    printf( "\t Real Part                Imaginary Part\n" );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            printf( "\t cr[%3d] = %8.3g          ci[%3d] = %8.3g\n",
                i*incn+j*incm, cr[i*incn+j*incm],
                i*incn+j*incm, ci[i*incn+j*incm] );
        }
    }

    free( cr );
    free( ci );
    free( trigs );
    free( wk );

    return 0;
}

```

(d) Output results

```

*** ASL_qfcmbf , ASL_pfcmbf ***
** Input **

```



```

n      = 8
m      = 4
incn   = 1
incm   = 9
nt     = 2

Real Part
cr[ 0] =      1
cr[ 1] =      2
cr[ 2] =      3
cr[ 3] =      4
cr[ 4] =      4
cr[ 5] =      3
cr[ 6] =      2
cr[ 7] =      1
cr[ 9] =      1
cr[10] =      1
cr[11] =      2
cr[12] =      2
cr[13] =      2
cr[14] =      2
cr[15] =      1
cr[16] =      1
cr[18] =      1
cr[19] =      1
cr[20] =      1
cr[21] =      1
cr[22] =      2
cr[23] =      2
cr[24] =      2
cr[25] =      2
cr[27] =      1
cr[28] =      1
cr[29] =      1
cr[30] =      1
cr[31] =      1
cr[32] =      1
cr[33] =      1
cr[34] =      1

Imaginary Part
ci[ 0] =      4
ci[ 1] =      3
ci[ 2] =      2
ci[ 3] =      1
ci[ 4] =      1
ci[ 5] =      2
ci[ 6] =      3
ci[ 7] =      4
ci[ 9] =      2
ci[10] =      2
ci[11] =      1
ci[12] =      1
ci[13] =      1
ci[14] =      1
ci[15] =      2
ci[16] =      2
ci[18] =      2
ci[19] =      2
ci[20] =      2
ci[21] =      2
ci[22] =      1
ci[23] =      1
ci[24] =      1
ci[25] =      1
ci[27] =      1
ci[28] =      1
ci[29] =      1
ci[30] =      1
ci[31] =      1
ci[32] =      1
ci[33] =      1
ci[34] =      1
    
```

\*\* Output \*\*

< Forward Transform >  
 ierr = 0

Solution

```

Real Part
cr[ 0] =      2.5
cr[ 1] =     -1.03
cr[ 2] =      0
cr[ 3] =    -0.0732
cr[ 4] =      0
cr[ 5] =     0.0303
cr[ 6] =      0
cr[ 7] =    -0.427
cr[ 9] =      1.5
cr[10] =    -0.427
cr[11] =      0
cr[12] =     0.177
cr[13] =      0
cr[14] =   -0.0732
cr[15] =      0
cr[16] =    -0.177
cr[18] =      1.5
cr[19] =     0.177
cr[20] =      0
cr[21] =   -0.0732
cr[22] =      0
cr[23] =    -0.177
cr[24] =      0
cr[25] =   -0.427
cr[27] =      1
cr[28] =      0
cr[29] =      0
cr[30] =      0
cr[31] =      0
cr[32] =      0
cr[33] =      0
cr[34] =      0

Imaginary Part
ci[ 0] =      2.5
ci[ 1] =     0.427
ci[ 2] =      0
ci[ 3] =   -0.0303
ci[ 4] =      0
ci[ 5] =     0.0732
ci[ 6] =      0
ci[ 7] =      1.03
ci[ 9] =      1.5
ci[10] =     0.177
ci[11] =      0
ci[12] =     0.0732
ci[13] =      0
ci[14] =   -0.177
ci[15] =      0
ci[16] =     0.427
ci[18] =      1.5
ci[19] =     0.427
ci[20] =      0
ci[21] =     0.177
ci[22] =      0
ci[23] =     0.0732
ci[24] =      0
ci[25] =   -0.177
ci[27] =      1
ci[28] =      0
ci[29] =      0
ci[30] =      0
ci[31] =      0
ci[32] =      0
ci[33] =      0
ci[34] =      0
    
```

< Backward Transform >  
 ierr = 0

Solution

```

Real Part
cr[ 0] =      1
cr[ 1] =      2
cr[ 2] =      3
cr[ 3] =      4
cr[ 4] =      4
cr[ 5] =      3
cr[ 6] =      2

Imaginary Part
ci[ 0] =      4
ci[ 1] =      3
ci[ 2] =      2
ci[ 3] =      1
ci[ 4] =      1
ci[ 5] =      2
ci[ 6] =      3
    
```

```
cr [ 7] = 1      ci [ 7] = 4
cr [ 9] = 1      ci [ 9] = 2
cr [10] = 1      ci [10] = 2
cr [11] = 2      ci [11] = 1
cr [12] = 2      ci [12] = 1
cr [13] = 2      ci [13] = 1
cr [14] = 2      ci [14] = 1
cr [15] = 1      ci [15] = 2
cr [16] = 1      ci [16] = 2
cr [18] = 1      ci [18] = 2
cr [19] = 1      ci [19] = 2
cr [20] = 1      ci [20] = 2
cr [21] = 1      ci [21] = 2
cr [22] = 2      ci [22] = 1
cr [23] = 2      ci [23] = 1
cr [24] = 2      ci [24] = 1
cr [25] = 2      ci [25] = 1
cr [27] = 1      ci [27] = 1
cr [28] = 1      ci [28] = 1
cr [29] = 1      ci [29] = 1
cr [30] = 1      ci [30] = 1
cr [31] = 1      ci [31] = 1
cr [32] = 1      ci [32] = 1
cr [33] = 1      ci [33] = 1
cr [34] = 1      ci [34] = 1
```

---

## 6.3 MULTIPLE ONE-DIMENSIONAL COMPLEX FOURIER TRANSFORM (COMPLEX ARGUMENT TYPE)

### 6.3.1 [DEPRECATED]ASL\_hfcmfb, ASL\_gfcmfb

Multiple One-Dimensional Complex Fourier Transforms (Include Initialization)

(1) **Function**

**Forward transform**

ASL\_hfcmfb or ASL\_gfcmfb computes the  $m$ -fold one-dimensional complex Fourier forward transform (arbitrary radix) for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ).

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

**Backward transform**

ASL\_hfcmfb or ASL\_gfcmfb computes the  $m$ -fold one-dimensional complex Fourier backward transform (arbitrary radix) for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ).

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

(2) **Usage**

Double precision:

ierr = ASL\_hfcmfb (n, m, c, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfcmfb (n, m, c, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	c	$\begin{cases} Z^* \\ C^* \end{cases}$	See Contents	Input	Input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Output data $d_{j,l}$ (See Notes (b) and (c))
4	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
5	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
6	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	20	Output	Factorization results and number of factors (See Note (d))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times n$	Output	Trigonometric function table (See Note (d))
9	wk	$\begin{cases} Z^* \\ C^* \end{cases}$	$m \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n \geq 1, m \geq 1$
- (b)  $\text{incn} \geq 1, \text{incm} \geq 1$
- (c)  $\text{incn} \geq m \times \text{gcm}(\text{incn}, \text{incm})$  or  
 $\text{incm} \geq n \times \text{gcm}(\text{incn}, \text{incm})$   
where,  $\text{gcm}(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $\text{isw} \in \{0, 1, -1\}$
- (e)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	n=1.	Input data is output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) When the number of transformed data  $n$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n = 289 (= 17^2)$ , it is usually more efficient to set  $n = 300 (= 2^2 \times 3 \times 5^2)$ ,  $n = 320 (= 2^6 \times 5)$ ,  $n = 384 (= 2^7 \times 3)$  or the like.

- (b) The complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ) and elements of array  $c$  are associated as follows.

$$c_{k,l} \leftrightarrow c[\text{incn} * k + \text{incm} * (l - 1)]$$

For example, if we let  $\text{incn}=1$  and  $\text{incm}=n$ , then the associations are as follows:

$$c_{k,l} \leftrightarrow c[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $\text{incn}=m$  and  $\text{incm}=1$ , then the associations are as follows:

$$c_{k,l} \leftrightarrow c[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . Similarly, for the complex data  $d_{j,l}$  ( $j = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ). Values in areas where the data of array  $c$  is not stored do not change when this function is called.

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ) be represented by  $\hat{c}_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ), then the following relationship holds.

$$\hat{c}_{k,l} = nc_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data  $n$ , you should call this function once, and then use the after-initialization transform 6.3.2  $\left\{ \begin{array}{l} \text{ASL\_hfcmfb} \\ \text{ASL\_gfcmfb} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays  $\text{ifax}$  and

trigs so they can be used as input to the function 6.3.2  $\left\{ \begin{array}{l} \text{ASL\_hfcmbf} \\ \text{ASL\_gfcmbf} \end{array} \right\}$ .

To perform initialization only by setting `isw=0`, you need not set input data for array `c`.

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.3.2 (7).

### 6.3.2 [DEPRECATED]ASL\_hfcmbf, ASL\_gfcmbf Multiple One-Dimensional Complex Fourier Transforms (After Initialization)

(1) **Function**

**Forward transform**

ASL\_hfcmbf or ASL\_gfcmbf computes the  $m$ -fold one-dimensional complex Fourier forward transform (arbitrary radix) for the complex data  $c_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ .

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

**Backward transform**

ASL\_hfcmbf or ASL\_gfcmbf computes the  $m$ -fold one-dimensional complex Fourier backward transform (arbitrary radix) for the complex data  $c_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ .

$$d_{j,l} = \sum_{k=0}^{n-1} c_{k,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, n - 1; l = 1, \dots, m)$$

(2) **Usage**

Double precision:

ierr = ASL\_hfcmbf (n, m, c, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfcmbf (n, m, c, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	c	$\begin{cases} Z^* \\ C^* \end{cases}$	See Contents	Input	Input data $c_{k,l}$ (See Note (b)) <b>Size:</b> $\text{incn} \times (n - 1) + \text{incm} \times (m - 1) + 1$
				Output	Output data $d_{j,l}$ (See Notes (b) and (c))
4	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
5	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
6	isw	I	1	Input	Processing switch isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	20	Input	Factorization results and number of factors (See Note (a))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times n$	Input	Trigonometric function table (See Note (a))
9	wk	$\begin{cases} Z^* \\ C^* \end{cases}$	$m \times n$	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n \geq 1, m \geq 1$
- (b)  $\text{incn} \geq 1, \text{incm} \geq 1$
- (c)  $\text{incn} \geq m \times \text{gcm}(\text{incn}, \text{incm})$  or  
 $\text{incm} \geq n \times \text{gcm}(\text{incn}, \text{incm})$   
where,  $\text{gcm}(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $\text{isw} \in \{1, -1\}$
- (e)  $\text{nt} \geq 1$



(5) Error indicator (Return Value)

ier value	Meaning	Processing
0	Normal termination.	
1000	n=1.	Input data is output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) Notes

- (a) This function can be used to repeatedly compute the transform for the same number of transformed data  $n$  after the including-initialization function 6.3.1  $\left\{ \begin{matrix} \text{ASL\_hfcmbf} \\ \text{ASL\_gfcmbf} \end{matrix} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.

- (b) The complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ) and elements of array  $c$  are associated as follows.

$$c_{k,l} \leftrightarrow c[\text{incn} * k + \text{incm} * (l - 1)]$$

For example, if we let  $\text{incn}=1$  and  $\text{incm}=n$ , then the associations are as follows:

$$c_{k,l} \leftrightarrow c[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $\text{incn}=m$  and  $\text{incm}=1$ , then the associations are as follows:

$$c_{k,l} \leftrightarrow c[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . Similarly, for the complex data  $d_{j,l}$  ( $j = 0, \dots, n - 1; l = 1, \dots, m$ ). Values in areas where the data of array  $c$  is not stored do not change when this function is called.

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ) be represented by  $\hat{c}_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ), then the following relationship holds.

$$\hat{c}_{k,l} = nc_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that

is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

- (a) Problem

Compute the multiple one-dimensional complex Fourier transform using the following sequence of numbers as input data.

c[ 0]= (1.000, 4.000)

c[ 1]= (2.000, 3.000)

c[ 2]= (3.000, 2.000)

c[ 3]= (4.000, 1.000)

c[ 4]= (4.000, 1.000)

c[ 5]= (3.000, 2.000)

c[ 6]= (2.000, 3.000)

c[ 7]= (1.000, 4.000)

c[ 9]= (1.000, 2.000)

c[10]= (1.000, 2.000)

c[11]= (2.000, 1.000)

c[12]= (2.000, 1.000)

c[13]= (2.000, 1.000)

c[14]= (2.000, 1.000)

c[15]= (1.000, 2.000)

c[16]= (1.000, 2.000)

c[18]= (1.000, 2.000)

c[19]= (1.000, 2.000)

c[20]= (1.000, 2.000)

c[21]= (1.000, 2.000)

c[22]= (2.000, 1.000)

c[23]= (2.000, 1.000)

c[24]= (2.000, 1.000)

c[25]= (2.000, 1.000)

c[27]= (1.000, 1.000)

c[28]= (1.000, 1.000)

c[29]= (1.000, 1.000)

c[30]= (1.000, 1.000)

c[31]= (1.000, 1.000)

c[32]= (1.000, 1.000)

c[33]= (1.000, 1.000)

c[34]= (1.000, 1.000)

- (b) Input data

Array c, n=8, m=4, incn=1, incm=9, isw=1 (forward transform), isw=-1 (backward transform) and

nt=2.

(c) Main program

```

/*      C interface example for ASL_hfcmbf , ASL_hfcmbf */
#include <stdio.h>
#include <stdlib.h>
#include <complex.h>
#include <asl.h>

int main()
{
    int ld=35;
    int n;      int m;
    double _Complex *c;
    int incn;   int incm;
    int isw;
    int ifax[20]; double *trigs;
    double _Complex *wk;
    int nt;
    int ierr;
    int i,j;
    FILE *fp;

    fp = fopen( "hfcmbf.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    printf( "      *** ASL_hfcmbf , ASL_hfcmbf ***\n" );
    printf( "\n      ** Input **\n\n" );

    c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * ld ));
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*ld) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * ld ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    fscanf( fp, "%d,%d,%d,%d,%d", &n, &m, &incn, &incm, &nt );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            double tmp_re, tmp_im;
            fscanf( fp, "%lf,%lf", &tmp_re, &tmp_im );
            c[i*incn+j*incm] = tmp_re + tmp_im * _Complex_I;
        }
    }

    printf("\t  n   = %d \n", n );
    printf("\t  m   = %d \n", m );
    printf("\t  incn = %d \n", incn );
    printf("\t  incm = %d \n", incm );
    printf("\t  nt  = %d \n\n", nt );

    printf( "\t Real Part                                Imaginary Part\n" );
    for( j=0 ; j<m ; j++ )
    {
        for( i=0 ; i<n ; i++ )
        {
            printf( "\t  creal(c[%3d]) = %8.3g          cimag(c[%3d]) = %8.3g\n",
                i*incn+j*incm, creal(c[i*incn+j*incm]),
                i*incn+j*incm, cimag(c[i*incn+j*incm]) );
        }
    }

    fclose( fp );

    printf( "\n      ** Output **\n" );

```

```

isw = 1;
ierr = ASL_hfcmbf(n, m, c, incn, incm, isw, ifax, trigs, wk, nt);
for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n ; i++ )
    {
        c[i*incn+j*incm] /= n ;
    }
}

printf( "\n\t< Forward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tSolution\n\n" );
printf( "\t Real Part                Imaginary Part\n" );
for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t creal(c[%3d]) = %8.3g          cimag(c[%3d]) = %8.3g\n",
                i*incn+j*incm, creal(c[i*incn+j*incm]),
                i*incn+j*incm, cimag(c[i*incn+j*incm]) );
    }
}

isw = -1;
ierr = ASL_hfcmbf(n, m, c, incn, incm, isw, ifax, trigs, wk, nt);

printf( "\n\t< Backward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\n\tSolution\n\n" );
printf( "\t Real Part                Imaginary Part\n" );
for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t creal(c[%3d]) = %8.3g          cimag(c[%3d]) = %8.3g\n",
                i*incn+j*incm, creal(c[i*incn+j*incm]),
                i*incn+j*incm, cimag(c[i*incn+j*incm]) );
    }
}

free( c );
free( trigs );
free( wk );

return 0;
}

```

(d) Output results

```

*** ASL_hfcmbf , ASL_hfcmbf ***

** Input **

n      = 8
m      = 4
incn   = 1
incm   = 9
nt     = 2

Real Part                Imaginary Part
creal(c[ 0]) =          1          cimag(c[ 0]) =          4
creal(c[ 1]) =          2          cimag(c[ 1]) =          3
creal(c[ 2]) =          3          cimag(c[ 2]) =          2
creal(c[ 3]) =          4          cimag(c[ 3]) =          1
creal(c[ 4]) =          4          cimag(c[ 4]) =          1
creal(c[ 5]) =          3          cimag(c[ 5]) =          2
creal(c[ 6]) =          2          cimag(c[ 6]) =          3
creal(c[ 7]) =          1          cimag(c[ 7]) =          4
creal(c[ 9]) =          1          cimag(c[ 9]) =          2
creal(c[10]) =          1          cimag(c[10]) =          2
creal(c[11]) =          2          cimag(c[11]) =          1
creal(c[12]) =          2          cimag(c[12]) =          1
creal(c[13]) =          2          cimag(c[13]) =          1
creal(c[14]) =          2          cimag(c[14]) =          1
creal(c[15]) =          1          cimag(c[15]) =          2
creal(c[16]) =          1          cimag(c[16]) =          2
creal(c[18]) =          1          cimag(c[18]) =          2
creal(c[19]) =          1          cimag(c[19]) =          2
creal(c[20]) =          1          cimag(c[20]) =          2
creal(c[21]) =          1          cimag(c[21]) =          2

```

```

creal(c[ 22]) =      2      cimag(c[ 22]) =      1
creal(c[ 23]) =      2      cimag(c[ 23]) =      1
creal(c[ 24]) =      2      cimag(c[ 24]) =      1
creal(c[ 25]) =      2      cimag(c[ 25]) =      1
creal(c[ 27]) =      1      cimag(c[ 27]) =      1
creal(c[ 28]) =      1      cimag(c[ 28]) =      1
creal(c[ 29]) =      1      cimag(c[ 29]) =      1
creal(c[ 30]) =      1      cimag(c[ 30]) =      1
creal(c[ 31]) =      1      cimag(c[ 31]) =      1
creal(c[ 32]) =      1      cimag(c[ 32]) =      1
creal(c[ 33]) =      1      cimag(c[ 33]) =      1
creal(c[ 34]) =      1      cimag(c[ 34]) =      1

```

\*\* Output \*\*

< Forward Transform >  
 ierr = 0

Solution

Real Part		Imaginary Part	
creal(c[ 0]) =	2.5	cimag(c[ 0]) =	2.5
creal(c[ 1]) =	-1.03	cimag(c[ 1]) =	0.427
creal(c[ 2]) =	0	cimag(c[ 2]) =	0
creal(c[ 3]) =	-0.0732	cimag(c[ 3]) =	-0.0303
creal(c[ 4]) =	0	cimag(c[ 4]) =	0
creal(c[ 5]) =	0.0303	cimag(c[ 5]) =	0.0732
creal(c[ 6]) =	0	cimag(c[ 6]) =	0
creal(c[ 7]) =	-0.427	cimag(c[ 7]) =	1.03
creal(c[ 9]) =	1.5	cimag(c[ 9]) =	1.5
creal(c[ 10]) =	-0.427	cimag(c[ 10]) =	0.177
creal(c[ 11]) =	0	cimag(c[ 11]) =	0
creal(c[ 12]) =	0.177	cimag(c[ 12]) =	0.0732
creal(c[ 13]) =	0	cimag(c[ 13]) =	0
creal(c[ 14]) =	-0.0732	cimag(c[ 14]) =	-0.177
creal(c[ 15]) =	0	cimag(c[ 15]) =	0
creal(c[ 16]) =	-0.177	cimag(c[ 16]) =	0.427
creal(c[ 18]) =	1.5	cimag(c[ 18]) =	1.5
creal(c[ 19]) =	0.177	cimag(c[ 19]) =	0.427
creal(c[ 20]) =	0	cimag(c[ 20]) =	0
creal(c[ 21]) =	-0.0732	cimag(c[ 21]) =	0.177
creal(c[ 22]) =	0	cimag(c[ 22]) =	0
creal(c[ 23]) =	-0.177	cimag(c[ 23]) =	0.0732
creal(c[ 24]) =	0	cimag(c[ 24]) =	0
creal(c[ 25]) =	-0.427	cimag(c[ 25]) =	-0.177
creal(c[ 27]) =	1	cimag(c[ 27]) =	1
creal(c[ 28]) =	0	cimag(c[ 28]) =	0
creal(c[ 29]) =	0	cimag(c[ 29]) =	0
creal(c[ 30]) =	0	cimag(c[ 30]) =	0
creal(c[ 31]) =	0	cimag(c[ 31]) =	0
creal(c[ 32]) =	0	cimag(c[ 32]) =	0
creal(c[ 33]) =	0	cimag(c[ 33]) =	0
creal(c[ 34]) =	0	cimag(c[ 34]) =	0

< Backward Transform >  
 ierr = 0

Solution

Real Part		Imaginary Part	
creal(c[ 0]) =	1	cimag(c[ 0]) =	4
creal(c[ 1]) =	2	cimag(c[ 1]) =	3
creal(c[ 2]) =	3	cimag(c[ 2]) =	2
creal(c[ 3]) =	4	cimag(c[ 3]) =	1
creal(c[ 4]) =	4	cimag(c[ 4]) =	1
creal(c[ 5]) =	3	cimag(c[ 5]) =	2
creal(c[ 6]) =	2	cimag(c[ 6]) =	3
creal(c[ 7]) =	1	cimag(c[ 7]) =	4
creal(c[ 9]) =	1	cimag(c[ 9]) =	2
creal(c[ 10]) =	1	cimag(c[ 10]) =	2
creal(c[ 11]) =	2	cimag(c[ 11]) =	1
creal(c[ 12]) =	2	cimag(c[ 12]) =	1
creal(c[ 13]) =	2	cimag(c[ 13]) =	1
creal(c[ 14]) =	2	cimag(c[ 14]) =	1
creal(c[ 15]) =	1	cimag(c[ 15]) =	2
creal(c[ 16]) =	1	cimag(c[ 16]) =	2
creal(c[ 18]) =	1	cimag(c[ 18]) =	2
creal(c[ 19]) =	1	cimag(c[ 19]) =	2
creal(c[ 20]) =	1	cimag(c[ 20]) =	2
creal(c[ 21]) =	1	cimag(c[ 21]) =	2
creal(c[ 22]) =	2	cimag(c[ 22]) =	1

<code>creal(c[ 23]) =</code>	<code>2</code>	<code>cimag(c[ 23]) =</code>	<code>1</code>
<code>creal(c[ 24]) =</code>	<code>2</code>	<code>cimag(c[ 24]) =</code>	<code>1</code>
<code>creal(c[ 25]) =</code>	<code>2</code>	<code>cimag(c[ 25]) =</code>	<code>1</code>
<code>creal(c[ 27]) =</code>	<code>1</code>	<code>cimag(c[ 27]) =</code>	<code>1</code>
<code>creal(c[ 28]) =</code>	<code>1</code>	<code>cimag(c[ 28]) =</code>	<code>1</code>
<code>creal(c[ 29]) =</code>	<code>1</code>	<code>cimag(c[ 29]) =</code>	<code>1</code>
<code>creal(c[ 30]) =</code>	<code>1</code>	<code>cimag(c[ 30]) =</code>	<code>1</code>
<code>creal(c[ 31]) =</code>	<code>1</code>	<code>cimag(c[ 31]) =</code>	<code>1</code>
<code>creal(c[ 32]) =</code>	<code>1</code>	<code>cimag(c[ 32]) =</code>	<code>1</code>
<code>creal(c[ 33]) =</code>	<code>1</code>	<code>cimag(c[ 33]) =</code>	<code>1</code>
<code>creal(c[ 34]) =</code>	<code>1</code>	<code>cimag(c[ 34]) =</code>	<code>1</code>

---

## 6.4 MULTIPLE ONE-DIMENSIONAL REAL FOURIER TRANSFORM

### 6.4.1 [DEPRECATED]ASL\_qfrmfb, ASL\_pfrmfb

#### Multiple One-Dimensional Real Fourier Transforms (Including Initialization)

##### (1) Function

###### Forward transform

ASL\_qfrmfb or ASL\_pfrmfb obtains a half period of the  $m$ -fold one-dimensional Fourier forward transform (arbitrary radix) for the real data  $r_{k,l}$  ( $k = 0, \dots, n-1$ ;  $l = 1, \dots, m$ ).

$$c_{j,l} = \sum_{k=0}^{n-1} r_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, \lfloor \frac{n}{2} \rfloor; l = 1, \dots, m)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationship.

$$c_{n-j,l}^* = c_{j,l}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

###### Backward transform

Given the half period  $c_{j,l}$  ( $j = 0, \dots, \lfloor \frac{n}{2} \rfloor$ ;  $l = 1, \dots, m$ ) for  $n$  complex data groups  $c_{j,l}$  ( $j = 0, \dots, n-1$ ;  $l = 1, \dots, m$ ) satisfying  $c_{n-j,l}^* = c_{j,l}$ , ASL\_qfrmfb or ASL\_pfrmfb obtains the  $m$ -fold one-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$\begin{aligned} r_{k,l} &= \sum_{j=0}^{n-1} c_{j,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \\ &= c_{0,l} + (-1)^k \hat{c}_{\frac{n}{2},l} + 2 \sum_{j=1}^{\lfloor \frac{n}{2} \rfloor - 1} \Re\{c_{j,l} e^{2\pi\sqrt{-1}\frac{jk}{n}}\} \\ &= c_{0,l} + (-1)^k \hat{c}_{\frac{n}{2},l} + 2 \sum_{j=1}^{\lfloor \frac{n}{2} \rfloor - 1} \left[ \Re\{c_{j,l}\} \cos(2\pi\frac{jk}{n}) - \Im\{c_{j,l}\} \sin(2\pi\frac{jk}{n}) \right] \\ &\quad (k = 0, \dots, n-1; l = 1, \dots, m) \end{aligned}$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  and  $\Im\{z\}$  represent the real and imaginary parts of the complex number  $z$ , respectively. Also, when  $n$  is odd,  $\hat{c}_{\frac{n}{2},l} = 0$ , and when  $n$  is even,  $\hat{c}_{\frac{n}{2},l} = c_{\frac{n}{2},l}$ .

##### (2) Usage

Double precision:

ierr = ASL\_qfrmfb (n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfrmfb (n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	r	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Input	Input data $r_{k,l}$ (Forward transform) or $c_{j,l}$ (Backward transform) (See Note (b)). <b>Size:</b> $\text{incn} \times (n) + \text{incm} \times (m - 1) + 1$ , where $n$ is an odd, or $\text{incn} \times (n + 1) + \text{incm} \times (m - 1) + 1$ , where $n$ is an even.
				Output	Output results $c_{j,l}$ (Forward transform), or $r_{k,l}$ (Backward transform) (See Notes (b) and (c))
4	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
5	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
6	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	20	Output	Factorization results and number of factors (See Note (d))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	n	Output	Trigonometric function table (See Note (d))
9	wk	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Work	Work area <b>Size:</b> $(n+1) \times m$ , where $n$ is an odd, or $(n+2) \times m$ , where $n$ is an even.
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)



(4) **Restrictions**

- (a)  $n \geq 1, m \geq 1$
- (b)  $\text{incn} \geq 1, \text{incm} \geq 1$
- (c)  $\text{incn} \geq m \times \text{gcm}(\text{incn}, \text{incm})$  or:  
 In the case where  $n$  is an odd:  
 $\text{incm} \geq (n+1) \times \text{gcm}(\text{incn}, \text{incm})$   
 or if  $n$  is an even:  
 $\text{incm} \geq (n+2) \times \text{gcm}(\text{incn}, \text{incm})$   
 where,  $\text{gcm}(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $\text{isw} \in \{0, 1, -1\}$
- (e)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	Input-time contents are output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) When the number of data  $n$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n = 289(=17^2)$ , it is usually more efficient to set  $n = 300(=2^2 \times 3 \times 5^2)$ ,  $n = 320(=2^6 \times 5)$ ,  $n = 384(=2^7 \times 3)$  or the like.
- (b) The real data  $r_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$  and elements of array  $r$  are associated as follows.

$$r_{k,l} \leftrightarrow r[\text{incn} * k + \text{incm} * (l - 1)]$$

For example, if we let  $\text{incn}=1$  and  $\text{incm}=n$ , then the associations are as follows:

$$r_{k,l} \leftrightarrow r[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $\text{incn}=m$  and  $\text{incm}=1$ , then the associations are as follows:

$$r_{k,l} \leftrightarrow r[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . When computing the backward transform, if  $n(=n)$  is odd, then  $r[\text{incn} * n + \text{incm} * (l - 1)] = 0$ , and when  $n$  is even, then  $r[\text{incn} * n + \text{incm} * (l - 1)] = r[\text{incn} * (n + 1) + \text{incm} * (l - 1)] = 0$ . If we let the real and imaginary parts of the complex data  $c_{j,l}(j = 0, \dots, \lfloor \frac{n}{2} \rfloor; l = 1, \dots, m)$  be  $\Re\{c_{j,l}\}$  and  $\Im\{c_{j,l}\}$ , respectively, the  $c_{j,l}$  and

elements of array  $r$  are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned}\Re\{c_{j,l}\} &\leftrightarrow r[\text{incn} * (2j) + \text{incm} * (1 - 1)] \\ \Im\{c_{j,l}\} &\leftrightarrow r[\text{incn} * (2j + 1) + \text{incm} * (1 - 1)]\end{aligned}$$

From the properties of a real Fourier transform, when  $n$  is odd,  $\Im\{c_{0,l}\} = 0$ , and when  $n$  is even,  $\Im\{c_{0,l}\} = \Im\{c_{\frac{n}{2},l}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array  $r$ , they are considered to be zero when processing is performed. Since the elements  $c_{j,l}$  ( $j = \lfloor \frac{n}{2} \rfloor + 1, \dots, n - 1$ ;  $l = 1, \dots, m$ ) can be obtained according to the following relationship from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$c_{n-j,l} = c_{j,l}^*$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ) be represented by  $\hat{r}_{k,l}$  ( $k = 0, \dots, n - 1$ ;  $l = 1, \dots, m$ ), then the following relationship holds.

$$\hat{r}_{k,l} = nr_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data  $n$ , you should call this function once, and then use the after-initialization transform 6.4.2  $\left\{ \begin{array}{l} \text{ASL\_qfrmbf} \\ \text{ASL\_pfrmbf} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays `ifax` and `trigs` so they can be used as input to the function 6.4.2  $\left\{ \begin{array}{l} \text{ASL\_qfrmbf} \\ \text{ASL\_pfrmbf} \end{array} \right\}$ .  
 To perform initialization only by setting `isw=0`, you need not set input data for array  $r$ .

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.4.2 (7).

### 6.4.2 [DEPRECATED]ASL\_qfrmbf, ASL\_pfrmbf Multiple One-Dimensional Real Fourier Transforms (After Initialization)

(1) **Function**

**Forward transform**

ASL\_qfrmbf or ASL\_pfrmbf obtains a half period of the  $m$ -fold one-dimensional Fourier forward transform (arbitrary radix) for the real data  $r_{k,l}(k = 0, \dots, n - 1; l = 1, \dots, m)$ .

$$c_{j,l} = \sum_{k=0}^{n-1} r_{k,l} e^{-2\pi\sqrt{-1}\frac{jk}{n}} \quad (j = 0, \dots, \lfloor \frac{n}{2} \rfloor; l = 1, \dots, m)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationship.

$$c_{n-j,l}^* = c_{j,l}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

**Backward transform**

Given the half period  $c_{j,l}(j = 0, \dots, \lfloor \frac{n}{2} \rfloor; l = 1, \dots, m)$  for  $n$  complex data groups  $c_{j,l}(j = 0, \dots, n - 1; l = 1, \dots, m)$  satisfying  $c_{n-j,l}^* = c_{j,l}$ , ASL\_qfrmbf or ASL\_pfrmbf obtains the  $m$ -fold one-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$\begin{aligned} r_{k,l} &= \sum_{j=0}^{n-1} c_{j,l} e^{2\pi\sqrt{-1}\frac{jk}{n}} \\ &= c_{0,l} + (-1)^k \hat{c}_{\frac{n}{2},l} + 2 \sum_{j=1}^{\lfloor \frac{n}{2} \rfloor - 1} \Re\{c_{j,l} e^{2\pi\sqrt{-1}\frac{jk}{n}}\} \\ &= c_{0,l} + (-1)^k \hat{c}_{\frac{n}{2},l} + 2 \sum_{j=1}^{\lfloor \frac{n}{2} \rfloor - 1} \left[ \Re\{c_{j,l}\} \cos(2\pi\frac{jk}{n}) - \Im\{c_{j,l}\} \sin(2\pi\frac{jk}{n}) \right] \\ &\quad (k = 0, \dots, n - 1; l = 1, \dots, m) \end{aligned}$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  and  $\Im\{z\}$  represent the real and imaginary parts of the complex number  $z$ , respectively. Also, when  $n$  is odd,  $\hat{c}_{\frac{n}{2},l} = 0$ , and when  $n$  is even,  $\hat{c}_{\frac{n}{2},l} = c_{\frac{n}{2},l}$ .

(2) **Usage**

Double precision:

ierr = ASL\_qfrmbf (n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfrmbf (n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	n	I	1	Input	Number of transformed data values $n$ (See Note (a))
2	m	I	1	Input	Multiplicity $m$
3	r	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Input	Input data $r_{k,l}$ (Forward transform) or $c_{j,l}$ (Backward transform) (See Note (b)). <b>Size:</b> $\text{incn} \times (n) + \text{incm} \times (m - 1) + 1$ , where $n$ is an odd, or $\text{incn} \times (n + 1) + \text{incm} \times (m - 1) + 1$ , where $n$ is an even.
				Output	Output results $c_{j,l}$ (Forward transform), or $r_{k,l}$ (Backward transform) (See Notes (b) and (c))
4	incn	I	1	Input	The stride between each transformed datum in storage (See Note (b))
5	incm	I	1	Input	The stride between the first elements of each transformed data in storage (See Note (b))
6	isw	I	1	Input	Processing switch isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	20	Input	Factorization results and number of factors (See Note (a))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	n	Input	Trigonometric function table (See Note (a))
9	wk	$\begin{cases} D^* \\ R^* \end{cases}$	See Contents	Work	Work area <b>Size:</b> $(n+1) \times m$ , where $n$ is an odd, or $(n+2) \times m$ , where $n$ is an even.
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $n \geq 1, m \geq 1$
- (b)  $\text{incn} \geq 1, \text{incm} \geq 1$
- (c)  $\text{incn} \geq m \times \text{gcm}(\text{incn}, \text{incm})$  or  
 In the case where  $n$  is an odd:  
 $\text{incm} \geq (n+1) \times \text{gcm}(\text{incn}, \text{incm})$   
 or if  $n$  is an even:  
 $\text{incm} \geq (n+2) \times \text{gcm}(\text{incn}, \text{incm})$   
 where,  $\text{gcm}(i, j)$  is the greatest common measure between  $i$  and  $j$ .
- (d)  $\text{isw} \in \{1, -1\}$
- (e)  $\text{nt} \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$n$ was equal to 1.	Input-time contents are output unchanged.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) This function can be used to repeatedly compute the transform for the same number of transformed data  $n$  after the including-initialization function 6.4.1  $\left\{ \begin{matrix} \text{ASL\_qfrmbf} \\ \text{ASL\_pfrmbf} \end{matrix} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.
- (b) The real data  $r_{k,l} (k = 0, \dots, n - 1; l = 1, \dots, m)$  and elements of array  $r$  are associated as follows.

$$r_{k,l} \leftrightarrow r[\text{incn} * k + \text{incm} * (l - 1)]$$

For example, if we let  $\text{incn}=1$  and  $\text{incm}=n$ , then the associations are as follows:

$$r_{k,l} \leftrightarrow r[k + n * (l - 1)]$$

and the data is stored so that it is packed consecutively for subscript  $k$ . If we let  $\text{incn}=m$  and  $\text{incm}=1$ , then the associations are as follows:

$$r_{k,l} \leftrightarrow r[(l - 1) + m * k]$$

and the data is stored so that it is packed consecutively for subscript  $l$ . When computing the backward transform, if  $n(=n)$  is odd, then  $r[\text{incn} * n + \text{incm} * (l - 1)] = 0$ , and when  $n$  is even, then  $r[\text{incn} * n + \text{incm} * (l - 1)] = r[\text{incn} * (n + 1) + \text{incm} * (l - 1)] = 0$ . If we let the real and imaginary parts of the complex data  $c_{j,l} (j = 0, \dots, \lfloor \frac{n}{2} \rfloor; l = 1, \dots, m)$  be  $\Re\{c_{j,l}\}$  and  $\Im\{c_{j,l}\}$ , respectively, the  $c_{j,l}$  and

elements of array  $r$  are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned}\Re\{c_{j,l}\} &\leftrightarrow r[\text{incn} * (2j) + \text{incm} * (1 - 1)] \\ \Im\{c_{j,l}\} &\leftrightarrow r[\text{incn} * (2j + 1) + \text{incm} * (1 - 1)]\end{aligned}$$

From the properties of a real Fourier transform, when  $n$  is odd,  $\Im\{c_{0,l}\} = 0$ , and when  $n$  is even,  $\Im\{c_{0,l}\} = \Im\{c_{\frac{n}{2},l}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array  $r$ , they are considered to be zero when processing is performed. Since the elements  $c_{j,l}$  ( $j = \lfloor \frac{n}{2} \rfloor + 1, \dots, n - 1; l = 1, \dots, m$ ) can be obtained according to the following relationship from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$c_{n-j,l} = c_{j,l}^*$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of transformed data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ) be represented by  $\hat{r}_{k,l}$  ( $k = 0, \dots, n - 1; l = 1, \dots, m$ ), then the following relationship holds.

$$\hat{r}_{k,l} = nr_{k,l} \quad (k = 0, \dots, n - 1; l = 1, \dots, m)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

(a) Problem

Compute the multiple one-dimensional real Fourier forward and backward transforms using the following sequence of numbers as input data.

```
r[ 0]= 1.000
r[ 1]= 2.000
r[ 2]= 3.000
r[ 3]= 4.000
r[ 4]= 5.000
r[ 5]= 6.000
r[ 6]= 7.000
r[ 7]= 8.000
r[12]= 1.000
r[13]= 1.000
r[14]= 2.000
r[15]= 2.000
r[16]= 3.000
r[17]= 3.000
r[18]= 4.000
r[19]= 4.000
r[24]= 1.000
r[25]= 1.000
r[26]= 1.000
r[27]= 1.000
r[28]= 2.000
r[29]= 2.000
r[30]= 2.000
r[31]= 2.000
r[36]= 1.000
r[37]= 1.000
r[38]= 1.000
r[39]= 1.000
r[40]= 1.000
r[41]= 1.000
r[42]= 1.000
r[43]= 1.000
```

(b) Input data

Array r, n=8, m=4, incn=1, incm=12, isw=1(Forward transform), isw=-1 (Backward transform) and nt=2.

(c) Main program

```
/*      C interface example for ASL_qfrmbf , ASL_pfrmbf */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int ld=46;
```

```

int n;      int m;
double *r;
int incn;  int incm;
int isw;
int ifax[20];  double *trigs;
double *wk;
int nt;
int ierr;
int i,j;
FILE *fp;

fp = fopen( "qfrmbf.dat", "r" );
if( fp == NULL )
{
    printf( "file open error\n" );
    return -1;
}

printf( "    *** ASL_qfrmbf , ASL_qfrmbf ***\n" );
printf( "\n    ** Input **\n\n" );

r = ( double * )malloc((size_t)( sizeof(double) * ld ));
if( r == NULL )
{
    printf( "no enough memory for array r\n" );
    return -1;
}

trigs = ( double * )malloc((size_t)( sizeof(double) * ld ));
if( trigs == NULL )
{
    printf( "no enough memory for array trigs\n" );
    return -1;
}

wk = ( double * )malloc((size_t)( sizeof(double) * ld ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

fscanf( fp, "%d,%d,%d,%d,%d", &n, &m, &incn, &incm, &nt );
for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n ; i++ )
    {
        fscanf( fp, "%lf", &r[i*incn+j*incm] );
    }
}

printf("\t  n    = %d \n", n );
printf("\t  m    = %d \n", m );
printf("\t  incn = %d \n", incn );
printf("\t  incm = %d \n", incm );
printf("\t  nt   = %d \n\n", nt );

printf( "\t Real Part\n" );
for( j=0 ; j<m ; j++ )
{
    printf( "\t" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "    r[%3d] =%4.1f",
            i*incn+j*incm, r[i*incn+j*incm] );
        if((i+1)%4==0) printf( "\n\t" );
    }
    printf( "\n" );
}

fclose( fp );

printf( "\n    ** Output **\n" );

isw = 1;
ierr = ASL_qfrmbf(n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n+2 ; i++ )
    {
        r[i*incn+j*incm] /= n ;
    }
}

printf( "\n\t< Forward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

```



```

printf( "\n\tSolution\n\n" );
printf( "\t Real Part                Imaginary Part\n" );
for( j=0 ; j<m ; j++ )
{
    for( i=0 ; i<n+2 ; i=i+2 )
    {
        printf( "\t r[%3d] = %8.3g      r[%3d] = %8.3g\n",
                i*incn+j*incm, r[i*incn+j*incm],
                (i+1)*incn+j*incm, r[(i+1)*incn+j*incm] );
    }
    printf( "\n" );
}

isw = -1;
ierr = ASL_qfrmbf(n, m, r, incn, incm, isw, ifax, trigs, wk, nt);

printf( "\n\t< Backward Transform >\n" );
printf( "\t ierr = %6d\n", ierr );

printf( "\n\tSolution\n\n" );
printf( "\t Real Part\n" );
for( j=0 ; j<m ; j++ )
{
    printf( "\t" );
    for( i=0 ; i<n+2 ; i++ )
    {
        printf( "      r[%3d] =%4.1f",
                i*incn+j*incm, r[i*incn+j*incm] );
        if((i+1)%4==0) printf( "\n\t" );
    }
    printf( "\n" );
}

free( r );
free( trigs );
free( wk );

return 0;
}

```

(d) Output results

```

*** ASL_qfrmbf , ASL_qfrmbf ***

** Input **

n      = 8
m      = 4
incn   = 1
incm   = 12
nt     = 2

Real Part
r[  0] = 1.0   r[  1] = 2.0   r[  2] = 3.0   r[  3] = 4.0
r[  4] = 5.0   r[  5] = 6.0   r[  6] = 7.0   r[  7] = 8.0

r[ 12] = 1.0   r[ 13] = 1.0   r[ 14] = 2.0   r[ 15] = 2.0
r[ 16] = 3.0   r[ 17] = 3.0   r[ 18] = 4.0   r[ 19] = 4.0

r[ 24] = 1.0   r[ 25] = 1.0   r[ 26] = 1.0   r[ 27] = 1.0
r[ 28] = 2.0   r[ 29] = 2.0   r[ 30] = 2.0   r[ 31] = 2.0

r[ 36] = 1.0   r[ 37] = 1.0   r[ 38] = 1.0   r[ 39] = 1.0
r[ 40] = 1.0   r[ 41] = 1.0   r[ 42] = 1.0   r[ 43] = 1.0

** Output **

< Forward Transform >
ierr =      0

Solution

Real Part                Imaginary Part
r[  0] =      4.5        r[  1] =      0
r[  2] =     -0.5        r[  3] =     1.21
r[  4] =     -0.5        r[  5] =      0.5
r[  6] =     -0.5        r[  7] =     0.207
r[  8] =     -0.5        r[  9] =      0

r[ 12] =      2.5        r[ 13] =      0
r[ 14] =     -0.25       r[ 15] =     0.604
r[ 16] =     -0.25       r[ 17] =     0.25
r[ 18] =     -0.25       r[ 19] =     0.104
r[ 20] =      0          r[ 21] =      0

r[ 24] =      1.5        r[ 25] =      0
r[ 26] =     -0.125      r[ 27] =     0.302

```

```
r[ 28] =      0          r[ 29] =      0
r[ 30] =   -0.125       r[ 31] =   0.0518
r[ 32] =      0          r[ 33] =      0

r[ 36] =      1          r[ 37] =      0
r[ 38] =      0          r[ 39] =      0
r[ 40] =      0          r[ 41] =      0
r[ 42] =      0          r[ 43] =      0
r[ 44] =      0          r[ 45] =      0
```

```
< Backward Transform >
ierr =      0
```

Solution

```
Real Part
r[  0] = 1.0    r[  1] = 2.0    r[  2] = 3.0    r[  3] = 4.0
r[  4] = 5.0    r[  5] = 6.0    r[  6] = 7.0    r[  7] = 8.0
r[  8] = 0.0    r[  9] = 0.0
r[ 12] = 1.0    r[ 13] = 1.0    r[ 14] = 2.0    r[ 15] = 2.0
r[ 16] = 3.0    r[ 17] = 3.0    r[ 18] = 4.0    r[ 19] = 4.0
r[ 20] = 0.0    r[ 21] = 0.0
r[ 24] = 1.0    r[ 25] = 1.0    r[ 26] = 1.0    r[ 27] = 1.0
r[ 28] = 2.0    r[ 29] = 2.0    r[ 30] = 2.0    r[ 31] = 2.0
r[ 32] = 0.0    r[ 33] = 0.0
r[ 36] = 1.0    r[ 37] = 1.0    r[ 38] = 1.0    r[ 39] = 1.0
r[ 40] = 1.0    r[ 41] = 1.0    r[ 42] = 1.0    r[ 43] = 1.0
r[ 44] = 0.0    r[ 45] = 0.0
```

---

## 6.5 TWO-DIMENSIONAL COMPLEX FOURIER TRANSFORM (REAL ARGUMENT TYPE)

### 6.5.1 [DEPRECATED]ASL\_qfc2fb, ASL\_pfc2fb

#### Two-Dimensional Complex Fourier Transform (Including Initialization)

##### (1) Function

###### Forward transform

ASL\_qfc2fb or ASL\_pfc2fb computes the two-dimensional complex Fourier forward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

###### Backward transform

ASL\_qfc2fb or ASL\_pfc2fb computes the two-dimensional complex Fourier backward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

##### (2) Usage

Double precision:

ierr = ASL\_qfc2fb (nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfc2fb (nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	cr	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lx×ly	Input	Real part of input data $c_{k_x, k_y}$ (See Note (b))
				Output	Real part of output data $d_{j_x, j_y}$ (See Notes (b) and (c))
4	ci	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lx×ly	Input	Imaginary part of input data $c_{k_x, k_y}$ (See Note (b))
				Output	Imaginary part of output results $d_{j_x, j_y}$ (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array cr and ci (See Note (b))
6	ly	I	1	Input	Second dimension of array cr and ci (See Note (b))
7	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
8	ifax	I*	40	Output	Factorization results and number of factors (See Note (d))
9	trigs	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times (n_x + n_y)$	Output	Trigonometric function table (See Note (d))
10	wk	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times l_x \times l_y$	Work	Work area
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y$
- (c)  $isw \in \{0, 1, -1\}$
- (d)  $nt \geq 1$

(5) Error indicator (Return Value)

ier value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) Notes

- (a) When the number of data  $n_x$  or  $n_y$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n_x = 289 (=17^2)$ , it is usually more efficient to set  $n_x = 300 (=2^2 \times 3 \times 5^2)$ ,  $n_x = 320 (=2^6 \times 5)$ ,  $n_x = 384 (=2^7 \times 3)$  or the like.
- (b) If we let the real and imaginary parts of the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) be  $\Re\{c_{k_x, k_y}\}$  and  $\Im\{c_{k_x, k_y}\}$ , respectively, the  $c_{k_x, k_y}$  and elements of arrays  $cr$  and  $ci$  are associated as follows.

$$\begin{aligned} \Re\{c_{k_x, k_y}\} &\leftrightarrow cr[k_x + lx * k_y] \\ \Im\{c_{k_x, k_y}\} &\leftrightarrow ci[k_x + lx * k_y] \end{aligned}$$

Similarly, for the complex data  $d_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ).

**The adjustable dimensions  $lx$  and  $ly$  of arrays  $cr$  and  $ci$  should be set to odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is even,  $lx = n_x + 1$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{c}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y} = n_x n_y c_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data ( $n_x, n_y$ ), you should call this function once, and then use the after-initialization transform 6.5.2  $\left\{ \begin{matrix} ASL\_qfc2bf \\ ASL\_pfc2bf \end{matrix} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays  $ifax$  and  $trigs$  so they can be used as input to the function 6.5.2  $\left\{ \begin{matrix} ASL\_qfc2bf \\ ASL\_pfc2bf \end{matrix} \right\}$ .  
 To perform initialization only by setting  $isw=0$ , you need not set input data for arrays  $cr$  and  $ci$ .
- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period,

the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

(f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.5.2 (7).

## 6.5.2 [DEPRECATED]ASL\_qfc2bf, ASL\_pfc2bf Two-Dimensional Complex Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_qfc2bf or ASL\_pfc2bf computes the two-dimensional complex Fourier forward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

#### Backward transform

ASL\_qfc2bf or ASL\_pfc2bf computes the two-dimensional complex Fourier backward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

### (2) Usage

Double precision:

ierr = ASL\_qfc2bf (nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfc2bf (nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);

### (3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	cr	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Input	Real part of input data $c_{k_x, k_y}$ (See Note (b))
				Output	Real part of output data $d_{j_x, j_y}$ (See Notes (b) and (c))
4	ci	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Input	Imaginary part of input data $c_{k_x, k_y}$
				Output	Imaginary part of output results $d_{j_x, j_y}$ (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array cr and ci (See Note (b))
6	ly	I	1	Input	Second dimension of array cr and ci (See Note (b))
7	isw	I	1	Input	Processing switch isw= 1:Forward transform isw=-1:Backward transform
8	ifax	I*	40	Input	Factorization results and number of factors (See Note (a))
9	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y)$	Input	Trigonometric function table (See Note (a))
10	wk	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times l_x \times l_y$	Work	Work area
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)

### (4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y$
- (c)  $isw \in \{1, -1\}$
- (d)  $nt \geq 1$



(5) **Error indicator (Return Value)**

ier value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

(a) This function can be used to repeatedly compute the transform for the same number of data ( $n_x, n_y$ ) after the including-initialization function 6.5.1  $\left\{ \begin{array}{l} \text{ASL\_qfc2fb} \\ \text{ASL\_pfc2fb} \end{array} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.

(b) If we let the real and imaginary parts of the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ) be  $\Re\{c_{k_x, k_y}\}$  and  $\Im\{c_{k_x, k_y}\}$ , respectively, the  $c_{k_x, k_y}$  and elements of arrays cr and ci are associated as follows.

$$\begin{aligned} \Re\{c_{k_x, k_y}\} &\leftrightarrow \text{cr}[k_x + l_x * k_y] \\ \Im\{c_{k_x, k_y}\} &\leftrightarrow \text{ci}[k_x + l_x * k_y] \end{aligned}$$

Similarly, for the complex data  $d_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1$ ).

**The adjustable dimensions  $l_x$  and  $l_y$  of arrays cr and ci should be set to odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is even,  $l_x = n_x + 1$  is set.**

(c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{c}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y} = n_x n_y c_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

(d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c (t - iT)}{\pi(t - iT)}$$

(e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) Example

(a) Problem

Compute the two-dimensional complex Fourier forward and backward transforms using

$$c_{k_x, k_y} = (k_x + 1) + (k_y + 1) + \sqrt{-1} \frac{(k_x + 1)(k_y + 1)}{n_x n_y}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

as input data.

(b) Input data

Array cr and ci, nx=5, ny=4, lx=5, ly=5, isw=1 (forward transform), isw=-1 (backward transform) and nt=2.

(c) Main program

```

/*      C Interface example for ASL_qfc2fb , ASL_qfc2bf */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    int nx = 5; int ny = 4;
    int lx = 5; int ly = 5;
    double *cr; double *ci;
    int isw;
    int ifax[40];
    double *trigs;
    double *wk;
    int nt = 2;
    int ierr;
    int i, j;

    printf( "      *** ASL_qfc2fb , ASL_qfc2bf ***\n" );
    printf( "\n      ** Input **\n" );

    cr = ( double * )malloc((size_t)( sizeof(double) * (lx*ly) ));
    if( cr == NULL )
    {
        printf( "no enough memory for array cr\n" );
        return -1;
    }

    ci = ( double * )malloc((size_t)( sizeof(double) * (lx*ly) ));
    if( ci == NULL )
    {
        printf( "no enough memory for array ci\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*(nx+ny)) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double * )malloc((size_t)( sizeof(double) * (2*lx*ly) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    printf( "\tnx = %6d\n", nx );
    printf( "\tny = %6d\n", ny );
    printf( "\tnt = %6d\n", nt );

    for( j=1 ; j<=ny ; j++ )
    {
        for( i=1 ; i<=nx ; i++ )
        {
            cr[(i-1)+lx*(j-1)]=(double)(i+j) ;
            ci[(i-1)+lx*(j-1)]=(double)(i*j)/(double)(nx*ny) ;
        }
    }

    printf( "\tcr[ix][iy]\n" );
    for( i=0 ; i<nx ; i++ )
    {

```

```

        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j], ci[i+lx*j] );
        }
        printf( "\n" );
    }

    isw = 1;
    ierr = ASL_qfc2fb(nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);
    for( i=0 ; i<lx*ly ; i++ )
    {
        cr[i] /= (double)(nx*ny);
        ci[i] /= (double)(nx*ny);
    }

    printf( "\n    ** Output **\n" );
    printf( "\t< Forward Transform >\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\tcr[ix][iy]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j], ci[i+lx*j] );
        }
        printf( "\n" );
    }

    isw = -1;
    ierr = ASL_qfc2bf(nx, ny, cr, ci, lx, ly, isw, ifax, trigs, wk, nt);

    printf( "\t< Backward Transform >\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\tcr[ix][iy]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j], ci[i+lx*j] );
        }
        printf( "\n" );
    }

    free( cr );
    free( ci );
    free( trigs );
    free( wk );

    return 0;
}

```

(d) Output results

```

*** ASL_qfc2fb , ASL_qfc2bf ***

** Input **
nx =      5
ny =      4
nt =      2
cr[ix][iy]
(      2,      0.05) (      3,      0.1) (      4,      0.15) (      5,      0.2)
(      3,      0.1) (      4,      0.2) (      5,      0.3) (      6,      0.4)
(      4,      0.15) (      5,      0.3) (      6,      0.45) (      7,      0.6)
(      5,      0.2) (      6,      0.4) (      7,      0.6) (      8,      0.8)
(      6,      0.25) (      7,      0.5) (      8,      0.75) (      9,      1)

** Output **
< Forward Transform >
ierr =      0
cr[ix][iy]
(      5.5,      0.375) (      -0.575,      0.425) (      -0.5,      -0.075) (      -0.425,      -0.575)
(      -0.586,      0.626) (      0.0297,      -0.0047) (      0.0172,      0.0125) (      0.0047,      0.0297)
(      -0.52,      0.1) (      0.0166,      0.00844) (      0.00406,      0.0125) (      -0.00844,      0.0166)
(      -0.48,      -0.225) (      0.00844,      0.0166) (      -0.00406,      0.0125) (      -0.0166,      0.00844)
(      -0.414,      -0.751) (      -0.0047,      0.0297) (      -0.0172,      0.0125) (      -0.0297,      -0.0047)
< Backward Transform >
ierr =      0
cr[ix][iy]
(      2,      0.05) (      3,      0.1) (      4,      0.15) (      5,      0.2)
(      3,      0.1) (      4,      0.2) (      5,      0.3) (      6,      0.4)
(      4,      0.15) (      5,      0.3) (      6,      0.45) (      7,      0.6)
(      5,      0.2) (      6,      0.4) (      7,      0.6) (      8,      0.8)
(      6,      0.25) (      7,      0.5) (      8,      0.75) (      9,      1)

```

---

## 6.6 TWO-DIMENSIONAL COMPLEX FOURIER TRANSFORM (COMPLEX ARGUMENT TYPE)

### 6.6.1 [DEPRECATED] ASL\_hfc2fb, ASL\_gfc2fb

#### Two-Dimensional Complex Fourier Transform (Including Initialization)

(1) **Function**

**Forward transform**

ASL\_hfc2fb or ASL\_gfc2fb computes the two-dimensional complex Fourier forward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

**Backward transform**

ASL\_hfc2fb or ASL\_gfc2fb computes the two-dimensional complex Fourier backward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

(2) **Usage**

Double precision:

ierr = ASL\_hfc2fb (nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfc2fb (nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	c	$\begin{cases} Z^* \\ C^* \end{cases}$	lx×ly	Input	Input data $c_{k_x,k_y}$ (See Note (b))
				Output	Output results $d_{j_x,j_y}$ (See Notes (b) and (c))
4	lx	I	1	Input	Adjustable dimension of array c (See Note (b))
5	ly	I	1	Input	Second dimension of array c (See Note (b))
6	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	40	Output	Factorization results and number of factors (See Note (d))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y)$	Output	Trigonometric function table (See Note (d))
9	wk	$\begin{cases} Z^* \\ C^* \end{cases}$	lx × ly	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y$
- (c)  $isw \in \{0, 1, -1\}$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ier value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

- (a) When the number of data  $n_x$  or  $n_y$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n_x = 289 (=17^2)$ , it is usually more efficient to set  $n_x = 300 (=2^2 \times 3 \times 5^2)$ ,  $n_x = 320 (=2^6 \times 5)$ ,  $n_x = 384 (=2^7 \times 3)$  or the like.

- (b) The complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) and elements of array  $c$  are associated as follows.

$$c_{k_x, k_y} \leftrightarrow c[k_x + l_x * k_y]$$

Similarly, for the complex data  $d_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ).

**The adjustable dimensions  $l_x$  and  $l_y$  of array  $c$  should be set to odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is even,  $l_x = n_x + 1$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{c}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y} = n_x n_y c_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data ( $n_x, n_y$ ), you should call this function once, and then use the after-initialization transform 6.6.2  $\left\{ \begin{matrix} \text{ASL\_hfc2bf} \\ \text{ASL\_gfc2bf} \end{matrix} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays  $ifax$  and  $trigs$  so they can be used as input to the function 6.6.2  $\left\{ \begin{matrix} \text{ASL\_hfc2bf} \\ \text{ASL\_gfc2bf} \end{matrix} \right\}$ .

To perform initialization only by setting  $isw=0$ , you need not set input data for array  $c$ .

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time

function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

(f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.6.2 (7).

## 6.6.2 [DEPRECATED]ASL\_hfc2bf, ASL\_gfc2bf Two-Dimensional Complex Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_hfc2bf or ASL\_gfc2bf computes the two-dimensional complex Fourier forward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

#### Backward transform

ASL\_hfc2bf or ASL\_gfc2bf computes the two-dimensional complex Fourier backward transform (arbitrary radix) for the two-dimensional complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$d_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} c_{k_x, k_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1)$$

### (2) Usage

Double precision:

ierr = ASL\_hfc2bf (nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfc2bf (nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);



(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	c	$\begin{cases} Z^* \\ C^* \end{cases}$	lx×ly	Input	Input data $c_{k_x,k_y}$ (See Note (b))
				Output	Output results $d_{j_x,j_y}$ (See Notes (b) and (c))
4	lx	I	1	Input	Adjustable dimension of array c (See Note (b))
5	ly	I	1	Input	Second dimension of array c (See Note (b))
6	isw	I	1	Input	Processing switch isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	40	Input	Factorization results and number of factors (See Note (a))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y)$	Input	Trigonometric function table (See Note (a))
9	wk	$\begin{cases} Z^* \\ C^* \end{cases}$	lx × ly	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y$
- (c)  $isw \in \{1, -1\}$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ier value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

(a) This function can be used to repeatedly compute the transform for the same number of data ( $n_x, n_y$ ) after the including-initialization function 6.6.1  $\left\{ \begin{array}{l} \text{ASL\_hfc2fb} \\ \text{ASL\_gfc2fb} \end{array} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.

(b) The complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ) and elements of array c are associated as follows.

$$c_{k_x, k_y} \leftrightarrow c[k_x + l_x * k_y]$$

Similarly, for the complex data  $d_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1$ ).

**The adjustable dimensions  $l_x$  and  $l_y$  of array c should be set to odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is even,  $l_x = n_x + 1$  is set.**

(c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{c}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y} = n_x n_y c_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

(d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c (t - iT)}{\pi (t - iT)}$$

(e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) Example

(a) Problem

Compute the two-dimensional complex Fourier forward and backward transforms using

$$c_{k_x, k_y} = (k_x + 1) + (k_y + 1) + \sqrt{-1} \frac{(k_x + 1)(k_y + 1)}{n_x n_y}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

as input data.

(b) Input data

Array c, nx=5, ny=4, lx=5, ly=5, isw=1(Forward transform), isw=-1 (Backward transform) and nt=2.

(c) Main program

```

/*      C Interface example for ASL_hfc2fb , ASL_hfc2bf */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <complex.h>
#include <asl.h>

int main()
{
    int nx = 5; int ny = 4;
    int lx = 5; int ly = 5;
    double _Complex *c;
    int isw;
    int ifax[40];
    double *trigs;
    double _Complex *wk;
    int nt = 2;
    int ierr;
    int i,j;

    printf( "      *** ASL_hfc2fb , ASL_hfc2bf ***\n" );
    printf( "\n      ** Input **\n" );

    c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lx*ly) ));
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*(nx+ny)) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lx*ly) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    printf( "\tnx = %6d\n", nx );
    printf( "\tny = %6d\n", ny );
    printf( "\tnt = %6d\n", nt );

    for( j=1 ; j<=ny ; j++ )
    {
        for( i=1 ; i<=nx ; i++ )
        {
            c[(i-1)+lx*(j-1)]=(double)(i+j) + (double)(i*j)/(double)(nx*ny) * _Complex_I;
        }
    }

    printf( "\tc[ix][iy]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j]), cimag(c[i+lx*j]) );
        }
        printf( "\n" );
    }
}

```

```

isw = 1;
ierr = ASL_hfc2fb(nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);

for( i=0 ; i<lx*ly ; i++)
{
    c[i] /= (double)(nx*ny);
}

printf( "\n    ** Output **\n" );

printf( "\t< Forward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tc[ix][iy]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j]), cimag(c[i+lx*j]) );
    }
    printf( "\n" );
}

isw = -1;
ierr = ASL_hfc2bf(nx, ny, c, lx, ly, isw, ifax, trigs, wk, nt);

printf( "\t< Backward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tc[ix][iy]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j]), cimag(c[i+lx*j]) );
    }
    printf( "\n" );
}

free( c );
free( trigs );
free( wk );

return 0;
}

```

(d) Output results

```

*** ASL_hfc2fb , ASL_hfc2bf ***

** Input **
nx =      5
ny =      4
nt =      2
c[ix][iy]
(      2,      0.05) (      3,      0.1) (      4,      0.15) (      5,      0.2)
(      3,      0.1) (      4,      0.2) (      5,      0.3) (      6,      0.4)
(      4,      0.15) (      5,      0.3) (      6,      0.45) (      7,      0.6)
(      5,      0.2) (      6,      0.4) (      7,      0.6) (      8,      0.8)
(      6,      0.25) (      7,      0.5) (      8,      0.75) (      9,      1)

** Output **
< Forward Transform >
ierr =      0
c[ix][iy]
(      5.5,      0.375) ( -0.575,      0.425) ( -0.5, -0.075) ( -0.425, -0.575)
( -0.586,      0.626) (      0.0297, -0.0047) (      0.0172,      0.0125) (      0.0047,      0.0297)
( -0.52,      0.1) (      0.0166,      0.00844) (      0.00406,      0.0125) (-0.00844,      0.0166)
( -0.48, -0.225) (      0.00844,      0.0166) (-0.00406,      0.0125) (-0.0166,      0.00844)
( -0.414, -0.751) ( -0.0047,      0.0297) ( -0.0172,      0.0125) ( -0.0297, -0.0047)
< Backward Transform >
ierr =      0
c[ix][iy]
(      2,      0.05) (      3,      0.1) (      4,      0.15) (      5,      0.2)
(      3,      0.1) (      4,      0.2) (      5,      0.3) (      6,      0.4)
(      4,      0.15) (      5,      0.3) (      6,      0.45) (      7,      0.6)
(      5,      0.2) (      6,      0.4) (      7,      0.6) (      8,      0.8)
(      6,      0.25) (      7,      0.5) (      8,      0.75) (      9,      1)

```

---

## 6.7 TWO-DIMENSIONAL REAL FOURIER TRANSFORM

### 6.7.1 [DEPRECATED]ASL\_qfr2fb, ASL\_pfr2fb

#### Two-Dimensional Real Fourier Transform (Including Initialization)

##### (1) Function

###### Forward transform

ASL\_qfr2fb or ASL\_pfr2fb obtains a half period of the two-dimensional Fourier forward transform (arbitrary radix) for the two-dimensional real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$c_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} r_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationships.

$$\begin{aligned} c_{n_x-j_x, n_y-j_y}^* &= c_{j_x, j_y} \\ c_{n_x-j_x, j_y}^* &= c_{j_x, n_y-j_y} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

###### Backward transform

Given the half period  $c_{j_x, j_y}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ) for  $n_x n_y$  complex data  $c_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) satisfying  $c_{n_x-j_x, n_y-j_y}^* = c_{j_x, j_y}$  and  $c_{n_x-j_x, j_y}^* = c_{j_x, n_y-j_y}$ , ASL\_qfr2fb or ASL\_pfr2fb obtains the two-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$\begin{aligned} r_{k_x, k_y} &= \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} c_{j_x, j_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \\ &= \sum_{j_y=0}^{n_y-1} \{c_{0, j_y} + (-1)^{k_x} \hat{c}_{\frac{n_x}{2}, j_y}\} e^{2\pi\sqrt{-1}\frac{j_y k_y}{n_y}} + 2 \sum_{j_y=0}^{n_y-1} \sum_{j_x=1}^{\lfloor \frac{n_x}{2} \rfloor - 1} \Re\{c_{j_x, j_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)}\} \\ &\quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1) \end{aligned}$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  represents the real part of the complex number  $z$ . Also, when  $n_x$  is odd,  $\hat{c}_{\frac{n_x}{2}, j_y} = 0$ , and when  $n_x$  is even,  $\hat{c}_{\frac{n_x}{2}, j_y} = c_{\frac{n_x}{2}, j_y}$ .

##### (2) Usage

Double precision:

ierr = ASL\_qfr2fb (nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfr2fb (nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	r	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Input	Input data $r_{k_x,k_y}$ (Forward transform), or $c_{j_x,j_y}$ (Backward transform) (See Note (b))
				Output	Output results $c_{j_x,j_y}$ (Forward transform), or $r_{k_x,k_y}$ (Backward transform) (See Notes (b) and (c))
4	lx	I	1	Input	Adjustable dimension of array r (See Note (b))
5	ly	I	1	Input	Second dimension of array r (See Note (b))
6	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
7	ifax	I*	40	Output	Factorization results and number of factors (See Note (d))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	nx+2×ny	Output	Trigonometric function table (See Note (d))
9	wk	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b) In the case where  $n_x$  is an odd,  $n_x+1 \leq lx, n_y \leq ly$   
or if  $n_x$  is an even,  $n_x+2 \leq lx, n_y \leq ly$
- (c)  $isw \in \{0, 1, -1\}$
- (d)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) Notes

- (a) When the number of data  $n_x$  or  $n_y$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n_x = 289 (=17^2)$ , it is usually more efficient to set  $n_x = 300 (=2^2 \times 3 \times 5^2)$ ,  $n_x = 320 (=2^6 \times 5)$ ,  $n_x = 384 (=2^7 \times 3)$  or the like.
- (b) The real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) and elements of array  $r$  are associated as follows.

$$r_{k_x, k_y} \leftrightarrow r[k_x + l_x * k_y]$$

When computing the backward transform, if  $n_x (=n_x)$  is odd, then  $r[n_x + l_x * k_y] = 0$ , and when  $n_x$  is even, then  $r[n_x + l_x * k_y] = r[n_x + 1 + l_x * k_y] = 0$ . Also, when entering the real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) into array  $r$ , the corresponding zeros mentioned above need not be specifically stored.

If we let the real and imaginary parts of the complex data  $c_{j_x, j_y}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ) be  $\Re\{c_{j_x, j_y}\}$  and  $\Im\{c_{j_x, j_y}\}$ , respectively, the  $c_{j_x, j_y}$  and elements of array  $r$  are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned} \Re\{c_{j_x, j_y}\} &\leftrightarrow r[2 * j_x + l_x * j_y] \\ \Im\{c_{j_x, j_y}\} &\leftrightarrow r[2 * j_x + 1 + l_x * j_y] \end{aligned}$$

From the properties of a real Fourier transform,  $\Im\{c_{0,0}\} = 0$ , and when  $n_x$  and  $n_y$  are both even,  $\Im\{c_{\frac{n_x}{2}, \frac{n_y}{2}}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array  $r$ , they are considered to be zero when processing is performed. Since the elements  $c_{j_x, j_y}$  ( $j_x = \lfloor \frac{n_x}{2} \rfloor + 1, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) can be obtained according to the following relationships from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$\begin{aligned} c_{n_x - j_x, n_y - j_y}^* &= c_{j_x, j_y} \\ c_{n_x - j_x, j_y}^* &= c_{j_x, n_y - j_y} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ . **The adjustable dimensions of array  $r$  should be set so that  $l_x/2$  and  $l_y$  are odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is (a multiple of 4)+2,  $l_x = n_x + 4$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{r}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{r}_{k_x, k_y} = n_x n_y r_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data ( $n_x, n_y$ ), you should call this function once, and then use the after-initialization transform 6.7.2  $\left\{ \begin{array}{l} \text{ASL\_qfr2fb} \\ \text{ASL\_pfr2fb} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays ifax and trigs so they can be used as input to the function 6.7.2  $\left\{ \begin{array}{l} \text{ASL\_qfr2fb} \\ \text{ASL\_pfr2fb} \end{array} \right\}$ .

To perform initialization only by setting isw=0, you need not set input data for array r.

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.7.2 (7).



## 6.7.2 [DEPRECATED]ASL\_qfr2bf, ASL\_pfr2bf Two-Dimensional Real Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_qfr2bf or ASL\_pfr2bf obtains a half period of the two-dimensional Fourier forward transform (arbitrary radix) for the two-dimensional real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ).

$$c_{j_x, j_y} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} r_{k_x, k_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \quad (j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationships.

$$\begin{aligned} c_{n_x-j_x, n_y-j_y}^* &= c_{j_x, j_y} \\ c_{n_x-j_x, j_y}^* &= c_{j_x, n_y-j_y} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

#### Backward transform

Given the half period  $c_{j_x, j_y}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ) for  $n_x n_y$  complex data  $c_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) satisfying  $c_{n_x-j_x, n_y-j_y}^* = c_{j_x, j_y}$  and  $c_{n_x-j_x, j_y}^* = c_{j_x, n_y-j_y}$ , ASL\_qfr2bf or ASL\_pfr2bf obtains the two-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$\begin{aligned} r_{k_x, k_y} &= \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} c_{j_x, j_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \\ &= \sum_{j_y=0}^{n_y-1} \{c_{0, j_y} + (-1)^{k_x} \hat{c}_{\frac{n_x}{2}, j_y}\} e^{2\pi\sqrt{-1}\frac{j_y k_y}{n_y}} + 2 \sum_{j_y=0}^{n_y-1} \sum_{j_x=1}^{\lfloor \frac{n_x}{2} \rfloor - 1} \Re\{c_{j_x, j_y} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)}\} \\ &\quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1) \end{aligned}$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  represents the real part of the complex number  $z$ . Also, when  $n_x$  is odd,  $\hat{c}_{\frac{n_x}{2}, j_y} = 0$ , and when  $n_x$  is even,  $\hat{c}_{\frac{n_x}{2}, j_y} = c_{\frac{n_x}{2}, j_y}$ .

### (2) Usage

Double precision:

ierr = ASL\_qfr2bf (nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfr2bf (nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	r	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Input	Input data $r_{k_x,k_y}$ (Forward transform), or $c_{j_x,j_y}$ (Backward transform) (See Note (b))
				Output	Output results $c_{j_x,j_y}$ (Forward transform), or $r_{k_x,k_y}$ (Backward transform) (See Notes (b) and (c))
4	lx	I	1	Input	Adjustable dimension of array r (See Note (b))
5	ly	I	1	Input	Second dimension of array r (See Note (b))
6	isw	I	1	Input	Processing switch isw= 1: Forward transform isw=-1: Backward transform
7	ifax	I*	40	Input	Factorization results and number of factors (See Note (a))
8	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	nx+2×ny	Input	Trigonometric function table (See Note (a))
9	wk	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly	Work	Work area
10	nt	I	1	Input	Number of tasks to be generated
11	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2$
- (b) In the case where  $n_x$  is an odd,  $n_x+1 \leq lx, n_y \leq ly$   
or if  $n_x$  is an even,  $n_x+2 \leq lx, n_y \leq ly$
- (c)  $isw \in \{1, -1\}$
- (d)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) Notes

- (a) This function can be used to repeatedly compute the transform for the same number of data (nx, ny) after the including-initialization function 6.7.1  $\left\{ \begin{matrix} \text{ASL\_qfr2fb} \\ \text{ASL\_pfr2fb} \end{matrix} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.
- (b) The real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) and elements of array r are associated as follows.

$$r_{k_x, k_y} \leftrightarrow r[k_x + l_x * k_y]$$

When computing the backward transform, if  $n_x(=n_x)$  is odd, then  $r[n_x + l_x * k_y] = 0$ , and when  $n_x$  is even, then  $r[n_x + l_x * k_y] = r[n_x + 1 + l_x * k_y] = 0$ . Also, when entering the real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) into array r, the corresponding zeros mentioned above need not be specifically stored.

If we let the real and imaginary parts of the complex data  $c_{j_x, j_y}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ) be  $\Re\{c_{j_x, j_y}\}$  and  $\Im\{c_{j_x, j_y}\}$ , respectively, the  $c_{j_x, j_y}$  and elements of array r are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned} \Re\{c_{j_x, j_y}\} &\leftrightarrow r[2 * j_x + l_x * j_y] \\ \Im\{c_{j_x, j_y}\} &\leftrightarrow r[2 * j_x + 1 + l_x * j_y] \end{aligned}$$

From the properties of a real Fourier transform,  $\Im\{c_{0,0}\} = 0$ , and when  $n_x$  and  $n_y$  are both even,  $\Im\{c_{\frac{n_x}{2}, \frac{n_y}{2}}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array r, they are considered to be zero when processing is performed. Since the elements  $c_{j_x, j_y}$  ( $j_x = \lfloor \frac{n_x}{2} \rfloor + 1, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) can be obtained according to the following relationships from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$\begin{aligned} c_{n_x - j_x, n_y - j_y}^* &= c_{j_x, j_y} \\ c_{n_x - j_x, j_y}^* &= c_{j_x, n_y - j_y} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ . **The adjustable dimensions of array r should be set so that  $l_x/2$  and  $l_y$  are odd numbers to avoid bank conflict of main memory. Usually, when  $n_x$ , for example, is (a multiple of 4)+2,  $l_x=n_x+4$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ) be represented by  $\hat{r}_{k_x, k_y}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ), then the following relationship holds.

$$\hat{r}_{k_x, k_y} = n_x n_y r_{k_x, k_y} \quad (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

- (a) Problem

Compute the two-dimensional real Fourier forward and backward transforms using

$$r_{k_x, k_y} = \frac{n_x + n_y}{(k_x + 1) + (k_y + 1)}$$

$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1)$

as input data.

- (b) Input data

Array r, nx=6, ny=4, lx=10, ly=5, isw=1(Forward transform), isw=-1 (Backward transform) and nt=2.

- (c) Main program

```
/*      C Interface example for ASL_qfr2fb , ASL_qfr2bf */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int nx = 6;    int ny = 4;
    int lx = 10;   int ly = 5;
    double *r;
    int isw;
    int ifax[40];
    double *trigs;
    double *wk;
    int nt = 2;
    int ierr;
    int i,j;

    printf( "      *** ASL_qfr2fb , ASL_qfr2bf ***\n" );
    printf( "\n      ** Input **\n" );

    r = ( double * )malloc((size_t)( sizeof(double) * (lx*ly) ));
    if( r == NULL )
    {
        printf( "no enough memory for array r\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (nx+2*ny) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }
}
```

```

wk = ( double * )malloc((size_t)( sizeof(double) * (lx*ly) ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

printf( "\tnx = %6d\n", nx );
printf( "\tny = %6d\n", ny );
printf( "\tnt = %6d\n", nt );

for( j=1 ; j<=ny ; j++ )
{
    for( i=1 ; i<=nx ; i++ )
    {
        r[i-1+lx*(j-1)]=(double)(nx+ny)/(double)(i+j) ;
    }
}

printf( "\tr[ix][iy]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j] );
    }
    printf( "\n" );
}

isw = 1;
ierr = ASL_qfr2fb(nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

for( i=0 ; i<lx*ly ; i++)
{
    r[i] /= (double)(nx*ny);
}

printf( "\n    ** Output **\n" );
printf( "\t< Forward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tr[ix][iy]\n" );
for( i=0 ; i<nx+2 ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j] );
    }
    printf( "\n" );
}

isw = -1;
ierr = ASL_qfr2bf(nx, ny, r, lx, ly, isw, ifax, trigs, wk, nt);

printf( "\t< Backward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tr[ix][iy]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j] );
    }
    printf( "\n" );
}

free( r );
free( trigs );
free( wk );

return 0;
}

```

(d) Output results

```

*** ASL_qfr2fb , ASL_qfr2bf ***

** Input **
nx =      6
ny =      4
nt =      2
r[ix][iy]
      5          3.33          2.5          2
3.33          2.5          2          1.67
2.5          2          1.67          1.43

```

```

      2      1.67      1.43      1.25
    1.67      1.43      1.25      1.11
    1.43      1.25      1.11      1

** Output **
< Forward Transform >
ierr = 0
r[ix][iy]
  1.94      0.249      0.219      0.249
    0      -0.155      0      0.155
  0.296      0.0585      0.0761      0.119
 -0.247      -0.0939      -0.0447      -0.00945
  0.229      0.0557      0.058      0.0794
 -0.0928      -0.0535      -0.0186      0.0102
  0.219      0.0637      0.0547      0.0637
    0      -0.0301      0      0.0301
< Backward Transform >
ierr = 0
r[ix][iy]
    5      3.33      2.5      2
   3.33      2.5      2      1.67
   2.5      2      1.67      1.43
    2      1.67      1.43      1.25
   1.67      1.43      1.25      1.11
   1.43      1.25      1.11      1

```

---

## 6.8 THREE-DIMENSIONAL COMPLEX FOURIER TRANSFORM (REAL ARGUMENT TYPE)

### 6.8.1 [DEPRECATED]ASL\_qfc3fb, ASL\_pfc3fb

#### Three-Dimensional Complex Fourier Transform (Including Initialization)

##### (1) Function

ASL\_qfc3fb or ASL\_pfc3fb computes the three-dimensional complex Fourier forward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

##### Backward transform

ASL\_qfc3fb or ASL\_pfc3fb computes the three-dimensional complex Fourier backward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

##### (2) Usage

Double precision:

ierr = ASL\_qfc3fb (nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfc3fb (nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	cr	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly×lz	Input	Real part of input data $c_{k_x, k_y, k_z}$ (See Note (b))
				Output	Real part of output data $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
5	ci	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly×lz	Input	Imaginary part of input data $c_{k, s, u}$ (See Note (b))
				Output	Imaginary part of output results $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
6	lx	I	1	Input	Adjustable dimension of array cr and ci (See Note (b))
7	ly	I	1	Input	Second dimension of array cr and ci (See Note (b))
8	lz	I	1	Input	Third dimension of array cr and ci (See Note (b))
9	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
10	ifax	I*	60	Output	Factorization results and number of factors (See Note (d))
11	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y + n_z)$	Output	Trigonometric function table (See Note (d))
12	wk	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times l_x \times l_y \times l_z$	Work	Work area
13	nt	I	1	Input	Number of tasks to be generated
14	ierr	I	1	Output	Error indicator (Return Value)



(4) **Restrictions**

- (a)  $nx \geq 2$ ,  $ny \geq 2$  or  $nz \geq 2$
- (b)  $nx \leq lx$ ,  $ny \leq ly$ , or  $nz \leq lz$
- (c)  $isw \in \{0, 1, -1\}$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

- (a) When the number of data  $nx$ ,  $ny$  or  $nz$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $nx = 289 (=17^2)$ , it is usually more efficient to set  $nx = 300 (=2^2 \times 3 \times 5^2)$ ,  $nx = 320 (=2^6 \times 5)$ ,  $nx = 384 (=2^7 \times 3)$  or the like.
- (b) If we let the real and imaginary parts of the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be  $\Re\{c_{k_x, k_y, k_z}\}$  and  $\Im\{c_{k_x, k_y, k_z}\}$ , respectively, the  $c_{k_x, k_y, k_z}$  and elements of arrays  $cr$  and  $ci$  are associated as follows.

$$\begin{aligned} \Re\{c_{k_x, k_y, k_z}\} &\leftrightarrow cr[k_x + lx * (k_y + ly * k_z)] \\ \Im\{c_{k_x, k_y, k_z}\} &\leftrightarrow ci[k_x + lx * (k_y + ly * k_z)] \end{aligned}$$

Similarly, for the complex data  $d_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ).

**The adjustable dimensions  $lx$ ,  $ly$ , and  $lz$  of arrays  $cr$  and  $ci$  should be set to odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within arrays  $cr$  and  $ci$ . Usually, when  $nx$ , for example, is even,  $lx=nx+1$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be represented by  $\hat{c}_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ), then the following relationship holds.

$$\begin{aligned} \hat{c}_{k_x, k_y, k_z} &= n_x n_y n_z c_{k_x, k_y, k_z} \\ &(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1, k_z = 0, \dots, n_z - 1) \end{aligned}$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

(d) To repeatedly compute the transform for the same number of data ( $n_x, n_y, n_z$ ), you should call this function once, and then use the after-initialization transform 6.8.2  $\left\{ \begin{array}{l} \text{ASL\_qfc3fb} \\ \text{ASL\_pfc3fb} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays ifax and trigs so they can be used as input to the function 6.8.2  $\left\{ \begin{array}{l} \text{ASL\_qfc3fb} \\ \text{ASL\_pfc3fb} \end{array} \right\}$ .

To perform initialization only by setting isw=0, you need not set input data for arrays cr and ci.

(e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c (t - iT)}{\pi(t - iT)}$$

(f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.8.2 (7).

## 6.8.2 [DEPRECATED]ASL\_qfc3bf, ASL\_pfc3bf Three-Dimensional Complex Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_qfc3bf or ASL\_pfc3bf computes the three-dimensional complex Fourier forward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

#### Backward transform

ASL\_qfc3bf or ASL\_pfc3bf computes the three-dimensional complex Fourier backward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

### (2) Usage

Double precision:

ierr = ASL\_qfc3bf (nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfc3bf (nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	cr	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly×lz	Input	Real part of input data $c_{k_x, k_y, k_z}$ (See Note (b))
				Output	Real part of output data $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
5	ci	$\begin{cases} D^* \\ R^* \end{cases}$	lx×ly×lz	Input	Imaginary part of input data $c_{k_x, k_y, k_z}$ (See Note (b))
				Output	Imaginary part of output results $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
6	lx	I	1	Input	Adjustable dimension of array cr and ci (See Note (b))
7	ly	I	1	Input	Second dimension of array cr and ci (See Note (b))
8	lz	I	1	Input	Third dimension of array cr and ci (See Note (b))
9	isw	I	1	Input	Processing switch isw= 1 :Forward transform isw=-1 :Backward transform
10	ifax	I*	60	Input	Factorization results and number of factors (See Note (d))
11	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y + n_z)$	Input	Trigonometric function table (See Note (d))
12	wk	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times l_x \times l_y \times l_z$	Work	Work area
13	nt	I	1	Input	Number of tasks to be generated
14	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2, n_z \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y, n_z \leq l_z$
- (c)  $isw \in \{1, -1\}$
- (d)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) Notes

(a) This function can be used to repeatedly compute the transform for the same number of data (nx, ny, nz) after the including-initialization function 6.8.1  $\left\{ \begin{matrix} \text{ASL\_qfc3fb} \\ \text{ASL\_pfc3fb} \end{matrix} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.

(b) If we let the real and imaginary parts of the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be  $\Re\{c_{k_x, k_y, k_z}\}$  and  $\Im\{c_{k_x, k_y, k_z}\}$ , respectively, the  $c_{k_x, k_y, k_z}$  and elements of arrays cr and ci are associated as follows.

$$\begin{aligned} \Re\{c_{k_x, k_y, k_z}\} &\leftrightarrow \text{cr}[k_x + l_x * (k_y + l_y * k_z)] \\ \Im\{c_{k_x, k_y, k_z}\} &\leftrightarrow \text{ci}[k_x + l_x * (k_y + l_y * k_z)] \end{aligned}$$

Similarly, for the complex data  $d_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ).

**The adjustable dimensions lx, ly, and lz of arrays cr and ci should be set to odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within arrays cr and ci. Usually, when nx, for example, is even, lx=nx+1 is set.**

(c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be represented by  $\hat{c}_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ), then the following relationship holds.

$$\begin{aligned} \hat{c}_{k_x, k_y, k_z} &= n_x n_y n_z c_{k_x, k_y, k_z} \\ &(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1, k_z = 0, \dots, n_z - 1) \end{aligned}$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

(d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

- (a) Problem

Compute the three-dimensional complex Fourier forward and backward transforms using

$$C_{k_x, k_y, k_z} = \frac{n_x + n_y + n_z}{(k_x + 1) + (k_y + 1) + (k_z + 1)} + \sqrt{-1} \frac{(k_x + 1)(k_y + 1)(k_z + 1)}{n_x n_y n_z}$$

$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$

as input data.

- (b) Input data

Array cr and ci, nx=5, ny=4, nz=3, lx=5, ly=5, lz=3, isw=1 (Forward transform), isw=-1 (Backward transform) and nt=2.

- (c) Main program

```

/*      C Interface example for ASL_qfc3fb , ASL_qfc3bf */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    int nx = 5; int ny = 4; int nz = 3;
    int lx = 5; int ly = 5; int lz = 3;
    double *cr; double *ci;
    int isw;
    int ifax[60];
    double *trigs;
    double *wk;
    int nt = 2;
    int ierr;
    int i,j,k;

    printf( "      *** ASL_qfc3fb , ASL_qfc3bf ***\n" );
    printf( "\n      ** Input **\n" );

    cr = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*lz) ));
    if( cr == NULL )
    {
        printf( "no enough memory for array cr\n" );
        return -1;
    }

    ci = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*lz) ));
    if( ci == NULL )
    {
        printf( "no enough memory for array ci\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*(nx+ny+nz)) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double * )malloc((size_t)( sizeof(double) * (2*lx*ly*lz) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    printf( "\tnx = %6d\n", nx );
    printf( "\tny = %6d\n", ny );
    printf( "\tnz = %6d\n", nz );
    printf( "\tnt = %6d\n", nt );

    for( k=1 ; k<=nz ; k++ )
    {
        for( j=1 ; j<=ny ; j++ )
        {
            for( i=1 ; i<=nx ; i++ )

```

```

                cr[i-1+lx*(j-1)+lx*ly*(k-1)]=(double)(nx+ny+nz)/(double)(i+j+k) ;
                ci[i-1+lx*(j-1)+lx*ly*(k-1)]=(double)(i*j*k)/(double)(nx*ny*nz) ;
            }
        }
    printf( "\tcr[ix][iy][1] ci[ix][iy][1]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j          ], ci[i+lx*j          ] );
        }
        printf( "\n" );
    }

    printf( "\tcr[ix][iy][2] ci[ix][iy][2]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*1], ci[i+lx*j+lx*ly*1] );
        }
        printf( "\n" );
    }

    printf( "\tcr[ix][iy][3] ci[ix][iy][3]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*2], ci[i+lx*j+lx*ly*2] );
        }
        printf( "\n" );
    }

    isw = 1;
    ierr = ASL_qfc3fb(nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

    for( i=0 ; i<lx*ly*lz ; i++ )
    {
        cr[i] /= (double)(nx*ny*nz);
        ci[i] /= (double)(nx*ny*nz);
    }

    printf( "\n    ** Output **\n" );
    printf( "\t< Forward Transform >\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\tcr[ix][iy][1] ci[ix][iy][1]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j          ], ci[i+lx*j          ] );
        }
        printf( "\n" );
    }

    printf( "\tcr[ix][iy][2] ci[ix][iy][2]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*1], ci[i+lx*j+lx*ly*1] );
        }
        printf( "\n" );
    }

    printf( "\tcr[ix][iy][3] ci[ix][iy][3]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*2], ci[i+lx*j+lx*ly*2] );
        }
        printf( "\n" );
    }

    isw = -1;
    ierr = ASL_qfc3bf(nx, ny, nz, cr, ci, lx, ly, lz, isw, ifax, trigs, wk, nt);

    printf( "\t< Backward Transform >\n" );
    printf( "\tierr = %6d\n", ierr );

    printf( "\tcr[ix][iy][1] ci[ix][iy][1]\n" );

```

```

for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", cr[i+lx*j      ], ci[i+lx*j      ] );
    }
    printf( "\n" );
}

printf( "\tcr[ix][iy][2] ci[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*1], ci[i+lx*j+lx*ly*1] );
    }
    printf( "\n" );
}

printf( "\tcr[ix][iy][3] ci[ix][iy][3]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", cr[i+lx*j+lx*ly*2], ci[i+lx*j+lx*ly*2] );
    }
    printf( "\n" );
}

free( cr );
free( ci );
free( trigs );
free( wk );

return 0;
}
    
```

(d) Output results

```

*** ASL_qfc3fb , ASL_qfc3bf ***

** Input **
nx = 5
ny = 4
nz = 3
nt = 2
cr[ix][iy][1] ci[ix][iy][1]
( 4, 0.0167) ( 3, 0.0333) ( 2.4, 0.05) ( 2, 0.0667)
( 3, 0.0333) ( 2.4, 0.0667) ( 2, 0.1) ( 1.71, 0.133)
( 2.4, 0.05) ( 2, 0.1) ( 1.71, 0.15) ( 1.5, 0.2)
( 2, 0.0667) ( 1.71, 0.133) ( 1.5, 0.2) ( 1.33, 0.267)
( 1.71, 0.0833) ( 1.5, 0.167) ( 1.33, 0.25) ( 1.2, 0.333)
cr[ix][iy][2] ci[ix][iy][2]
( 3, 0.0333) ( 2.4, 0.0667) ( 2, 0.1) ( 1.71, 0.133)
( 2.4, 0.0667) ( 2, 0.133) ( 1.71, 0.2) ( 1.5, 0.267)
( 2, 0.1) ( 1.71, 0.2) ( 1.5, 0.3) ( 1.33, 0.4)
( 1.71, 0.133) ( 1.5, 0.267) ( 1.33, 0.4) ( 1.2, 0.533)
( 1.5, 0.167) ( 1.33, 0.333) ( 1.2, 0.5) ( 1.09, 0.667)
cr[ix][iy][3] ci[ix][iy][3]
( 2.4, 0.05) ( 2, 0.1) ( 1.71, 0.15) ( 1.5, 0.2)
( 2, 0.1) ( 1.71, 0.2) ( 1.5, 0.3) ( 1.33, 0.4)
( 1.71, 0.15) ( 1.5, 0.3) ( 1.33, 0.45) ( 1.2, 0.6)
( 1.5, 0.2) ( 1.33, 0.4) ( 1.2, 0.6) ( 1.09, 0.8)
( 1.33, 0.25) ( 1.2, 0.5) ( 1.09, 0.75) ( 1, 1)

** Output **
< Forward Transform >
ierr = 0
cr[ix][iy][1] ci[ix][iy][1]
( 1.74, 0.25) ( 0.102, -0.16) ( 0.137, -0.05) ( 0.202, 0.06)
( 0.108, -0.189) ( 0.0379, -0.0469) ( 0.0406, -0.0125) ( 0.0525, 0.016)
( 0.125, -0.0784) ( 0.034, -0.0168) ( 0.0261, 0.00288) ( 0.0254, 0.0209)
( 0.152, -0.00492) ( 0.0366, 0.00116) ( 0.0207, 0.0138) ( 0.012, 0.028)
( 0.223, 0.106) ( 0.0462, 0.0236) ( 0.0177, 0.0292) (-0.00166, 0.0406)
cr[ix][iy][2] ci[ix][iy][2]
( 0.106, -0.127) ( 0.0407, -0.0223) ( 0.0315, 0.00295) ( 0.0297, 0.0255)
( 0.0419, -0.0317) (-0.00167, -0.00877) ( 0.0025, -0.00799) ( 0.00901, -0.00976)
( 0.0317, -0.00698) ( 0.00134, -0.00743) ( 0.00423, -0.00524) ( 0.00924, -0.00424)
( 0.0297, 0.0084) ( 0.00473, -0.00711) ( 0.00655, -0.00336) ( 0.0108, 0.0001)
( 0.0318, 0.0285) ( 0.0112, -0.00921) ( 0.0118, -0.00179) ( 0.016, 0.00627)
cr[ix][iy][3] ci[ix][iy][3]
( 0.178, 0.00231) ( 0.0403, 0.014) ( 0.017, 0.022) ( 0.00125, 0.0329)
( 0.0484, 0.00885) ( 0.00516, -0.0104) ( 0.00849, -0.00569) ( 0.0153, -0.0016)
( 0.0244, 0.0163) ( 0.00692, -0.0061) ( 0.0076, -0.00159) ( 0.0107, 0.00322)
( 0.0129, 0.0239) ( 0.00961, -0.0029) ( 0.00799, 0.00185) ( 0.00849, 0.0078)
( 0.00117, 0.036) ( 0.0163, 0.000768) ( 0.0106, 0.00714) ( 0.00733, 0.0161)
< Backward Transform >
ierr = 0
cr[ix][iy][1] ci[ix][iy][1]
    
```



```

(      4, 0.0167) (      3, 0.0333) (      2.4, 0.05) (      2, 0.0667)
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)
(      2, 0.0667) (      1.71, 0.133) (      1.5, 0.2) (      1.33, 0.267)
(      1.71, 0.0833) (      1.5, 0.167) (      1.33, 0.25) (      1.2, 0.333)
cr[ix][iy][2] ci[ix][iy][2]
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.0667) (      2, 0.133) (      1.71, 0.2) (      1.5, 0.267)
(      2, 0.1) (      1.71, 0.2) (      1.5, 0.3) (      1.33, 0.4)
(      1.71, 0.133) (      1.5, 0.267) (      1.33, 0.4) (      1.2, 0.533)
(      1.5, 0.167) (      1.33, 0.333) (      1.2, 0.5) (      1.09, 0.667)
cr[ix][iy][3] ci[ix][iy][3]
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)
(      2, 0.1) (      1.71, 0.2) (      1.5, 0.3) (      1.33, 0.4)
(      1.71, 0.15) (      1.5, 0.3) (      1.33, 0.45) (      1.2, 0.6)
(      1.5, 0.2) (      1.33, 0.4) (      1.2, 0.6) (      1.09, 0.8)
(      1.33, 0.25) (      1.2, 0.5) (      1.09, 0.75) (      1, 1)
    
```

---

## 6.9 THREE-DIMENSIONAL COMPLEX FOURIER TRANSFORM (COMPLEX ARGUMENT TYPE)

### 6.9.1 [DEPRECATED] ASL\_hfc3fb, ASL\_gfc3fb

#### Three-Dimensional Complex Fourier Transform (Including Initialization)

##### (1) Function

###### Forward transform

ASL\_hfc3fb or ASL\_gfc3fb computes the three-dimensional complex Fourier forward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

###### Backward transform

ASL\_hfc3fb or ASL\_gfc3fb computes the three-dimensional complex Fourier backward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

##### (2) Usage

Double precision:

ierr = ASL\_hfc3fb (nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfc3fb (nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	c	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	lx×ly×lz	Input	Input data $c_{k_x, k_y, k_z}$ (See Note (b))
				Output	Output results $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array c (See Note (b))
6	ly	I	1	Input	Second dimension of array c (See Note (b))
7	lz	I	1	Input	Third dimension of array c (See Note (b))
8	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
9	ifax	I*	60	Output	Factorization results and number of factors (See Note (d))
10	trigs	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	$2 \times (n_x + n_y + n_z)$	Output	Trigonometric function table (See Note (d))
11	wk	$\begin{Bmatrix} Z^* \\ C^* \end{Bmatrix}$	lx × ly × lz	Work	Work area
12	nt	I	1	Input	Number of tasks to be generated
13	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2, n_z \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y, n_z \leq l_z$
- (c)  $isw \in \{0, 1, -1\}$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

(a) When the number of data  $n_x$ ,  $n_y$  or  $n_z$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $n_x = 289 (=17^2)$ , it is usually more efficient to set  $n_x = 300 (=2^2 \times 3 \times 5^2)$ ,  $n_x = 320 (=2^6 \times 5)$ ,  $n_x = 384 (=2^7 \times 3)$  or the like.

(b) The complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) and elements of array  $c$  are associated as follows.

$$c_{k_x, k_y, k_z} \leftrightarrow c[k_x + l_x * (k_y + l_y * k_z)]$$

Similarly, for the complex data  $d_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ).

**The adjustable dimensions  $l_x$ ,  $l_y$ , and  $l_z$  of array  $c$  should be set to odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within array  $c$ . Usually, when  $n_x$ , for example, is even,  $l_x = n_x + 1$  is set.**

(c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be represented by  $\hat{c}_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y, k_z} = n_x n_y n_z c_{k_x, k_y, k_z} \\
(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1, k_z = 0, \dots, n_z - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

(d) To repeatedly compute the transform for the same number of data ( $n_x$ ,  $n_y$ ,  $n_z$ ), you should call this function once, and then use the after-initialization transform 6.9.2  $\left\{ \begin{matrix} \text{ASL\_hfc3bf} \\ \text{ASL\_gfc3bf} \end{matrix} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays  $ifax$  and  $trigs$  so they can be used as input to the function 6.9.2  $\left\{ \begin{matrix} \text{ASL\_hfc3bf} \\ \text{ASL\_gfc3bf} \end{matrix} \right\}$ .

To perform initialization only by setting  $isw=0$ , you need not set input data for array  $c$ .

(e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period,

the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

(f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.9.2 (7).

## 6.9.2 [DEPRECATED]ASL\_hfc3bf, ASL\_gfc3bf Three-Dimensional Complex Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_hfc3bf or ASL\_gfc3bf computes the three-dimensional complex Fourier forward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

#### Backward transform

ASL\_hfc3bf or ASL\_gfc3bf computes the three-dimensional complex Fourier backward transform (arbitrary radix) for the three-dimensional complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$d_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} c_{k_x, k_y, k_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$(j_x = 0, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$

### (2) Usage

Double precision:

ierr = ASL\_hfc3bf (nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_gfc3bf (nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	c	$\begin{cases} Z^* \\ C^* \end{cases}$	lx×ly×lz	Input	Input data $c_{k_x, k_y, k_z}$ (See Note (b))
				Output	Output results $d_{j_x, j_y, j_z}$ (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array c (See Note (b))
6	ly	I	1	Input	Second dimension of array c (See Note (b))
7	lz	I	1	Input	Third dimension of array c (See Note (b))
8	isw	I	1	Input	Processing switch isw= 1 :Forward transform isw=-1 :Backward transform
9	ifax	I*	60	Input	Factorization results and number of factors (See Note (a))
10	trigs	$\begin{cases} D^* \\ R^* \end{cases}$	$2 \times (n_x + n_y + n_z)$	Input	Trigonometric function table (See Note (a))
11	wk	$\begin{cases} Z^* \\ C^* \end{cases}$	lx × ly × lz	Work	Work area
12	nt	I	1	Input	Number of tasks to be generated
13	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n_x \geq 2, n_y \geq 2, n_z \geq 2$
- (b)  $n_x \leq l_x, n_y \leq l_y, n_z \leq l_z$
- (c)  $isw \in \{1, -1\}$
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	

(6) **Notes**

- (a) This function can be used to repeatedly compute the transform for the same number of data ( $n_x, n_y, n_z$ ) after the including-initialization function 6.9.1  $\left\{ \begin{array}{l} \text{ASL\_hfc3bf} \\ \text{ASL\_gfc3bf} \end{array} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.
- (b) The complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) and elements of array c are associated as follows.

$$c_{k_x, k_y, k_z} \leftrightarrow c[k_x + l_x * (k_y + l_y * k_z)]$$

**The adjustable dimensions  $l_x, l_y,$  and  $l_z$  of array c should be set to odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within array c. Usually, when  $n_x,$  for example, is even,  $l_x = n_x + 1$  is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the complex data  $c_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) be represented by  $\hat{c}_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ), then the following relationship holds.

$$\hat{c}_{k_x, k_y, k_z} = n_x n_y n_z c_{k_x, k_y, k_z} \\ (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1, k_z = 0, \dots, n_z - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c (t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.



(7) Example

(a) Problem

Compute the three-dimensional complex Fourier forward and backward transforms using

$$c_{k_x, k_y, k_z} = \frac{n_x + n_y + n_z}{(k_x + 1) + (k_y + 1) + (k_z + 1)} + \sqrt{-1} \frac{(k_x + 1)(k_y + 1)(k_z + 1)}{n_x n_y n_z}$$

( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ )

as input data.

(b) Input data

Array cr and ci, nx=5, ny=4, nz=3, lx=5, ly=5, lz=3, isw=1 (Forward transform), isw=-1 (Backward transform) and nt=2.

(c) Main program

```

/*      C Interface example for ASL_hfc3fb , ASL_hfc3bf */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <complex.h>
#include <asl.h>

int main()
{
    int nx = 5; int ny = 4; int nz = 3;
    int lx = 5; int ly = 5; int lz = 3;
    double _Complex *c;
    int isw;
    int ifax[60];
    double *trigs;
    double _Complex *wk;
    int nt = 2;
    int ierr;
    int i,j,k;

    printf( "      *** ASL_hfc3fb , ASL_hfc3bf ***\n" );
    printf( "\n      ** Input **\n" );

    c = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lx*ly*lz) ));
    if( c == NULL )
    {
        printf( "no enough memory for array c\n" );
        return -1;
    }

    trigs = ( double * )malloc((size_t)( sizeof(double) * (2*(nx+ny+nz)) ));
    if( trigs == NULL )
    {
        printf( "no enough memory for array trigs\n" );
        return -1;
    }

    wk = ( double _Complex * )malloc((size_t)( sizeof(double _Complex) * (lx*ly*lz) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    printf( "\tnx = %6d\n", nx );
    printf( "\tny = %6d\n", ny );
    printf( "\tnz = %6d\n", nz );
    printf( "\tnt = %6d\n", nt );

    for( k=1 ; k<=nz ; k++ )
    {
        for( j=1 ; j<=ny ; j++ )
        {
            for( i=1 ; i<=nx ; i++ )
            {
                c[(i-1)+lx*(j-1)+lx*ly*(k-1)]=(double)(nx+ny+nz)/(double)(i+j+k)
                    +(double)(i*j*k)/(double)(nx*ny*nz) * _Complex_I;
            }
        }
    }

    printf( "\tc[ix][iy][1]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {

```

```

        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j      ]), cimag(c[i+lx*j      ] ) );
    }
    printf( "\n" );
}

printf( "\tc[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*1]), cimag(c[i+lx*j+lx*ly*1]) );
    }
    printf( "\n" );
}

printf( "\tc[ix][iy][3]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*2]), cimag(c[i+lx*j+lx*ly*2]) );
    }
    printf( "\n" );
}

isw = 1;
ierr = ASL_hfc3fb(nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);
for( i=0 ; i<lx*ly*lz ; i++ )
{
    c[i] /= (double)(nx*ny*nz);
}

printf( "\n    ** Output **\n" );
printf( "\t< Forward Transform >\n" );
printf( "\t(ierr = %6d\n", ierr );

printf( "\tc[ix][iy][1]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j      ]), cimag(c[i+lx*j      ] ) );
    }
    printf( "\n" );
}

printf( "\tc[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*1]), cimag(c[i+lx*j+lx*ly*1]) );
    }
    printf( "\n" );
}

printf( "\tc[ix][iy][3]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*2]), cimag(c[i+lx*j+lx*ly*2]) );
    }
    printf( "\n" );
}

isw = -1;
ierr = ASL_hfc3bf(nx, ny, nz, c, lx, ly, lz, isw, ifax, trigs, wk, nt);
printf( "\t< Backward Transform >\n" );
printf( "\t(ierr = %6d\n", ierr );

printf( "\tc[ix][iy][1]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j      ]), cimag(c[i+lx*j      ] ) );
    }
    printf( "\n" );
}

printf( "\tc[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{

```

```

        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*1]), cimag(c[i+lx*j+lx*ly*1]) );
        }
        printf( "\n" );
    }

    printf( "\tc[ix][iy][3]\n" );
    for( i=0 ; i<nx ; i++ )
    {
        for( j=0 ; j<ny ; j++ )
        {
            printf( "\t(%8.3g,%8.3g)", creal(c[i+lx*j+lx*ly*2]), cimag(c[i+lx*j+lx*ly*2]) );
        }
        printf( "\n" );
    }

    free( c );
    free( trigs );
    free( wk );

    return 0;
}

```

(d) Output results

```

*** ASL_hfc3fb , ASL_hfc3bf ***

** Input **
nx =      5
ny =      4
nz =      3
nt =      2
c[ix][iy][1]
(      4, 0.0167) (      3, 0.0333) (      2.4, 0.05) (      2, 0.0667)
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)
(      2, 0.0667) (      1.71, 0.133) (      1.5, 0.2) (      1.33, 0.267)
(      1.71, 0.0833) (      1.5, 0.167) (      1.33, 0.25) (      1.2, 0.333)
c[ix][iy][2]
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.0667) (      2, 0.133) (      1.71, 0.2) (      1.5, 0.267)
(      2, 0.1) (      1.71, 0.2) (      1.5, 0.3) (      1.33, 0.4)
(      1.71, 0.133) (      1.5, 0.267) (      1.33, 0.4) (      1.2, 0.533)
(      1.5, 0.167) (      1.33, 0.333) (      1.2, 0.5) (      1.09, 0.667)
c[ix][iy][3]
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)
(      2, 0.1) (      1.71, 0.2) (      1.5, 0.3) (      1.33, 0.4)
(      1.71, 0.15) (      1.5, 0.3) (      1.33, 0.45) (      1.2, 0.6)
(      1.5, 0.2) (      1.33, 0.4) (      1.2, 0.6) (      1.09, 0.8)
(      1.33, 0.25) (      1.2, 0.5) (      1.09, 0.75) (      1, 1)

** Output **
< Forward Transform >
ierr =      0
c[ix][iy][1]
(      1.74, 0.25) (      0.102, -0.16) (      0.137, -0.05) (      0.202, 0.06)
(      0.108, -0.189) (      0.0379, -0.0469) (      0.0406, -0.0125) (      0.0525, 0.016)
(      0.125, -0.0784) (      0.034, -0.0168) (      0.0261, 0.00288) (      0.0254, 0.0209)
(      0.152, -0.00492) (      0.0366, 0.00116) (      0.0207, 0.0138) (      0.012, 0.028)
(      0.223, 0.106) (      0.0462, 0.0236) (      0.0177, 0.0292) (-0.00166, 0.0406)
c[ix][iy][2]
(      0.106, -0.127) (      0.0407, -0.0223) (      0.0315, 0.00295) (      0.0297, 0.0255)
(      0.0419, -0.0317) (-0.00167, -0.00877) (      0.0025, -0.00799) (      0.00901, -0.00976)
(      0.0317, -0.00698) (      0.00134, -0.00743) (      0.00423, -0.00524) (      0.00924, -0.00424)
(      0.0297, 0.0084) (      0.00473, -0.00711) (      0.00655, -0.00336) (      0.0108, 0.0001)
(      0.0318, 0.0285) (      0.0112, -0.00921) (      0.0118, -0.00179) (      0.016, 0.00627)
c[ix][iy][3]
(      0.178, 0.00231) (      0.0403, 0.014) (      0.017, 0.022) (      0.00125, 0.0329)
(      0.0484, 0.00885) (      0.00516, -0.0104) (      0.00849, -0.00569) (      0.0153, -0.0016)
(      0.0244, 0.0163) (      0.00692, -0.0061) (      0.0076, -0.00159) (      0.0107, 0.00322)
(      0.0129, 0.0239) (      0.00961, -0.0029) (      0.00799, 0.00185) (      0.00849, 0.0078)
(      0.00117, 0.036) (      0.0163, 0.000768) (      0.0106, 0.00714) (      0.00733, 0.0161)
< Backward Transform >
ierr =      0
c[ix][iy][1]
(      4, 0.0167) (      3, 0.0333) (      2.4, 0.05) (      2, 0.0667)
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)
(      2, 0.0667) (      1.71, 0.133) (      1.5, 0.2) (      1.33, 0.267)
(      1.71, 0.0833) (      1.5, 0.167) (      1.33, 0.25) (      1.2, 0.333)
c[ix][iy][2]
(      3, 0.0333) (      2.4, 0.0667) (      2, 0.1) (      1.71, 0.133)
(      2.4, 0.0667) (      2, 0.133) (      1.71, 0.2) (      1.5, 0.267)
(      2, 0.1) (      1.71, 0.2) (      1.5, 0.3) (      1.33, 0.4)
(      1.71, 0.133) (      1.5, 0.267) (      1.33, 0.4) (      1.2, 0.533)
(      1.5, 0.167) (      1.33, 0.333) (      1.2, 0.5) (      1.09, 0.667)
c[ix][iy][3]
(      2.4, 0.05) (      2, 0.1) (      1.71, 0.15) (      1.5, 0.2)

```

```
(      2,      0.1) (  1.71,    0.2) (   1.5,    0.3) (  1.33,    0.4)
(  1.71,    0.15) (   1.5,    0.3) (  1.33,    0.45) (   1.2,    0.6)
(   1.5,    0.2) (  1.33,    0.4) (   1.2,    0.6) (  1.09,    0.8)
(  1.33,    0.25) (   1.2,    0.5) (   1.09,    0.75) (     1,     1)
```

## 6.10 THREE-DIMENSIONAL REAL FOURIER TRANSFORM

### 6.10.1 [DEPRECATED]ASL\_qfr3fb, ASL\_pfr3fb

#### Three-Dimensional Real Fourier Transform (Including Initialization)

##### (1) Function

###### Forward transform

ASL\_qfr3fb or ASL\_pfr3fb obtains a half period of the three-dimensional Fourier forward transform (arbitrary radix) for the three-dimensional real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$c_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} r_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$(j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationships.

$$c_{n_x-j_x, n_y-j_y, n_z-j_z}^* = c_{j_x, j_y, j_z}$$

$$c_{n_x-j_x, j_y, j_z}^* = c_{j_x, n_y-j_y, n_z-j_z}$$

$$c_{n_x-j_x, n_y-j_y, j_z}^* = c_{j_x, j_y, n_z-j_z}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

###### Backward transform

Given the half period  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) for  $n_x n_y n_z$  complex data  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) satisfying  $c_{n_x-j_x, n_y-j_y, n_z-j_z}^* = c_{j_x, j_y, j_z}$ ,  $c_{n_x-j_x, j_y, j_z}^* = c_{j_x, n_y-j_y, n_z-j_z}$ , and  $c_{n_x-j_x, n_y-j_y, j_z}^* = c_{j_x, j_y, n_z-j_z}$ , ASL\_qfr3fb or ASL\_pfr3fb obtains the three-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$r_{k_x, k_y, k_z} = \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \sum_{j_z=0}^{n_z-1} c_{j_x, j_y, j_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$= \sum_{j_z=0}^{n_z-1} \sum_{j_y=0}^{n_y-1} \{c_{0, j_y, j_z} + (-1)^{k_x} \hat{c}_{\frac{n_x}{2}, j_y, j_z}\} e^{2\pi\sqrt{-1}\left(\frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$+ 2 \sum_{j_z=0}^{n_z-1} \sum_{j_y=0}^{n_y-1} \sum_{j_x=1}^{\lceil \frac{n_x}{2} \rceil - 1} \Re\{c_{j_x, j_y, j_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}\}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  represents the real part of the complex number  $z$ . Also, when  $n_x$  is odd,  $\hat{c}_{\frac{n_x}{2}, j_y, j_z} = 0$ , and when  $n_x$  is even,  $\hat{c}_{\frac{n_x}{2}, j_y, j_z} = c_{\frac{n_x}{2}, j_y, j_z}$ .

##### (2) Usage

Double precision:

ierr = ASL\_qfr3fb (nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfr3fb (nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int} \text{ as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	r	$\left\{ \begin{array}{l} D* \\ R* \end{array} \right\}$	lx×ly×lz	Input	Input data $r_{k_x, k_y, k_z}$ (Forward transform), or $c_{j_x, j_y, j_z}$ (Backward transform) (See Note (b))
				Output	Output results $c_{j_x, j_y, j_z}$ (Forward transform), or $r_{k_x, k_y, k_z}$ (Backward transform) (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array r (See Note (b))
6	ly	I	1	Input	Second dimension of array r (See Note (b))
7	lz	I	1	Input	Third dimension of array r (See Note (b))
8	isw	I	1	Input	Processing switch (See Note (d)) isw= 0:Initialization only isw= 1:Forward transform isw=-1:Backward transform
9	ifax	I*	60	Output	Factorization results and number of factors (See Note (d))
10	trigs	$\left\{ \begin{array}{l} D* \\ R* \end{array} \right\}$	nx + 2 × (ny + nz)	Output	Trigonometric function table (See Note (d))
11	wk	$\left\{ \begin{array}{l} D* \\ R* \end{array} \right\}$	lx × ly × lz	Work	Work area
12	nt	I	1	Input	Number of tasks to be generated
13	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $nx \geq 2, ny \geq 2, nz \geq 2$
- (b) In the case where  $nx$  is an odd,  $nx+1 \leq lx, ny \leq ly, nz \leq lz$   
 or if  $nx$  is an even,  $nx+2 \leq lx, ny \leq ly, nz \leq lz$
- (c)  $lx$  should be even number.
- (d)  $isw \in \{0, 1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) When the number of data  $nx, ny$  or  $nz$  can be adjusted, the calculations can be performed more efficiently by setting a number for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc.). For example, rather than setting  $nx = 289 (=17^2)$ , it is usually more efficient to set  $nx = 300 (=2^2 \times 3 \times 5^2)$ ,  $nx = 320 (=2^6 \times 5)$ ,  $nx = 384 (=2^7 \times 3)$  or the like.
- (b) The real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1$ ) and elements of array  $r$  are associated as follows.

$$r_{k_x, k_y, k_z} \leftrightarrow r[k_x + lx * (k_y + ly * k_z)]$$

When computing the backward transform, if  $nx (=n_x)$  is odd, then  $r[nx + lx * (k_y + ly * k_z)] = 0$ , and when  $nx$  is even, then  $r[nx + lx * (k_y + ly * k_z)] = r[nx + 1 + lx * (k_y + ly * k_z)] = 0$ . Also, when entering the real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1$ ) into array  $r$ , the corresponding zeros mentioned above need not be specifically stored.

If we let the real and imaginary parts of the complex data  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1$ ) be  $\Re\{c_{j_x, j_y, j_z}\}$  and  $\Im\{c_{j_x, j_y, j_z}\}$ , respectively, the  $c_{j_x, j_y, j_z}$  and elements of array  $r$  are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned} \Re\{c_{j_x, j_y, j_z}\} &\leftrightarrow r[2 * j_x + lx * (j_y + ly * j_z)] \\ \Im\{c_{j_x, j_y, j_z}\} &\leftrightarrow r[2 * j_x + 1 + lx * (j_y + ly * j_z)] \end{aligned}$$

From the properties of a real Fourier transform,  $\Im\{c_{0,0,0}\} = 0$ , and when  $nx, ny$  and  $nz$  are all even,  $\Im\{c_{\frac{n_x}{2}, \frac{n_y}{2}, \frac{n_z}{2}}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array  $r$ , they are considered to be zero when processing is performed. Since the elements  $c_{j_x, j_y, j_z}$  ( $j_x = \lfloor \frac{n_x}{2} \rfloor + 1, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1$ ) can be obtained according to the following relationships from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$\begin{aligned} c_{n_x - j_x, n_y - j_y, n_z - j_z}^* &= c_{j_x, j_y, j_z} \\ c_{n_x - j_x, j_y, j_z}^* &= c_{j_x, n_y - j_y, n_z - j_z} \\ c_{n_x - j_x, n_y - j_y, j_z}^* &= c_{j_x, j_y, n_z - j_z} \end{aligned}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ . **The adjustable dimensions of array r should be set so that lx/2, ly, and lz are odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within array r. Usually, when nx, for example, is (a multiple of 4)+2, lx=nx+4 is set.**

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k_x, k_y, k_z}(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$  be represented by  $\hat{r}_{k_x, k_y, k_z}(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$ , then the following relationship holds.

$$\hat{r}_{k_x, k_y, k_z} = n_x n_y n_z r_{k_x, k_y, k_z} \\ (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) To repeatedly compute the transform for the same number of data ( $n_x, n_y, n_z$ ), you should call this function once, and then use the after-initialization transform 6.10.2  $\left\{ \begin{array}{l} \text{ASL\_qfr3bf} \\ \text{ASL\_pfr3bf} \end{array} \right\}$ , thereafter. This enables processing to be performed more efficiently since initialization (factorization or the creation of trigonometric tables) is performed only once. However, in this case, you must retain the contents of arrays ifax and trigs so they can be used as input to the function 6.10.2  $\left\{ \begin{array}{l} \text{ASL\_qfr3bf} \\ \text{ASL\_pfr3bf} \end{array} \right\}$ . To perform initialization only by setting isw=0, you need not set input data for array r.

- (e) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (f) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

(7) **Example**

See the example in Section 6.10.2 (7).



## 6.10.2 [DEPRECATED]ASL\_qfr3bf, ASL\_pfr3bf Three-Dimensional Real Fourier Transform (After Initialization)

### (1) Function

#### Forward transform

ASL\_qfr3bf or ASL\_pfr3bf obtains a half period of the three-dimensional Fourier forward transform (arbitrary radix) for the three-dimensional real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1$ ;  $k_y = 0, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ).

$$c_{j_x, j_y, j_z} = \sum_{k_x=0}^{n_x-1} \sum_{k_y=0}^{n_y-1} \sum_{k_z=0}^{n_z-1} r_{k_x, k_y, k_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$(j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ . The remaining half period is obtained from the following relationships.

$$c_{n_x-j_x, n_y-j_y, n_z-j_z}^* = c_{j_x, j_y, j_z}$$

$$c_{n_x-j_x, j_y, j_z}^* = c_{j_x, n_y-j_y, n_z-j_z}$$

$$c_{n_x-j_x, n_y-j_y, j_z}^* = c_{j_x, j_y, n_z-j_z}$$

Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .

#### Backward transform

Given the half period  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) for  $n_x n_y n_z$  complex data  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) satisfying  $c_{n_x-j_x, n_y-j_y, n_z-j_z}^* = c_{j_x, j_y, j_z}$ ,  $c_{n_x-j_x, j_y, j_z}^* = c_{j_x, n_y-j_y, n_z-j_z}$ , and  $c_{n_x-j_x, n_y-j_y, j_z}^* = c_{j_x, j_y, n_z-j_z}$ , ASL\_qfr3bf or ASL\_pfr3bf obtains the three-dimensional Fourier backward transform (arbitrary radix) defined as follows.

$$r_{k_x, k_y, k_z} = \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \sum_{j_z=0}^{n_z-1} c_{j_x, j_y, j_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$= \sum_{j_z=0}^{n_z-1} \sum_{j_y=0}^{n_y-1} \{c_{0, j_y, j_z} + (-1)^{k_x} \hat{c}_{\frac{n_x}{2}, j_y, j_z}\} e^{2\pi\sqrt{-1}\left(\frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}$$

$$+ 2 \sum_{j_z=0}^{n_z-1} \sum_{j_y=0}^{n_y-1} \sum_{j_x=1}^{\lfloor \frac{n_x}{2} \rfloor - 1} \Re\{c_{j_x, j_y, j_z} e^{2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)}\}$$

$$(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

Here,  $\lceil x \rceil$  represents the minimum integer greater than or equal to  $x$ , and  $\Re\{z\}$  represents the real part of the complex number  $z$ . Also, when  $n_x$  is odd,  $\hat{c}_{\frac{n_x}{2}, j_y, j_z} = 0$ , and when  $n_x$  is even,  $\hat{c}_{\frac{n_x}{2}, j_y, j_z} = c_{\frac{n_x}{2}, j_y, j_z}$ .

### (2) Usage

Double precision:

ierr = ASL\_qfr3bf (nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

Single precision:

ierr = ASL\_pfr3bf (nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Number of data values in the first dimension, $n_x$ (See Note (a))
2	ny	I	1	Input	Number of data values in the second dimension, $n_y$ (See Note (a))
3	nz	I	1	Input	Number of data values in the third dimension, $n_z$ (See Note (a))
4	r	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	lx×ly×lz	Input	Input data $r_{k_x, k_y, k_z}$ (Forward transform), or $c_{j_x, j_y, j_z}$ (Backward transform) (See Note (b))
				Output	Output results $c_{j_x, j_y, j_z}$ (Forward transform), or $r_{k_x, k_y, k_z}$ (Backward transform) (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array r (See Note (b))
6	ly	I	1	Input	Second dimension of array r (See Note (b))
7	lz	I	1	Input	Third dimension of array r (See Note (b))
8	isw	I	1	Input	Processing switch isw= 1: Forward transform isw=-1: Backward transform
9	ifax	I*	60	Input	Factorization results and number of factors (See Note (a))
10	trigs	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	nx + 2 × (ny + nz)	Input	Trigonometric function table (See Note (a))
11	wk	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	lx × ly × lz	Work	Work area
12	nt	I	1	Input	Number of tasks to be generated
13	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $n_x \geq 2, n_y \geq 2, n_z \geq 2$
- (b) In the case where  $n_x$  is an odd,  $n_x + 1 \leq l_x, n_y \leq l_y, n_z \leq l_z$   
 or if  $n_x$  is an even,  $n_x + 2 \leq l_x, n_y \leq l_y, n_z \leq l_z$
- (c)  $l_x$  should be even number.
- (d)  $isw \in \{1, -1\}$
- (e)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	

(6) **Notes**

- (a) This function can be used to repeatedly compute the transform for the same number of data ( $n_x, n_y, n_z$ ) after the including-initialization function 6.10.1  $\left\{ \begin{array}{l} \text{ASL\_qfr3fb} \\ \text{ASL\_pfr3fb} \end{array} \right\}$  has been used. In this case, you must retain the contents of arrays ifax and trigs so they can be used as input in this function.
- (b) The real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1$ ) and elements of array r are associated as follows.

$$r_{k_x, k_y, k_z} \leftrightarrow r[k_x + l_x * (k_y + l_y * k_z)]$$

When computing the backward transform, if  $n_x (=n_x)$  is odd, then  $r[n_x + l_x * (k_y + l_y * k_z)] = 0$ , and when  $n_x$  is even, then  $r[n_x + l_x * (k_y + l_y * k_z)] = r[n_x + 1 + l_x * (k_y + l_y * k_z)] = 0$ . Also, when entering the real data  $r_{k_x, k_y, k_z}$  ( $k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1$ ) into array r, the corresponding zeros mentioned above need not be specifically stored.

If we let the real and imaginary parts of the complex data  $c_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, \lfloor \frac{n_x}{2} \rfloor; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1$ ) be  $\Re\{c_{j_x, j_y, j_z}\}$  and  $\Im\{c_{j_x, j_y, j_z}\}$ , respectively, the  $c_{j_x, j_y, j_z}$  and elements of array r are associated as follows. Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

$$\begin{aligned} \Re\{c_{j_x, j_y, j_z}\} &\leftrightarrow r[2 * j_x + l_x * (j_y + l_y * j_z)] \\ \Im\{c_{j_x, j_y, j_z}\} &\leftrightarrow r[2 * j_x + 1 + l_x * (j_y + l_y * j_z)] \end{aligned}$$

From the properties of a real Fourier transform,  $\Im\{c_{0,0,0}\} = 0$ , and when  $n_x, n_y$  and  $n_z$  are all even,  $\Im\{c_{\frac{n_x}{2}, \frac{n_y}{2}, \frac{n_z}{2}}\} = 0$ . Therefore, even if nonzero values are set for the corresponding elements of array r, they are considered to be zero when processing is performed. Since the elements  $c_{j_x, j_y, j_z}$  ( $j_x = \lfloor \frac{n_x}{2} \rfloor + 1, \dots, n_x - 1; j_y = 0, \dots, n_y - 1; j_z = 0, \dots, n_z - 1$ ) can be obtained according to the following relationships from the symmetry of the real Fourier transform, they need not be assigned as input when computing the backward transform. Also, they are not output when computing the forward transform.

$$\begin{aligned} c_{n_x - j_x, n_y - j_y, n_z - j_z}^* &= c_{j_x, j_y, j_z} \\ c_{n_x - j_x, j_y, j_z}^* &= c_{j_x, n_y - j_y, n_z - j_z} \\ c_{n_x - j_x, n_y - j_y, j_z}^* &= c_{j_x, j_y, n_z - j_z} \end{aligned}$$

The adjustable dimensions of array `r` should be set so that `lx/2`, `ly`, and `lz` are odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within array `r`. Usually, when `nx`, for example, is (a multiple of 4)+2, `lx=nx+4` is set.

- (c) When this function is used to compute the backward transform immediately following the forward transform, the values of the data obtained will be the original data multiplied by the number of data. For example, if we let the data obtained by computing the backward transform immediately following the forward transform for the real data  $r_{k_x, k_y, k_z}(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$  be represented by  $\hat{r}_{k_x, k_y, k_z}(k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$ , then the following relationship holds.

$$\hat{r}_{k_x, k_y, k_z} = n_x n_y n_z r_{k_x, k_y, k_z} \\ (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

Therefore, normalization must be performed for the result of either the forward transform or the backward transform. Note that in some of the entries in the Reference Bibliography, the definitions of the forward and backward transforms are reversed from those in this book, and in some of the entries a normalized result is defined.

- (d) Since a discrete Fourier transform is assumed to be a periodic function for which the data sequences before and after the transform are assumed to have the number of data ( $n_x$  or  $n_y$  or  $n_z$ ) as the period, the number of samples or sampling interval must be set with this taken into account when sampling to approximate the continuous Fourier transform. According to **the sampling theorem**, for a time function  $h(t)$  that is bandwidth limited by the frequency  $f_c$ , if the sampling interval is taken as  $T = \frac{1}{2f_c}$ , then  $h(t)$  can be reconstructed from knowledge of only a sequence of sample values  $\{h(iT)\}$  as follows.

$$h(t) = T \sum_{i=-\infty}^{\infty} h(iT) \frac{\sin 2\pi f_c(t - iT)}{\pi(t - iT)}$$

- (e) **DEPRECATED** This function will be removed in the future. Use **ASL Unified Interface**, the higher performance alternative implementation instead.

## (7) Example

- (a) Problem

Compute the three-dimensional real Fourier forward and backward transforms using

$$r_{k_x, k_y, k_z} = \frac{(k_x + 1)(k_y + 1)(k_z + 1)}{n_x n_y n_z} \\ (k_x = 0, \dots, n_x - 1; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

as input data.

- (b) Input data

Array `r`, `nx=6`, `ny=4`, `nz=3`, `lx=10`, `ly=5`, `lz=3`, `isw=1` (Forward transform), `isw=-1` (Backward transform) and `nt=2`.

- (c) Main program

```
/*      C Interface example for ASL_qfr3fb , ASL_qfr3bf */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int nx = 6;    int ny = 4;    int nz = 3;
    int lx = 10;  int ly = 5;    int lz = 3;
```

```

double *r;
int isw;
int ifax[60];
double *trigs;
double *wk;
int nt = 2;
int ierr;
int i,j,k;

printf( "    *** ASL_qfr3fb , ASL_qfr3bf ***\n" );
printf( "\n    ** Input **\n" );

r = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*lz) ));
if( r == NULL )
{
    printf( "no enough memory for array r\n" );
    return -1;
}

trigs = ( double * )malloc((size_t)( sizeof(double) * (nx+2*(ny+nz)) ));
if( trigs == NULL )
{
    printf( "no enough memory for array trigs\n" );
    return -1;
}

wk = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*lz) ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}

printf( "\tnx = %6d\n", nx );
printf( "\tny = %6d\n", ny );
printf( "\tnz = %6d\n", nz );
printf( "\tnt = %6d\n", nt );

for( k=1 ; k<=nz ; k++ )
{
    for( j=1 ; j<=ny ; j++ )
    {
        for( i=1 ; i<=nx ; i++ )
        {
            r[(i-1)+lx*(j-1)+lx*ly*(k-1)]=(double)(i*j*k)/(double)(nx*ny*nz) ;
        }
    }
}

printf( "\tr[ix][iy][1]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*1] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][3]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*2] );
    }
    printf( "\n" );
}

isw = 1;
ierr = ASL_qfr3fb(nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

for( i=0 ; i<lx*ly*lz ; i++)
{
    r[i] /= (double)(nx*ny*nz);
}

printf( "\n    ** Output **\n" );

```

```

printf( "\t< Forward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tr[ix][iy][1]\n" );
for( i=0 ; i<nx+2 ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j      ] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][2]\n" );
for( i=0 ; i<nx+2 ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*1] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][3]\n" );
for( i=0 ; i<nx+2 ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*2] );
    }
    printf( "\n" );
}

isw = -1;
ierr = ASL_qfr3bf(nx, ny, nz, r, lx, ly, lz, isw, ifax, trigs, wk, nt);

printf( "\t< Backward Transform >\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tr[ix][iy][1]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j      ] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][2]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*1] );
    }
    printf( "\n" );
}

printf( "\tr[ix][iy][3]\n" );
for( i=0 ; i<nx ; i++ )
{
    for( j=0 ; j<ny ; j++ )
    {
        printf( "\t%8.3g", r[i+lx*j+lx*ly*2] );
    }
    printf( "\n" );
}

free( r );
free( trigs );
free( wk );

return 0;
}

```

(d) Output results

```

*** ASL_qfr3fb , ASL_qfr3bf ***
** Input **
nx =      6
ny =      4
nz =      3
nt =      2
r[ix][iy][1]
0.0139      0.0278      0.0417      0.0556

```

```

    0.0278    0.0556    0.0833    0.111
    0.0417    0.0833    0.125    0.167
    0.0556    0.111    0.167    0.222
    0.0694    0.139    0.208    0.278
    0.0833    0.167    0.25    0.333
r[ix][iy][2]
    0.0278    0.0556    0.0833    0.111
    0.0556    0.111    0.167    0.222
    0.0833    0.167    0.25    0.333
    0.111    0.222    0.333    0.444
    0.139    0.278    0.417    0.556
    0.167    0.333    0.5    0.667
r[ix][iy][3]
    0.0417    0.0833    0.125    0.167
    0.0833    0.167    0.25    0.333
    0.125    0.25    0.375    0.5
    0.167    0.333    0.5    0.667
    0.208    0.417    0.625    0.833
    0.25    0.5    0.75    1

** Output **
< Forward Transform >
ierr = 0
r[ix][iy][1]
    0.243    -0.0486    -0.0486    -0.0486
    0    0.0486    0    -0.0486
    -0.0347    -0.00508    0.00694    0.019
    0.0601    -0.019    -0.012    -0.00508
    -0.0347    0.00294    0.00694    0.011
    0.02    -0.011    -0.00401    0.00294
    -0.0347    0.00694    0.00694    0.00694
    0    -0.00694    0    0.00694
r[ix][iy][2]
    -0.0608    0.00514    0.0122    0.0192
    0.0351    -0.0192    -0.00702    0.00514
    4.63e-18    0.00401    -1.54e-18    -0.00401
    -0.02    0.00401    0.00401    0.00401
    0.00579    0.000847    -0.00116    -0.00316
    -0.01    0.00316    0.002    0.000847
    0.00868    -0.000734    -0.00174    -0.00274
    -0.00501    0.00274    0.001    -0.000734
r[ix][iy][3]
    -0.0608    0.0192    0.0122    0.00514
    -0.0351    -0.00514    0.00702    0.0192
    0.0174    -0.00147    -0.00347    -0.00548
    -0.01    0.00548    0.002    -0.00147
    0.0116    -0.00231    -0.00231    -0.00231
    0    0.00231    0    -0.00231
    0.00868    -0.00274    -0.00174    -0.000734
    0.00501    0.000734    -0.001    -0.00274
< Backward Transform >
ierr = 0
r[ix][iy][1]
    0.0139    0.0278    0.0417    0.0556
    0.0278    0.0556    0.0833    0.111
    0.0417    0.0833    0.125    0.167
    0.0556    0.111    0.167    0.222
    0.0694    0.139    0.208    0.278
    0.0833    0.167    0.25    0.333
r[ix][iy][2]
    0.0278    0.0556    0.0833    0.111
    0.0556    0.111    0.167    0.222
    0.0833    0.167    0.25    0.333
    0.111    0.222    0.333    0.444
    0.139    0.278    0.417    0.556
    0.167    0.333    0.5    0.667
r[ix][iy][3]
    0.0417    0.0833    0.125    0.167
    0.0833    0.167    0.25    0.333
    0.125    0.25    0.375    0.5
    0.167    0.333    0.5    0.667
    0.208    0.417    0.625    0.833
    0.25    0.5    0.75    1
    
```

## 6.11 CONVOLUTIONS

### 6.11.1 ASL\_qfcn2d, ASL\_pfcn2d

#### Two-Dimensional Convolutions

(1) **Function**

Assume that the two multiperiodic discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  of period  $(m_x, m_y)$  satisfying:

$$\begin{aligned} f(i_x, i_y) &= f(i_x + L_x m_x, i_y + L_y m_y), \\ g(j_x, j_y) &= g(j_x + L_x m_x, j_y + L_y m_y), \\ &(i_x, j_x = 0, \dots, m_x - 1; i_y, j_y = 0, \dots, m_y - 1) \end{aligned}$$

for arbitrary integers  $L_x$  and  $L_y$  take nonzero values within their basic periods only for  $(i_x, i_y) \in [0, n_x^{(f)} - 1] \times [0, n_y^{(f)} - 1]$  and  $(j_x, j_y) \in [0, n_x^{(g)} - 1] \times [0, n_y^{(g)} - 1]$ . Here,  $[0, a] \times [0, b]$  is the direct product region (region contained in the square for which the point  $(0, 0)$  and the point  $(a, b)$  are diagonal points) on the plane in which the plane coordinates  $(i, j)$  lie. At this time, ASL\_qfcn2d or ASL\_pfcn2d calculates the discrete convolution  $p(k_x, k_y)$  defined as follows:

$$\begin{aligned} p(k_x, k_y) &= \sum_{i_x=0}^{m_x-1} \sum_{i_y=0}^{m_y-1} f(i_x, i_y) g(k_x - i_x, k_y - i_y) \\ &= \sum_{j_x=0}^{m_x-1} \sum_{j_y=0}^{m_y-1} g(j_x, j_y) f(k_x - j_x, k_y - j_y) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1) \end{aligned}$$

Here,  $m_x = \min(n_x^{(f)} + n_x^{(g)} - 1, M_x)$  and  $m_y = \min(n_y^{(f)} + n_y^{(g)} - 1, M_y)$  and  $M_x$  and  $M_y$  are arbitrary integers satisfying  $M_x \geq \max(n_x^{(f)}, n_x^{(g)})$  and  $M_y \geq \max(n_y^{(f)}, n_y^{(g)})$ , respectively. The two-dimensional real Fourier transform of  $p(k_x, k_y)$  can also be obtained.

(2) **Usage**

Double precision:

ierr = ASL\_qfcn2d (nx1, ny1, nx2, ny2, r1, lx1, ly1, r2, lx2, ly2, mx, my, isw, iwk, wk, nt);

Single precision:

ierr = ASL\_pfcn2d (nx1, ny1, nx2, ny2, r1, lx1, ly1, r2, lx2, ly2, mx, my, isw, iwk, wk, nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx1	I	1	Input	Number of effective data in $i_x$ direction $n_x^{(f)}$ for discrete function $f(i_x, i_y)$
2	ny1	I	1	Input	Number of effective data in $i_y$ direction $n_y^{(f)}$ for discrete function $f(i_x, i_y)$



No.	Argument and Return Value	Type	Size	Input/Output	Contents
3	nx2	I	1	Input	Number of effective data in $j_x$ direction $n_x^{(g)}$ for discrete function $g(j_x, j_y)$
4	ny2	I	1	Input	Number of effective data in $j_y$ direction $n_y^{(g)}$ for discrete function $g(j_x, j_y)$
5	r1	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lx1×ly1	Input	Values of discrete function $f(i_x, i_y)$ (See Note (a))
				Output	When $isw \geq 1$ , result of two-dimensional real Fourier transform of discrete function $f(i_x, i_y)$ (period $(M_x, M_y)$ )
6	lx1	I	1	Input	Adjustable dimension of array r1
7	ly1	I	1	Input	Second dimension of array r1
8	r2	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	lx2×ly2	Input	Values of discrete function $g(j_x, j_y)$ (See Note (a))
				Output	Value of discrete function $p(k_x, k_y)$ or its two-dimensional real Fourier transform (See Note (b))
9	lx2	I	1	Input	Adjustable dimension of array r2
10	ly2	I	1	Input	Second dimension of array r2
11	mx	I	1	Input	Parameter $M_x$ corresponding to the period $(m_x, m_y)$ of discrete functions $f(i_x, i_y)$ , $g(j_x, j_y)$ , and $p(k_x, k_y)$ (See Note (c))
12	my	I	1	Input	Parameter $M_y$ corresponding to the period $(m_x, m_y)$ of discrete functions $f(i_x, i_y)$ , $g(j_x, j_y)$ , and $p(k_x, k_y)$ (See Note (c))
13	isw	I	1	Input	Processing switch (See Note (d)) isw= 0: Calculate the convolution according to the definition. isw= 1: Calculate the convolution according to the FFT method. isw= 2: Calculate the real Fourier transform of the convolution.
14	iwk	I*	See Contents	Work	Work area <b>Size:</b> 0 (When $isw = 0$ ) 40 (When $isw \geq 1$ )

No.	Argument and Return Value	Type	Size	Input/Output	Contents
15	wk	$\begin{Bmatrix} D_* \\ R_* \end{Bmatrix}$	See Contents	Work	Work area <b>Size:</b> $nx2 \times ny2$ ((When isw= 0 and nx2 is odd) $(nx2 + 1) \times ny2$ (When isw= 0 and nx2 is even) $mx + 2 \times my + \max(lx1 \times ly1, lx2 \times ly2)$ (When isw $\geq 1$ )
16	nt	I	1	Input	Number of tasks to be generated
17	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, 1, 2\}$
- (b)  $nx1 > 1$  and  $ny1 > 1$
- (c)  $nx2 > 1$  and  $ny2 > 1$
- (d)  $mx \geq \text{MAX}(nx1, nx2)$  and  $my \geq \text{MAX}(ny1, ny2)$
- (e)  $lx1 \geq nx1$  and  $ly1 \geq ny1$  (When isw=0)  
 $lx1 \geq mx + 1$  and  $ly1 \geq my$  (When isw  $\geq 1$  and mx is odd)  
 $lx1 \geq mx + 2$  and  $ly1 \geq my$  (When isw  $\geq 1$  and mx is even)
- (f)  $lx2 \geq mx$  and  $ly2 \geq my$  (When isw=0)  
 $lx2 \geq mx + 1$  and  $ly2 \geq my$  (When isw  $\geq 1$  and mx is odd)  
 $lx2 \geq mx + 2$  and  $ly2 \geq my$  (When isw  $\geq 1$  and mx is even)
- (g)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$mx < nx1 + nx2 - 1$ or $my < ny1 + ny2 - 1$	Overlapping occurred during the convolution calculation.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	
3050	Restriction (f) was not satisfied.	
3060	Restriction (g) was not satisfied.	

(6) Notes

- (a) The values of the discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  and the elements of arrays r1 and r2 are associated as follows.

$$\begin{aligned} f(i_x, i_y) &\leftrightarrow r1[i_x + lx1 * i_y] \\ g(j_x, j_y) &\leftrightarrow r2[j_x + lx2 * j_y] \end{aligned}$$

Here,  $i_x = 0, \dots, n_x^{(f)} - 1$ ;  $i_y = 0, \dots, n_y^{(f)} - 1$  and  $j_x = 0, \dots, n_x^{(g)} - 1$ ;  $j_y = 0, \dots, n_y^{(g)} - 1$ , and no values need be entered in other elements. **The adjustable dimensions of arrays r1 and r2 should be set so that  $lx1/2$ ,  $ly1$ ,  $lx2/2$ , and  $ly2$  are odd numbers to avoid bank conflict of main memory. Usually, when  $mx$ , for example, is a multiple of 4,  $lx1=mx+3$  is set.**

- (b) The values of the discrete convolution  $p(k_x, k_y)$  and the elements of array r2 are associated as follows.

$$p(k_x, k_y) \leftrightarrow r2[k_x + lx2 * k_y]$$

Here,  $k_x = 0, \dots, M_x - 1$ ;  $k_y = 0, \dots, M_y - 1$ . When  $isw=2$  is set to obtain the two-dimensional real Fourier transform  $P(j_x, j_y)$  of the discrete convolution  $p(k_x, k_y)$ , which is defined as follows ( $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ ):

$$\begin{aligned} P(j_x, j_y) &= \frac{1}{M_x M_y} \sum_{k_x=0}^{M_x-1} \sum_{k_y=0}^{M_y-1} p(k_x, k_y) e^{-2\pi\sqrt{-1}(\frac{j_x k_x}{M_x} + \frac{j_y k_y}{M_y})} \\ &\quad (j_x = 0, \dots, \lfloor \frac{M_x}{2} \rfloor; j_y = 0, \dots, \lfloor \frac{M_y}{2} \rfloor) \end{aligned}$$

the following associations are made:

$$\begin{aligned} \Re\{P(j_x, j_y)\} &\leftrightarrow r2[2 * j_x + lx2 * j_y] \\ \Im\{P(j_x, j_y)\} &\leftrightarrow r2[2 * j_x + 1 + lx2 * j_y] \end{aligned}$$

In this case, note that the Fourier transform that is obtained is normalized. The remaining half period of the Fourier transform can be obtained from the symmetry of the real Fourier transform as follows:

$$\begin{aligned} P(M_x - j_x, M_y - j_y)^* &= P(j_x, j_y) \\ P(M_x - j_x, j_y)^* &= P(j_x, M_y - j_y) \end{aligned}$$

(Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .)

- (c) If  $mx \geq nx1 + nx2 - 1$  and  $my \geq ny1 + ny2 - 1$  are set, the convolution can be calculated without causing an overlap with the convolution of the next period. When  $mx > nx1 + nx2 - 1$  or  $my > ny1 + ny2 - 1$ , the following correspondences are made:

$$p(k_x, k_y) \leftrightarrow r2[k_x + lx2 * k_y]$$

and values that match 0.0 within the error range are stored in elements corresponding to  $k_x = nx1 + nx2 - 1, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y = ny1 + ny2 - 1, \dots, my - 1$ . When  $isw=0$ ,  $mx = nx1 + nx2 - 1$  and  $my = ny1 + ny2 - 1$  should be set. When  $isw \geq 1$ , the calculations can be performed more efficiently by setting a value for  $mx$  or  $my$  for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, if  $nx1=nx2=145$ , then when  $isw=0$ ,  $mx = 289(=17^2)$  should be set. However, when  $isw \geq 1$ , it is usually more efficient to set  $mx = 300(=2^2 \times 3 \times 5^2)$ ,  $mx = 320(=2^6 \times 5)$ ,  $mx = 384(=2^7 \times 3)$  or the like.

- (d) **Usually, the calculations can be performed more efficiently by setting  $isw=1$  to calculate the FFT convolution.** However, to conserve work area or if there is a restriction on the method of selecting the parameter  $mx$  or  $my$ , the calculations should be performed by setting  $isw=0$ .

- (e) To calculate the convolution of discrete functions for which the starting position of the nonzero portions are separated from the origin, first perform the calculations by shifting the functions so that the starting positions are at the origin, and then shift the calculation results again to obtain the final results more efficiently. For example, when the nonzero portions of the discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  are the intervals  $[i_0, i_0 + n_x^{(f)} - 1]$  and  $[j_0, j_0 + n_x^{(g)} - 1]$  for  $i_x$  and  $j_x$ , respectively, let  $\hat{f}(i_x, i_y)$  and  $\hat{g}(j_x, j_y)$  be defined as follows:

$$\hat{f}(i_x, i_y) = f(i_x - i_0, i_y), \quad \hat{g}(j_x, j_y) = g(j_x - j_0, j_y)$$

and apply this function to  $\hat{f}(i_x, i_y)$  and  $\hat{g}(j_x, j_y)$ . Let  $\hat{p}(k_x, k_y)$  represent the result that was obtained, and the convolution  $p(k_x, k_y)$  of the original functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  is given as follows:

$$p(k_x, k_y) = \hat{p}(k_x + (i_0 + j_0), k_y)$$

That is, the desired results are obtained if you shift  $f(i_x, i_y)$  and  $g(j_x, j_y)$  in the negative directions of  $i_x$  and  $j_x$  by  $i_0$  and  $j_0$ , respectively, before calculating the discrete convolution, and then shift the calculated value of the convolution after applying this function by  $i_0 + j_0$  in the positive direction of  $k_x$ . Similarly, for  $i_y, j_y$ , and  $k_y$ .

- (f) The sampling interval squared multiplied by the discrete convolution calculated by this function is the square approximation (or approximation by using the trapezoidal formula) of the continuous convolution integral of a bandwidth-limited function. Therefore, to raise the approximation precision, you must take a smaller sampling interval and a larger number of sample data. To associate these results with a continuous convolution, it is easiest to let  $p(n_x^{(f)} + n_x^{(g)} - 1, k_y) = 0$  and  $p(k_x, n_y^{(f)} + n_y^{(g)} - 1) = 0$  and consider  $(n_x^{(f)} + n_x^{(g)})(n_y^{(f)} + n_y^{(g)})$  data of  $p(k_x, k_y)$  ( $k_x = 0, 1, \dots, n_x^{(f)} + n_x^{(g)} - 1$ ;  $k_y = 0, 1, \dots, n_y^{(f)} + n_y^{(g)} - 1$ ). In this case, the coordinate  $(0, 0)$  element usually is associated with  $p(0, 0)$ . However,

When  $isw=0$ ,

then  $lx1 = nx1, ly1 = ny1, lx2 = mx, ly2 = my$ , and

$nwk = nx2 \times ny2$  (when  $nx2$  is odd) or

$nwk = (nx2 + 1) \times ny2$  (when  $nx2$  is even)

When  $isw \geq 1$ ,

then  $lx1=lx2=mx+1$  (when  $mx$  is odd) or

$lx1=lx2=mx+2$  (when  $mx$  is even),

$ly1=ly2=my$ , and  $nwk = mx + (lx1 + 2) \times my$ .

## (7) Example

- (a) Problem

Use the sampling interval  $\Delta$  to discretize the two finite waveforms defined by the following equations and calculate the discrete convolution.

$$f(x, y) = \begin{cases} x & ((x, y) \in [0, x_f] \times [0, y_f]) \\ 0 & \text{(Otherwise)} \end{cases}$$

$$g(x, y) = \begin{cases} x_g - x & ((x, y) \in [0, x_g] \times [0, y_g]) \\ 0 & \text{(Otherwise)} \end{cases}$$

- (b) Input data

Sampling data

$r1[i_x + lx1 * i_y] = f(i_x \Delta, i_y \Delta)$  ( $i_x = 0, 1, \dots, nx1 - 1$ ;  $i_y = 0, 1, \dots, ny1 - 1$ ),

$r2[j_x + lx2 * j_y] = g(j_x \Delta, j_y \Delta)$  ( $j_x = 0, 1, \dots, nx2 - 1$ ;  $j_y = 0, 1, \dots, ny2 - 1$ ).

Here,  $\Delta = 0.5$ .

$nx1, ny1, nx2, ny2, mx, my$  and  $isw$ .

(c) Main program

```

/*      C interface example for ASL_qfcn2d */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int nx1;
    int ny1;
    int nx2;
    int ny2;
    double *r1;
    int m0=8;
    int lx1;
    int ly1;
    double *r2;
    int lx2;
    int ly2;
    int mx;
    int my;
    int isw;
    int *iwk;
    int niwk=40;
    double *wk;
    int nwk;
    int ierr;
    int i,j;
    int nt = 2;
    double t;
    double dt=0.5;
    double xf=2.0,yf=2.0;
    double xg=2.0,yg=2.0;

    printf( "      *** ASL_qfcn2d ***\n" );
    printf( "\n      ** Input **\n\n" );

    isw=1;
    nx1=(int) xf/dt;
    ny1=(int) yf/dt;
    nx2=(int) xg/dt;
    ny2=(int) yg/dt;
    mx=my=m0;
    lx1=lx2=m0+2;
    ly1=ly2=m0;
    nwk=mx+2*my+lx2*my;

    r1 = ( double * )malloc((size_t)( sizeof(double) * (lx1*ly1) ));
    if( r1 == NULL )
    {
        printf( "no enough memory for array r1\n" );
        return -1;
    }
    r2 = ( double * )malloc((size_t)( sizeof(double) * (lx2*ly2) ));
    if( r2 == NULL )
    {
        printf( "no enough memory for array r2\n" );
        return -1;
    }
    wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }
    iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
    if( iwk == NULL )
    {
        printf( "no enough memory for array iwk\n" );
        return -1;
    }
    }

    printf( "\t isw = %6d\n\t (nx1, ny1) = (%3d,%3d)\n",
            isw, nx1, ny1);
    printf( "\t (nx2, ny2) = (%3d,%3d)\n",
            nx2, ny2);
    printf( "\t (mx , my ) = (%3d,%3d)\n\n",
            mx, my);

    for( j=0 ; j<ny1 ; j++ )
        for( i=0 ; i<nx1 ; i++ )
        {
            t=i*dt;
            r1[i+lx1*j]=t;
        }
}

```

```

for( j=0 ; j<ny2 ; j++ )
  for( i=0 ; i<nx2 ; i++ )
  {
    t=i*dt;
    r2[i+lx2*j]=xg-t;
  }
printf( "\tData r1[i+3d*j]\n", lx1 );
printf( "\ti/j      0      1      2      3\n" );
printf( "\t-----\n" );
for( i=0 ; i<nx1 ; i++ )
{
  printf( "\t%3d", i );
  for( j=0 ; j<ny1 ; j++ )
    printf( "%8.3g", r1[i+lx1*j] );
  printf( "\n" );
}
printf( "\n" );
printf( "\tData r2[i+3d*j]\n", lx2 );
printf( "\ti/j      0      1      2      3\n" );
printf( "\t-----\n" );
for( i=0 ; i<nx2 ; i++ )
{
  printf( "\t%3d", i );
  for( j=0 ; j<ny2 ; j++ )
    printf( "%8.3g", r2[i+lx2*j] );
  printf( "\n" );
}

ierr = ASL_qfcn2d(nx1, ny1, nx2, ny2, r1, lx1, ly1,
  r2, lx2, ly2, mx, my, isw, iwk, wk, nt);

printf( "\n    ** Output **\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tConvolution r2[i+3d*j]\n", lx2 );
printf( "\ti/j      0      1      2      3      4" );
printf( "\n      5      6      7\n" );
printf( "\t-----" );
printf( "-----\n" );
for( i=0 ; i<mx ; i++ )
{
  printf( "\t%2d", i );
  for( j=0 ; j<my ; j++ )
    printf( "%7.2lf", r2[i+lx2*j] );
  printf( "\n" );
}
free( iwk );
free( wk );
free( r2 );
free( r1 );

return 0;
}

```

(d) Output results

```

*** ASL_qfcn2d ***

** Input **

isw =      1
(nx1, ny1) = ( 4, 4)
(nx2, ny2) = ( 4, 4)
(mx , my ) = ( 8, 8)

Data r1[i+ 10*j]
i/j      0      1      2      3
-----
0      0      0      0      0
1     0.5    0.5    0.5    0.5
2      1      1      1      1
3     1.5    1.5    1.5    1.5

Data r2[i+ 10*j]
i/j      0      1      2      3
-----
0      2      2      2      2
1     1.5    1.5    1.5    1.5
2      1      1      1      1
3     0.5    0.5    0.5    0.5

** Output **

ierr =      0
Convolution r2[i+ 10*j]

```

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	0.00	-0.00	-0.00
1	1.00	2.00	3.00	4.00	3.00	2.00	1.00	-0.00
2	2.75	5.50	8.25	11.00	8.25	5.50	2.75	-0.00
3	5.00	10.00	15.00	20.00	15.00	10.00	5.00	-0.00
4	3.50	7.00	10.50	14.00	10.50	7.00	3.50	-0.00
5	2.00	4.00	6.00	8.00	6.00	4.00	2.00	-0.00
6	0.75	1.50	2.25	3.00	2.25	1.50	0.75	-0.00
7	-0.00	0.00	0.00	0.00	-0.00	-0.00	-0.00	-0.00

### 6.11.2 ASL\_qfcn3d, ASL\_pfcn3d Three-Dimensional Convolutions

(1) **Function**

Assume that the two multiperiodic discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  of period  $(m_x, m_y, m_z)$  satisfying:

$$\begin{aligned} f(i_x, i_y, i_z) &= f(i_x + L_x m_x, i_y + L_y m_y, i_z + L_z m_z), \\ g(j_x, j_y, j_z) &= g(j_x + L_x m_x, j_y + L_y m_y, j_z + L_z m_z), \\ &(i_x, j_x = 0, \dots, m_x - 1; i_y, j_y = 0, \dots, m_y - 1; i_z, j_z = 0, \dots, m_z - 1) \end{aligned}$$

for arbitrary integers  $L_x, L_y,$  and  $L_z$  take nonzero values within their basic periods only for  $(i_x, i_y, i_z) \in [0, n_x^{(f)} - 1] \times [0, n_y^{(f)} - 1] \times [0, n_z^{(f)} - 1]$  and  $(j_x, j_y, j_z) \in [0, n_x^{(g)} - 1] \times [0, n_y^{(g)} - 1] \times [0, n_z^{(g)} - 1]$ . Here,  $[0, a] \times [0, b] \times [0, c]$  is the direct product region (region contained in the cube for which the point  $(0, 0, 0)$  and the point  $(a, b, c)$  are diagonal points) in the space in which the space coordinates  $(i, j, k)$  lie. At this time, ASL\_qfcn3d or ASL\_pfcn3d calculates the discrete convolution  $p(k_x, k_y, k_z)$  defined as follows:

$$\begin{aligned} p(k_x, k_y, k_z) &= \sum_{i_x=0}^{m_x-1} \sum_{i_y=0}^{m_y-1} \sum_{i_z=0}^{m_z-1} f(i_x, i_y, i_z) g(k_x - i_x, k_y - i_y, k_z - i_z) \\ &= \sum_{j_x=0}^{m_x-1} \sum_{j_y=0}^{m_y-1} \sum_{j_z=0}^{m_z-1} g(j_x, j_y, j_z) f(k_x - j_x, k_y - j_y, k_z - j_z) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1; k_z = 0, \dots, m_z - 1) \end{aligned}$$

Here,  $m_x = \min(n_x^{(f)} + n_x^{(g)} - 1, M_x)$ ,  $m_y = \min(n_y^{(f)} + n_y^{(g)} - 1, M_y)$ , and  $m_z = \min(n_z^{(f)} + n_z^{(g)} - 1, M_z)$  and  $M_x, M_y,$  and  $M_z$  are arbitrary integers satisfying  $M_x \geq \max(n_x^{(f)}, n_x^{(g)})$ ,  $M_y \geq \max(n_y^{(f)}, n_y^{(g)})$ , and  $M_z \geq \max(n_z^{(f)}, n_z^{(g)})$ , respectively. The three-dimensional real Fourier transform of  $p(k_x, k_y, k_z)$  can also be obtained.

(2) **Usage**

Double precision:

```
ierr = ASL_qfcn3d (nx1, ny1, nz1, nx2, ny2, nz2, r1, lx1, ly1, lz1, r2, lx2, ly2, lz2, mx, my, mz,
                 isw, iwk, wk, nt);
```

Single precision:

```
ierr = ASL_pfcn3d (nx1, ny1, nz1, nx2, ny2, nz2, r1, lx1, ly1, lz1, r2, lx2, ly2, lz2, mx, my, mz,
                 isw, iwk, wk, nt);
```



(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx1	I	1	Input	Number of effective data in $i_x$ direction $n_x^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
2	ny1	I	1	Input	Number of effective data in $i_y$ direction $n_y^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
3	nz1	I	1	Input	Number of effective data in $i_z$ direction $n_z^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
4	nx2	I	1	Input	Number of effective data in $j_x$ direction $n_x^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
5	ny2	I	1	Input	Number of effective data in $j_y$ direction $n_y^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
6	nz2	I	1	Input	Number of effective data in $j_z$ direction $n_z^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
7	r1	$\begin{cases} D* \\ R* \end{cases}$	$lx1 \times ly1 \times lz1$	Input	Values of discrete function $f(i_x, i_y, i_z)$ (See Note (a))
				Output	When $isw \geq 1$ , result of three-dimensional real Fourier transform of discrete function $f(i_x, i_y, i_z)$ (period $(M_x, M_y, M_z)$ )
8	lx1	I	1	Input	Adjustable dimension of array r1
9	ly1	I	1	Input	Second dimension of array r1
10	lz1	I	1	Input	Third dimension of array r1
11	r2	$\begin{cases} D* \\ R* \end{cases}$	$lx2 \times ly2 \times lz2$	Input	Values of discrete function $g(j_x, j_y, j_z)$ (See Note (a))
				Output	Value of discrete function $p(k_x, k_y, k_z)$ or its three-dimensional real Fourier transform (See Note (b))
12	lx2	I	1	Input	Adjustable dimension of array r2
13	ly2	I	1	Input	Second dimension of array r2
14	lz2	I	1	Input	Third dimension of array r2

No.	Argument and Return Value	Type	Size	Input/Output	Contents
15	mx	I	1	Input	Parameter $M_x$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $p(k_x, k_y, k_z)$ (See Note (c))
16	my	I	1	Input	Parameter $M_y$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $p(k_x, k_y, k_z)$ (See Note (c))
17	mz	I	1	Input	Parameter $M_z$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $p(k_x, k_y, k_z)$ (See Note (c))
18	isw	I	1	Input	Processing switch (See Note (d)) isw= 0: Calculate the convolution according to the definition. isw= 1: Calculate the convolution according to the FFT method. isw= 2: Calculate the real Fourier transform of the convolution.
19	iwk	I*	See Contents	Work	Work area <b>Size:</b> 0 (When isw= 0) 60 (When isw $\geq$ 1)
20	wk	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	See Contents	Work	Work area <b>Size:</b> $(nx2 + 1) \times (ny2 + 1) \times nz2$ (When isw= 0, nx2 is even and ny2 is even) $nx2 \times (ny2 + 1) \times nz2$ (When isw= 0, nx2 is odd and ny2 is even) $(nx2 + 1) \times ny2 \times nz2$ (When isw= 0, nx2 is even and ny2 is odd) $nx2 \times ny2 \times nz2$ (When isw= 0, nx2 is odd and ny2 is odd) $mx + 2 \times (my + mz) + \max(lx1 \times ly1 \times lz1, lx2 \times ly2 \times lz2)$ (When isw $\geq$ 1)
21	nt	I	1	Input	Number of tasks to be generated
22	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, 1, 2\}$
- (b)  $nx1 > 1$  and  $ny1 > 1$  and  $nz1 > 1$
- (c)  $nx2 > 1$  and  $ny2 > 1$  and  $nz2 > 1$
- (d)  $mx \geq \text{MAX}(nx1, nx2)$  and  $my \geq \text{MAX}(ny1, ny2)$  and  $mz \geq \text{MAX}(nz1, nz2)$
- (e)  $lx1 \geq nx1$  and  $ly1 \geq ny1$  and  $lz1 \geq nz1$  (when  $isw=0$ )  
 $lx1 \geq mx + 1$  and  $lx1$  is even and  $ly1 \geq my$  and  $lz1 \geq mz$  (when  $isw \geq 1$  and  $mx$  is odd)  
 $lx1 \geq mx + 2$  and  $lx1$  is even and  $ly1 \geq my$  and  $lz1 \geq mz$  (when  $isw \geq 1$  and  $mx$  is even)
- (f)  $lx2 \geq mx$  and  $ly2 \geq ny2$  and  $lz2 \geq nz2$  (when  $isw=0$ )  
 $lx2 \geq mx + 1$  and  $lx2$  is even and  $ly2 \geq my$  and  $lz2 \geq mz$  (when  $isw \geq 1$  and  $mx$  is odd)  
 $lx2 \geq mx + 2$  and  $lx2$  is even and  $ly2 \geq my$  and  $lz2 \geq mz$  (when  $isw \geq 1$  and  $mx$  is even)
- (g)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$mx < nx1 + nx2 - 1$ or $my < ny1 + ny2 - 1$ or $mz < nz1 + nz2 - 1$	Overlapping occurred during the convolution calculation.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	
3050	Restriction (f) was not satisfied.	
3060	Restriction (g) was not satisfied.	

(6) **Notes**

- (a) The values of the discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  and the elements of arrays r1 and r2 are associated as follows.

$$f(i_x, i_y, i_z) \leftrightarrow r1[i_x + lx1 * (i_y + ly1 * i_z)]$$

$$g(j_x, j_y, j_z) \leftrightarrow r2[j_x + lx2 * (j_y + ly2 * j_z)]$$

Here,  $i_x = 0, \dots, n_x^{(f)} - 1$ ;  $i_y = 0, \dots, n_y^{(f)} - 1$ ;  $i_z = 0, \dots, n_z^{(f)} - 1$  and  $j_x = 0, \dots, n_x^{(g)} - 1$ ;  $j_y = 0, \dots, n_y^{(g)} - 1$ ;  $j_z = 0, \dots, n_z^{(g)} - 1$ , and no values need be entered in other elements. **The adjustable dimensions of arrays r1 and r2 should be set so that  $lx1/2$ ,  $ly1$ ,  $lz1$ ,  $lx2/2$ ,  $ly2$ , and  $lz2$  are odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within arrays r1 and r2. Usually, when  $mx$ , for example, is (a multiple of 4)+2,  $lx1=mx+4$  is set.**

- (b) The values of the discrete convolution  $p(k_x, k_y, k_z)$  and the elements of array r2 are associated as follows.

$$p(k_x, k_y, k_z) \leftrightarrow r2[k_x + lx2 * (k_y + ly2 * k_z)]$$

Here,  $k_x = 0, \dots, M_x - 1$ ;  $k_y = 0, \dots, M_y - 1$ ;  $k_z = 0, \dots, M_z - 1$ . When  $isw=2$  is set to obtain the three-dimensional real Fourier transform  $P(j_x, j_y, j_z)$  of the discrete convolution  $p(k_x, k_y, k_z)$ , which is defined as follows ( $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ ):

$$P(j_x, j_y, j_z) = \frac{1}{M_x M_y M_z} \sum_{k_x=0}^{M_x-1} \sum_{k_y=0}^{M_y-1} \sum_{k_z=0}^{M_z-1} p(k_x, k_y, k_z) e^{-2\pi\sqrt{-1}(\frac{j_x k_x}{M_x} + \frac{j_y k_y}{M_y} + \frac{j_z k_z}{M_z})}$$

$$(j_x = 0, \dots, \lfloor \frac{M_x}{2} \rfloor; j_y = 0, \dots, \lfloor \frac{M_y}{2} \rfloor; j_z = 0, \dots, \lfloor \frac{M_z}{2} \rfloor)$$

the following associations are made:

$$\Re\{P(j_x, j_y, j_z)\} \leftrightarrow r2[2 * j_x + lx2 * (j_y + ly2 * j_z)]$$

$$\Im\{P(j_x, j_y, j_z)\} \leftrightarrow r2[2 * j_x + 1 + lx2 * (j_y + ly2 * j_z)]$$

In this case, note that the Fourier transform that is obtained is normalized. The remaining half period of the Fourier transform can be obtained from the symmetry of the real Fourier transform as follows:

$$P(M_x - j_x, M_y - j_y, M_z - j_z)^* = P(j_x, j_y, j_z)$$

$$P(M_x - j_x, j_y, j_z)^* = P(j_x, M_y - j_y, M_z - j_z)$$

$$P(M_x - j_x, M_y - j_y, j_z)^* = P(j_x, j_y, M_z - j_z)$$

(Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .)

- (c) If  $mx \geq nx1 + nx2 - 1$  and  $my \geq ny1 + ny2 - 1$  and  $mz \geq nz1 + nz2 - 1$  are set, the convolution can be calculated without causing an overlap with the convolution of the next period. When  $mx > nx1 + nx2 - 1$  or  $my > ny1 + ny2 - 1$  or  $mz > nz1 + nz2 - 1$ , the following correspondences are made:

$$p(k_x, k_y, k_z) \leftrightarrow r2[k_x + lx2 * (k_y + ly2 * k_z)]$$

and values that match 0.0 within the error range are stored in elements corresponding to  $k_x = nx1 + nx2 - 1, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$ ;  $k_z = 0, \dots, mz - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y = ny1 + ny2 - 1, \dots, my - 1$ ;  $k_z = 0, \dots, mz - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$ ;  $k_z = nz1 + nz2 - 1, \dots, mz - 1$ . When  $isw=0$ ,  $mx = nx1 + nx2 - 1$ ,  $my = ny1 + ny2 - 1$ , and  $mz = nz1 + nz2 - 1$  should be set. When  $isw \geq 1$ , the calculations can be performed more efficiently by setting a value for  $mx$ ,  $my$  or  $mz$  for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, if  $nx1=nx2=145$ , then when  $isw=0$ ,  $mx = 289(=17^2)$  should be set. However, when  $isw \geq 1$ , it is usually more efficient to set  $mx = 300(=2^2 \times 3 \times 5^2)$ ,  $mx = 320(=2^6 \times 5)$ ,  $mx = 384(=2^7 \times 3)$  or the like.

- (d) **Usually, the calculations can be performed more efficiently by setting  $isw=1$  to calculate the FFT convolution.** However, to conserve work area or if there is a restriction on the method of selecting the parameter  $mx$ ,  $my$  or  $mz$ , the calculations should be performed by setting  $isw=0$ .
- (e) To calculate the convolution of discrete functions for which the starting position of the nonzero portions are separated from the origin, first perform the calculations by shifting the functions so that the starting positions are at the origin, and then shift the calculation results again to obtain the final results more efficiently. For example, when the nonzero portions of the discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  are the intervals  $[i_0, i_0 + n_x^{(f)} - 1]$  and  $[j_0, j_0 + n_x^{(g)} - 1]$  for  $i_x$  and  $j_x$ , respectively, let  $\hat{f}(i_x, i_y, i_z)$  and  $\hat{g}(j_x, j_y, j_z)$  be defined as follows:

$$\hat{f}(i_x, i_y, i_z) = f(i_x - i_0, i_y, i_z), \quad \hat{g}(j_x, j_y, j_z) = g(j_x - j_0, j_y, j_z)$$

and apply this function to  $\hat{f}(i_x, i_y, i_z)$  and  $\hat{g}(j_x, j_y, j_z)$ . Let  $\hat{p}(k_x, k_y, k_z)$  represent the result that was obtained, and the convolution  $p(k_x, k_y, k_z)$  of the original functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  is given as follows:

$$p(k_x, k_y, k_z) = \hat{p}(k_x + (i_0 + j_0), k_y, k_z)$$

That is, the desired results are obtained if you shift  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  in the negative directions of  $i_x$  and  $j_x$  by  $i_0$  and  $j_0$ , respectively, before calculating the discrete convolution, and then shift the calculated value of the convolution after applying this function by  $i_0 + j_0$  in the positive direction of  $k_x$ . Similarly, for  $i_y, j_y$ , and  $k_y$  and  $i_z, j_z$ , and  $k_z$ .

- (f) The sampling interval cubed multiplied by the discrete convolution calculated by this function is the square approximation (or approximation by using the trapezoidal formula) of the continuous convolution integral of a bandwidth-limited function. Therefore, to raise the approximation precision, you must take a smaller sampling interval and a larger number of sample data. To associate these results with a continuous convolution, it is easiest to let  $p(n_x^{(f)} + n_x^{(g)} - 1, k_y, k_z) = 0$ ,  $p(k_x, n_y^{(f)} + n_y^{(g)} - 1, k_z) = 0$ , and  $p(k_x, k_y, n_z^{(f)} + n_z^{(g)} - 1) = 0$  and consider  $(n_x^{(f)} + n_x^{(g)})(n_y^{(f)} + n_y^{(g)})(n_z^{(f)} + n_z^{(g)})$  data of  $p(k_x, k_y, k_z)$  ( $k_x = 0, 1, \dots, n_x^{(f)} + n_x^{(g)} - 1$ ;  $k_y = 0, 1, \dots, n_y^{(f)} + n_y^{(g)} - 1$ ;  $k_z = 0, 1, \dots, n_z^{(f)} + n_z^{(g)} - 1$ ). In this case, the coordinate  $(0, 0, 0)$  element usually is associated with  $p(0, 0, 0)$ . However,

When  $isw=0$ ,

then  $lx1 = nx1, ly1 = ny1, lz1 = nz1, lx2 = mx, ly2 = my, lz2 = mz$ , and  
 $nwk = (nx2 + 1) \times (ny2 + 1) \times nz2$  (when  $nx2$  is even and  $ny2$  is even) or  
 $nwk = nx2 \times (ny2 + 1) \times nz2$  (when  $nx2$  is odd and  $ny2$  is even) or  
 $nwk = (nx2 + 1) \times ny2 \times nz2$  (when  $nx2$  is even and  $ny2$  is odd) or  
 $nwk = nx2 \times ny2 \times nz2$  (when  $nx2$  is odd and  $ny2$  is odd)

When  $isw \geq 1$

$lx1=lx2=mx+1$  (when  $mx$  is odd) or

$lx1=lx2=mx+2$  (when  $mx$  is even),

$ly1=ly2=my, lz1=lz2=mz$ , and  $nwk = mx + 2 \times (my + mz) + lx1 \times my \times mz$ .

## (7) Example

- (a) Problem

Use the sampling interval  $\Delta$  to discretize the two finite waveforms defined by the following equations and calculate the discrete convolution.

$$f(x, y, z) = \begin{cases} x & ((x, y, z) \in [0, x_f] \times [0, y_f] \times [0, z_f]) \\ 0 & \text{(Otherwise)} \end{cases}$$

$$g(x, y, z) = \begin{cases} x_g - x & ((x, y, z) \in [0, x_g] \times [0, y_g] \times [0, z_g]) \\ 0 & \text{(Otherwise)} \end{cases}$$

- (b) Input data

Sampling data

$r1[i_x + lx1 * (i_y + ly1 * i_z)] = f(i_x \Delta, i_y \Delta, i_z \Delta)$  ( $i_x = 0, 1, \dots, nx1 - 1$ ;  $i_y = 0, 1, \dots, ny1 - 1$ ;  $i_z = 0, 1, \dots, nz1 - 1$ ),

$r2[j_x + lx2 * (j_y + ly2 * j_z)] = g(j_x \Delta, j_y \Delta, j_z \Delta)$  ( $j_x = 0, 1, \dots, nx2 - 1$ ;  $j_y = 0, 1, \dots, ny2 - 1$ ;  $j_z = 0, 1, \dots, nz2 - 1$ ).

Here,  $\Delta = 0.5$ .

$nx1, ny1, nz1, nx2, ny2, nz2, mx, my, mz$  and  $isw$ .

- (c) Main program

```
/*      C interface example for ASL_qfcn3d */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
```

```

{
    int nx1;
    int ny1;
    int nz1;
    int nx2;
    int ny2;
    int nz2;
    double *r1;
    int m0=8;
    int lx1;
    int ly1;
    int lz1;
    double *r2;
    int lx2;
    int ly2;
    int lz2;
    int mx;
    int my;
    int mz;
    int isw;
    int *iwk;
    int niwk=60;
    double *wk;
    int nwk;
    int ierr;
    int i,j,k;
    int nt = 2;
    double t;
    double dt=0.5;
    double xf=2.0,yf=2.0,zf=2.0;
    double xg=2.0,yg=2.0,zg=2.0;

    printf( "    *** ASL_qfcn3d ***\n" );
    printf( "\n    ** Input **\n\n" );

    isw=1;
    nx1=(int) xf/dt;
    ny1=(int) yf/dt;
    nz1=(int) zf/dt;
    nx2=(int) xg/dt;
    ny2=(int) yg/dt;
    nz2=(int) zg/dt;
    mx=my=mz=m0;
    lx1=lx2=(m0+2)/2*2;
    ly1=ly2=my;
    lz1=lz2=mz;
    nwk=mx+2*(my+mz)+lx2*my*mz;

    r1 = ( double * )malloc((size_t)( sizeof(double) * (lx1*ly1+lz1) ));
    if( r1 == NULL )
    {
        printf( "no enough memory for array r1\n" );
        return -1;
    }
    r2 = ( double * )malloc((size_t)( sizeof(double) * (lx2*ly2+lz2) ));
    if( r2 == NULL )
    {
        printf( "no enough memory for array r2\n" );
        return -1;
    }
    wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }
    iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
    if( iwk == NULL )
    {
        printf( "no enough memory for array iwk\n" );
        return -1;
    }
    }

    printf( "\t isw = %6d\n\t (nx1, ny1, nz1) = (%3d,%3d,%3d)\n",
        isw, nx1, ny1, nz1);
    printf( "\t (nx2, ny2, nz2) = (%3d,%3d,%3d)\n",
        nx2, ny2, nz2);
    printf( "\t (mx , my , mz ) = (%3d,%3d,%3d)\n\n",
        mx, my, mz);

    for( k=0 ; k<nz1 ; k++ )
        for( j=0 ; j<ny1 ; j++ )
            for( i=0 ; i<nx1 ; i++ )
                {
                    t=i*dt;
                    r1[i+lx1*(j+ly1*k)]=t;
                }
    for( k=0 ; k<nz2 ; k++ )
        for( j=0 ; j<ny2 ; j++ )
            for( i=0 ; i<nx2 ; i++ )
                {

```

```

        t=i*dt;
        r2[i+lx2*(j+ly2*k)]=xg-t;
    }
    for( k=0 ; k<nz1 ; k++ )
    {
        printf( "\tData r1[i+%3d*(j+%3d*%3d)]\n", lx1, ly1, k );
        printf( "\ti/j      0      1      2      3\n" );
        printf( "\t-----\n" );
        for( i=0 ; i<nx1 ; i++ )
        {
            printf( "\t%3d", i );
            for( j=0 ; j<ny1 ; j++ )
                printf( "%8.3g", r1[i+lx1*(j+ly1*k)] );
            printf( "\n" );
        }
        printf( "\n" );
    }
    for( k=0 ; k<nz2 ; k++ )
    {
        printf( "\tData r2[i+%3d*(j+%3d*%3d)]\n", lx2, ly2, k );
        printf( "\ti/j      0      1      2      3\n" );
        printf( "\t-----\n" );
        for( i=0 ; i<nx2 ; i++ )
        {
            printf( "\t%3d", i );
            for( j=0 ; j<ny2 ; j++ )
                printf( "%8.3g", r2[i+lx2*(j+ly2*k)] );
            printf( "\n" );
        }
        printf( "\n" );
    }
}

ierr = ASL_qfcn3d(nx1, ny1, nz1, nx2, ny2, nz2,
    r1, lx1, ly1, lz1, r2, lx2, ly2, lz2,
    mx, my, mz, isw, iwk, wk, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

for( k=0 ; k<mz ; k++ )
{
    printf( "\tConvolution r2[i+%3d*(j+%3d*%3d)]\n", lx2, ly2, k );
    printf( "\ti/j      0      1      2      3      4" );
    printf( "      5      6      7\n" );
    printf( "\t-----" );
    printf( "-----\n" );
    for( i=0 ; i<mx ; i++ )
    {
        printf( "\t%2d", i );
        for( j=0 ; j<my ; j++ )
            printf( "%7.2lf", r2[i+lx2*(j+ly2*k)] );
        printf( "\n" );
    }
    printf( "\n" );
}
free( iwk );
free( wk );
free( r2 );
free( r1 );

return 0;
}

```

(d) Output results

```

*** ASL_qfcn3d ***

** Input **

isw =      1
(nx1, ny1, nz1) = ( 4, 4, 4)
(nx2, ny2, nz2) = ( 4, 4, 4)
(mx , my , mz ) = ( 8, 8, 8)

Data r1[i+ 10*(j+ 8* 0)]
i/j      0      1      2      3
-----
0      0      0      0      0
1     0.5    0.5    0.5    0.5
2      1      1      1      1
3     1.5    1.5    1.5    1.5

Data r1[i+ 10*(j+ 8* 1)]
i/j      0      1      2      3
-----

```

0	0	0	0	0
1	0.5	0.5	0.5	0.5
2	1	1	1	1
3	1.5	1.5	1.5	1.5

Data r1[i+ 10\*(j+ 8\* 2)]

i/j	0	1	2	3
0	0	0	0	0
1	0.5	0.5	0.5	0.5
2	1	1	1	1
3	1.5	1.5	1.5	1.5

Data r1[i+ 10\*(j+ 8\* 3)]

i/j	0	1	2	3
0	0	0	0	0
1	0.5	0.5	0.5	0.5
2	1	1	1	1
3	1.5	1.5	1.5	1.5

Data r2[i+ 10\*(j+ 8\* 0)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

Data r2[i+ 10\*(j+ 8\* 1)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

Data r2[i+ 10\*(j+ 8\* 2)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

Data r2[i+ 10\*(j+ 8\* 3)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

\*\* Output \*\*

ierr = 0

Convolution r2[i+ 10\*(j+ 8\* 0)]

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	-0.00	-0.00	0.00
1	1.00	2.00	3.00	4.00	3.00	2.00	1.00	0.00
2	2.75	5.50	8.25	11.00	8.25	5.50	2.75	0.00
3	5.00	10.00	15.00	20.00	15.00	10.00	5.00	0.00
4	3.50	7.00	10.50	14.00	10.50	7.00	3.50	0.00
5	2.00	4.00	6.00	8.00	6.00	4.00	2.00	0.00
6	0.75	1.50	2.25	3.00	2.25	1.50	0.75	-0.00
7	0.00	0.00	0.00	0.00	-0.00	-0.00	-0.00	0.00

Convolution r2[i+ 10\*(j+ 8\* 1)]

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	0.00	-0.00	-0.00
1	2.00	4.00	6.00	8.00	6.00	4.00	2.00	-0.00
2	5.50	11.00	16.50	22.00	16.50	11.00	5.50	-0.00
3	10.00	20.00	30.00	40.00	30.00	20.00	10.00	-0.00
4	7.00	14.00	21.00	28.00	21.00	14.00	7.00	-0.00
5	4.00	8.00	12.00	16.00	12.00	8.00	4.00	-0.00
6	1.50	3.00	4.50	6.00	4.50	3.00	1.50	-0.00
7	0.00	-0.00	-0.00	0.00	-0.00	-0.00	-0.00	-0.00

Convolution r2[i+ 10\*(j+ 8\* 2)]

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	0.00	0.00	-0.00
1	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
2	8.25	16.50	24.75	33.00	24.75	16.50	8.25	-0.00
3	15.00	30.00	45.00	60.00	45.00	30.00	15.00	-0.00
4	10.50	21.00	31.50	42.00	31.50	21.00	10.50	-0.00
5	6.00	12.00	18.00	24.00	18.00	12.00	6.00	-0.00
6	2.25	4.50	6.75	9.00	6.75	4.50	2.25	-0.00
7	0.00	-0.00	0.00	-0.00	0.00	0.00	-0.00	-0.00



Convolution r2[i+ 10\*(j+ 8\* 3)]

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	0.00	-0.00	0.00
1	4.00	8.00	12.00	16.00	12.00	8.00	4.00	0.00
2	11.00	22.00	33.00	44.00	33.00	22.00	11.00	-0.00
3	20.00	40.00	60.00	80.00	60.00	40.00	20.00	-0.00
4	14.00	28.00	42.00	56.00	42.00	28.00	14.00	-0.00
5	8.00	16.00	24.00	32.00	24.00	16.00	8.00	-0.00
6	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
7	-0.00	-0.00	0.00	-0.00	-0.00	0.00	-0.00	0.00

Convolution r2[i+ 10\*(j+ 8\* 4)]

i/j	0	1	2	3	4	5	6	7
0	0.00	-0.00	0.00	0.00	0.00	0.00	-0.00	-0.00
1	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
2	8.25	16.50	24.75	33.00	24.75	16.50	8.25	-0.00
3	15.00	30.00	45.00	60.00	45.00	30.00	15.00	-0.00
4	10.50	21.00	31.50	42.00	31.50	21.00	10.50	0.00
5	6.00	12.00	18.00	24.00	18.00	12.00	6.00	0.00
6	2.25	4.50	6.75	9.00	6.75	4.50	2.25	-0.00
7	-0.00	-0.00	0.00	0.00	0.00	-0.00	-0.00	-0.00

Convolution r2[i+ 10\*(j+ 8\* 5)]

i/j	0	1	2	3	4	5	6	7
0	0.00	0.00	0.00	0.00	0.00	0.00	-0.00	0.00
1	2.00	4.00	6.00	8.00	6.00	4.00	2.00	0.00
2	5.50	11.00	16.50	22.00	16.50	11.00	5.50	-0.00
3	10.00	20.00	30.00	40.00	30.00	20.00	10.00	-0.00
4	7.00	14.00	21.00	28.00	21.00	14.00	7.00	-0.00
5	4.00	8.00	12.00	16.00	12.00	8.00	4.00	-0.00
6	1.50	3.00	4.50	6.00	4.50	3.00	1.50	-0.00
7	0.00	-0.00	0.00	-0.00	-0.00	-0.00	-0.00	0.00

Convolution r2[i+ 10\*(j+ 8\* 6)]

i/j	0	1	2	3	4	5	6	7
0	-0.00	0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
1	1.00	2.00	3.00	4.00	3.00	2.00	1.00	-0.00
2	2.75	5.50	8.25	11.00	8.25	5.50	2.75	-0.00
3	5.00	10.00	15.00	20.00	15.00	10.00	5.00	-0.00
4	3.50	7.00	10.50	14.00	10.50	7.00	3.50	-0.00
5	2.00	4.00	6.00	8.00	6.00	4.00	2.00	-0.00
6	0.75	1.50	2.25	3.00	2.25	1.50	0.75	-0.00
7	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00

Convolution r2[i+ 10\*(j+ 8\* 7)]

i/j	0	1	2	3	4	5	6	7
0	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
1	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
2	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
3	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
4	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
5	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
6	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
7	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00

---

## 6.12 CORRELATIONS

### 6.12.1 ASL\_qfcr2d, ASL\_pfcr2d Two-Dimensional Correlations

#### (1) Function

Assume that the two multiperiodic discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  of period  $(m_x, m_y)$  satisfying:

$$\begin{aligned} f(i_x, i_y) &= f(i_x + L_x m_x, i_y + L_y m_y), \\ g(j_x, j_y) &= g(j_x + L_x m_x, j_y + L_y m_y), \\ &(i_x, j_x = 0, \dots, m_x - 1; i_y, j_y = 0, \dots, m_y - 1) \end{aligned}$$

for arbitrary integers  $L_x$  and  $L_y$  take nonzero values within their basic periods only for  $(i_x, i_y) \in [0, n_x^{(f)} - 1] \times [0, n_y^{(f)} - 1]$  and  $(j_x, j_y) \in [0, n_x^{(g)} - 1] \times [0, n_y^{(g)} - 1]$ . Here,  $[0, a] \times [0, b]$  is the direct product region (region contained in the square for which the point  $(0, 0)$  and the point  $(a, b)$  are diagonal points) on the plane in which the plane coordinates  $(i, j)$  lie. At this time, ASL\_qfcr2d or ASL\_pfcr2d calculates the quantity  $\tilde{q}(k_x, k_y)$  obtained by shifting the discrete correlation  $q(k_x, k_y)$ , which is defined as follows:

$$\begin{aligned} q(k_x, k_y) &= \sum_{i_x=0}^{m_x-1} \sum_{i_y=0}^{m_y-1} f(i_x, i_y) g(k_x + i_x, k_y + i_y) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1) \end{aligned}$$

by  $(n_x^{(f)} - 1, n_y^{(f)} - 1)$  in the positive direction for  $(k_x, k_y)$ , respectively.  $\tilde{q}(k_x, k_y)$  is defined as follows:

$$\begin{aligned} \tilde{q}(k_x, k_y) &= q(k_x - (n_x^{(f)} - 1), k_y - (n_y^{(f)} - 1)) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1) \end{aligned}$$

Here,  $m_x = \min(n_x^{(f)} + n_x^{(g)} - 1, M_x)$  and  $m_y = \min(n_y^{(f)} + n_y^{(g)} - 1, M_y)$  and  $M_x$  and  $M_y$  are arbitrary integers satisfying  $M_x \geq \max(n_x^{(f)}, n_x^{(g)})$  and  $M_y \geq \max(n_y^{(f)}, n_y^{(g)})$ , respectively. The two-dimensional real Fourier transform of  $q(k_x, k_y)$  can also be obtained.

#### (2) Usage

Double precision:

```
ierr = ASL_qfcr2d (nx1, ny1, nx2, ny2, r1, lx1, ly1, r2, lx2, ly2, mx, my, isw, iwk, wk, nt);
```

Single precision:

```
ierr = ASL_pfcr2d (nx1, ny1, nx2, ny2, r1, lx1, ly1, r2, lx2, ly2, mx, my, isw, iwk, wk, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx1	I	1	Input	Number of effective data in $i_x$ direction $n_x^{(f)}$ for discrete function $f(i_x, i_y)$
2	ny1	I	1	Input	Number of effective data in $i_y$ direction $n_y^{(f)}$ for discrete function $f(i_x, i_y)$
3	nx2	I	1	Input	Number of effective data in $j_x$ direction $n_x^{(g)}$ for discrete function $g(j_x, j_y)$
4	ny2	I	1	Input	Number of effective data in $j_y$ direction $n_y^{(g)}$ for discrete function $g(j_x, j_y)$
5	r1	$\begin{cases} D* \\ R* \end{cases}$	lx1×ly1	Input	Values of discrete function $f(i_x, i_y)$ (See Note (a))
				Output	When $isw \geq 1$ , result of two-dimensional real Fourier transform of discrete function $f(i_x, i_y)$ (period $(M_x, M_y)$ )
6	lx1	I	1	Input	Adjustable dimension of array r1
7	ly1	I	1	Input	Second dimension of array r1
8	r2	$\begin{cases} D* \\ R* \end{cases}$	lx2×ly2	Input	Values of discrete function $g(j_x, j_y)$ (See Note (a))
				Output	Value of discrete function $\tilde{q}(k_x, k_y)$ or the two-dimensional real Fourier transform of $q(k_x, k_y)$ (See Note (b))
9	lx2	I	1	Input	Adjustable dimension of array r2
10	ly2	I	1	Input	Second dimension of array r2
11	mx	I	1	Input	Parameter $M_x$ corresponding to the period $(m_x, m_y)$ of discrete functions $f(i_x, i_y)$ , $g(j_x, j_y)$ , and $\tilde{q}(k_x, k_y)$ (See Note (c))
12	my	I	1	Input	Parameter $M_y$ corresponding to the period $(m_x, m_y)$ of discrete functions $f(i_x, i_y)$ , $g(j_x, j_y)$ , and $\tilde{q}(k_x, k_y)$ (See Note (c))

No.	Argument and Return Value	Type	Size	Input/Output	Contents
13	isw	I	1	Input	Processing switch (See Note (d)) isw= 0: Calculate the correlation according to the definition. isw= 1: Calculate the correlation according to the FFT method. isw= 2: Calculate the real Fourier transform of the correlation.
14	iwk	I*	See Contents	Work	Work area <b>Size:</b> 0 (When isw= 0) 40 (When isw $\geq$ 1)
15	wk	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	See Contents	Work	Work area <b>Size:</b> $nx2 \times ny2$ ((When isw= 0 and $nx2$ is odd) $(nx2 + 1) \times ny2$ (When isw= 0 and $nx2$ is even) $mx + 2 \times my + \max(lx1 \times ly1, lx2 \times ly2)$ (When isw $\geq$ 1)
16	nt	I	1	Input	Number of tasks to be generated
17	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, 1, 2\}$
- (b)  $nx1 > 1$  and  $ny1 > 1$
- (c)  $nx2 > 1$  and  $ny2 > 1$
- (d)  $mx \geq \text{MAX}(nx1, nx2)$  and  $my \geq \text{MAX}(ny1, ny2)$
- (e)  $lx1 \geq nx1$  and  $ly1 \geq ny1$  (When isw=0)  
 $lx1 \geq mx + 1$  and  $ly1 \geq my$  (When isw  $\geq$  1 and  $mx$  is odd)  
 $lx1 \geq mx + 2$  and  $ly1 \geq my$  (When isw  $\geq$  1 and  $mx$  is even)
- (f)  $lx2 \geq mx$  and  $ly2 \geq my$  (When isw=0)  
 $lx2 \geq mx + 1$  and  $ly2 \geq my$  (When isw  $\geq$  1 and  $mx$  is odd)  
 $lx2 \geq mx + 2$  and  $ly2 \geq my$  (When isw  $\geq$  1 and  $mx$  is even)
- (g)  $nt \geq 1$

(5) Error indicator (Return Value)

ier value	Meaning	Processing
0	Normal termination.	
1000	$mx < nx1 + nx2 - 1$ or $my < ny1 + ny2 - 1$	Overlapping occurred during the correlation calculation.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	
3050	Restriction (f) was not satisfied.	
3060	Restriction (g) was not satisfied.	

(6) Notes

- (a) The values of the discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  and the elements of arrays r1 and r2 are associated as follows.

$$\begin{aligned} f(i_x, i_y) &\leftrightarrow r1[i_x + lx1 * i_y] \\ g(j_x, j_y) &\leftrightarrow r2[j_x + lx2 * j_y] \end{aligned}$$

Here,  $i_x = 0, \dots, n_x^{(f)} - 1$ ;  $i_y = 0, \dots, n_y^{(f)} - 1$  and  $j_x = 0, \dots, n_x^{(g)} - 1$ ;  $j_y = 0, \dots, n_y^{(g)} - 1$ , and no values need be entered in other elements. **The adjustable dimensions of arrays r1 and r2 should be set so that  $lx1/2$ ,  $ly1$ ,  $lx2/2$ , and  $ly2$  are odd numbers to avoid bank conflict of main memory. Usually, when  $mx$ , for example, is a multiple of 4,  $lx1=mx+3$  is set.**

- (b) The values of the discrete correlation  $\tilde{q}(k_x, k_y)$  and the elements of array r2 are associated as follows.

$$\tilde{q}(k_x, k_y) \leftrightarrow r2[k_x + lx2 * k_y]$$

Here,  $k_x = 0, \dots, M_x - 1$ ;  $k_y = 0, \dots, M_y - 1$ . When  $isw=2$  is set to obtain the two-dimensional real Fourier transform  $Q(j_x, j_y)$  of the discrete correlation  $q(k_x, k_y)$ , which is defined as follows ( $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ ):

$$\begin{aligned} Q(j_x, j_y) &= \frac{1}{M_x M_y} \sum_{k_x=0}^{M_x-1} \sum_{k_y=0}^{M_y-1} q(k_x, k_y) e^{-2\pi\sqrt{-1}(\frac{j_x k_x}{M_x} + \frac{j_y k_y}{M_y})} \\ &(j_x = 0, \dots, \lfloor \frac{M_x}{2} \rfloor; j_y = 0, \dots, \lfloor \frac{M_y}{2} \rfloor) \end{aligned}$$

the following associations are made:

$$\begin{aligned} \Re\{Q(j_x, j_y)\} &\leftrightarrow r2[2 * j_x + lx2 * j_y] \\ \Im\{Q(j_x, j_y)\} &\leftrightarrow r2[2 * j_x + 1 + lx2 * j_y] \end{aligned}$$

In this case, note that the Fourier transform that is obtained is normalized. The remaining half period of the Fourier transform can be obtained from the symmetry of the real Fourier transform as follows:

$$\begin{aligned} Q(M_x - j_x, M_y - j_y)^* &= Q(j_x, j_y) \\ Q(M_x - j_x, j_y)^* &= Q(j_x, M_y - j_y) \end{aligned}$$

(Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .) Now,  $Q(j_x, j_y)$  can be thought of as an estimate of the cross spectrum of the original two functions for which the correlation is to be calculated. In this case,  $M_x$  and  $M_y$  should be thought of as  $M_x = n_x^{(f)} + n_x^{(g)}$  and  $M_y = n_y^{(f)} + n_y^{(g)}$ . In particular, if the original two functions for which the correlation is to be calculated are the same

function,  $Q(j_x, j_y)$  corresponds to the raw Fourier periodogram (estimate of the power spectrum), and  $Q(j_x, j_y)$  is a real number.

- (c) If  $mx \geq nx1 + nx2 - 1$  and  $my \geq ny1 + ny2 - 1$  are set, the correlation can be calculated without causing an overlap with the correlation of the next period. When  $mx > nx1 + nx2 - 1$  or  $my > ny1 + ny2 - 1$ , the following correspondences are made:

$$\tilde{q}(k_x, k_y) \leftrightarrow r2[k_x + lx2 * k_y]$$

and values that match 0.0 within the error range are stored in elements corresponding to  $k_x = nx1 + nx2 - 1, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y = ny1 + ny2 - 1, \dots, my - 1$ . When  $isw=0$ ,  $mx = nx1 + nx2 - 1$  and  $my = ny1 + ny2 - 1$  should be set. When  $isw \geq 1$ , the calculations can be performed more efficiently by setting a value for  $mx$  or  $my$  for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, if  $nx1=nx2=145$ , then when  $isw=0$ ,  $mx = 289(=17^2)$  should be set. However, when  $isw \geq 1$ , it is usually more efficient to set  $mx = 300(=2^2 \times 3 \times 5^2)$ ,  $mx = 320(=2^6 \times 5)$ ,  $mx = 384(=2^7 \times 3)$  or the like.

- (d) **Usually, the calculations can be performed more efficiently by setting  $isw=1$  to calculate the FFT correlation.** However, to conserve work area or if there is a restriction on the method of selecting the parameter  $mx$  or  $my$ , the calculations should be performed by setting  $isw=0$ .
- (e) To calculate the correlation of discrete functions for which the starting position of the nonzero portions are separated from the origin, first perform the calculations by shifting the functions so that the starting positions are at the origin, and then shift the calculation results again to obtain the final results more efficiently. For example, when the nonzero portions of the discrete functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  are the intervals  $[i_0, i_0 + n_x^{(f)} - 1]$  and  $[j_0, j_0 + n_x^{(g)} - 1]$  for  $i_x$  and  $j_x$ , respectively, let  $\hat{f}(i_x, i_y)$  and  $\hat{g}(j_x, j_y)$  be defined as follows:

$$\hat{f}(i_x, i_y) = f(i_x - i_0, i_y), \quad \hat{g}(j_x, j_y) = g(j_x - j_0, j_y)$$

and apply this function to  $\hat{f}(i_x, i_y)$  and  $\hat{g}(j_x, j_y)$ . Let  $\tilde{q}(k_x, k_y)$  represent the result that was obtained, and the correlation  $q(k_x, k_y)$  of the original functions  $f(i_x, i_y)$  and  $g(j_x, j_y)$  is given as follows:

$$q(k_x, k_y) = \tilde{q}(k_x - (j_0 - i_0) + (n_x^{(f)} - 1), k_y)$$

Therefore, even when  $i_0 = j_0 = 0$ , to consider the correlation  $q(k_x, k_y)$  that conforms to the normal definition, you must consider shifting the result by  $n_x^{(f)} - 1$  in the negative direction of  $k_x$  after applying this function or if you shift  $f(i_x, i_y)$  and  $g(j_x, j_y)$  in the negative directions of  $i_x$  and  $j_x$  by  $i_0$  and  $j_0$ , respectively, before calculating the discrete correlation, you must then shift the calculation result again by  $j_0 - i_0$  in the positive direction of  $k_x$ . Similarly, for  $i_y, j_y$ , and  $k_y$ .

- (f) The sampling interval squared multiplied by the discrete correlation calculated by this function is the square approximation (or approximation by using the trapezoidal formula) of the continuous correlation integral of a bandwidth-limited function. Therefore, to raise the approximation precision, you must take a smaller sampling interval and a larger number of sample data. To associate these results with a continuous correlation, it is easiest to let  $q(-n_x^{(f)}, k_y) = \tilde{q}(-1, k_y) = 0$  and  $q(k_x, -n_y^{(f)}) = \tilde{q}(k_x, -1) = 0$  and consider  $(n_x^{(f)} + n_x^{(g)})(n_y^{(f)} + n_y^{(g)})$  data of  $q(k_x, k_y)$  ( $k_x = -n_x^{(f)}, \dots, -1, 0, 1, \dots, n_x^{(g)} - 1$ ;  $k_y = -n_y^{(f)}, \dots, -1, 0, 1, \dots, n_y^{(g)} - 1$ ). Of course, this is the same as letting  $q(n_x^{(f)} + n_x^{(g)}, k_y) = \tilde{q}(n_x^{(g)}, k_y) = 0$  and  $q(k_x, n_y^{(f)} + n_y^{(g)}) = \tilde{q}(k_x, n_y^{(g)}) = 0$  and considering  $q(k_x, k_y)$  ( $k_x = -(n_x^{(f)} - 1), \dots, -1, 0, 1, \dots, n_x^{(g)}$ ;  $k_y = -(n_y^{(f)} - 1), \dots, -1, 0, 1, \dots, n_y^{(g)}$ ). In this case, the coordinate  $(0, 0)$  element usually is associated with  $q(0, 0)$ . However,

When  $isw=0$ ,  
then  $lx1 = nx1, ly1 = ny1, lx2 = mx, ly2 = my$ , and  
 $nwk = nx2 \times ny2$  (when  $nx2$  is odd) or  
 $nwk = (nx2 + 1) \times ny2$  (when  $nx2$  is even)  
When  $isw \geq 1$ ,  
then  $lx1=lx2=mx+1$  (when  $mx$  is odd) or  
 $lx1=lx2=mx+2$  (when  $mx$  is even),  
 $ly1=ly2=my$ , and  $nwk = mx + (lx1 + 2) \times my$ .

(7) **Example**

(a) Problem

Use the sampling interval  $\Delta$  to discretize the two finite waveforms defined by the following equations and calculate the discrete correlation.

$$f(x, y) = \begin{cases} x & ((x, y) \in [0, x_f] \times [0, y_f]) \\ 0 & (\text{Otherwise}) \end{cases}$$

$$g(x, y) = \begin{cases} x_g - x & ((x, y) \in [0, x_g] \times [0, y_g]) \\ 0 & (\text{Otherwise}) \end{cases}$$

(b) Input data

Sampling data

$$r1[i_x + lx1 * i_y] = f(i_x \Delta, i_y \Delta) \quad (i_x = 0, 1, \dots, nx1 - 1; i_y = 0, 1, \dots, ny1 - 1),$$

$$r2[j_x + lx2 * j_y] = g(j_x \Delta, j_y \Delta) \quad (j_x = 0, 1, \dots, nx2 - 1; j_y = 0, 1, \dots, ny2 - 1).$$

Here,  $\Delta = 0.5$ .

$nx1, ny1, nx2, ny2, mx, my$  and  $isw$ .

(c) Main program

```

/*      C interface example for ASL_qfcr2d */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int nx1;
    int ny1;
    int nx2;
    int ny2;
    double *r1;
    int m0=8;
    int lx1;
    int ly1;
    double *r2;
    int lx2;
    int ly2;
    int mx;
    int my;
    int isw;
    int *iwk;
    int niwk=40;
    double *wkw;
    int nwk;
    int ierr;
    int i,j;
    int nt = 2;
    double t;
    double dt=0.5;
    double xf=2.0,yf=2.0;
    double xg=2.0,yg=2.0;

    printf( "      *** ASL_qfcr2d ***\n" );
    printf( "\n      ** Input **\n\n" );

    isw=1;
    nx1=(int) xf/dt;
    ny1=(int) yf/dt;

```

```

nx2=(int) xg/dt;
ny2=(int) yg/dt;
mx=my*m0;
lx1=lx2=m0+2;
ly1=ly2=m0;
nwk=mx+2*my+lx2*my;

r1 = ( double * )malloc((size_t)( sizeof(double) * (lx1*ly1) ));
if( r1 == NULL )
{
    printf( "no enough memory for array r1\n" );
    return -1;
}
r2 = ( double * )malloc((size_t)( sizeof(double) * (lx2*ly2) ));
if( r2 == NULL )
{
    printf( "no enough memory for array r2\n" );
    return -1;
}
wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}
iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
if( iwk == NULL )
{
    printf( "no enough memory for array iwk\n" );
    return -1;
}

printf( "\t isw = %6d\n\t (nx1, ny1) = (%3d,%3d)\n",
        isw, nx1, ny1);
printf( "\t (nx2, ny2) = (%3d,%3d)\n",
        nx2, ny2);
printf( "\t (mx , my ) = (%3d,%3d)\n\n",
        mx, my);

for( j=0 ; j<ny1 ; j++ )
    for( i=0 ; i<nx1 ; i++ )
    {
        t=i*dt;
        r1[i+lx1*j]=t;
    }
for( j=0 ; j<ny2 ; j++ )
    for( i=0 ; i<nx2 ; i++ )
    {
        t=i*dt;
        r2[i+lx2*j]=xg-t;
    }
printf( "\tData r1[i+%3d*j]\n", lx1 );
printf( "\ti/j      0      1      2      3\n" );
printf( "\t-----\n" );
for( i=0 ; i<nx1 ; i++ )
{
    printf( "\t%3d", i );
    for( j=0 ; j<ny1 ; j++ )
        printf( "%8.3g", r1[i+lx1*j] );
    printf( "\n" );
}
printf( "\n" );
printf( "\tData r2[i+%3d*j]\n", lx2 );
printf( "\ti/j      0      1      2      3\n" );
printf( "\t-----\n" );
for( i=0 ; i<nx2 ; i++ )
{
    printf( "\t%3d", i );
    for( j=0 ; j<ny2 ; j++ )
        printf( "%8.3g", r2[i+lx2*j] );
    printf( "\n" );
}

ierr = ASL_qfcr2d(nx1, ny1, nx2, ny2, r1, lx1, ly1,
                 r2, lx2, ly2, mx, my, isw, iwk, wk, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

printf( "\tCorrelation r2[i+%3d*j]\n", lx2 );
printf( "\ti/j      0      1      2      3      4" );
printf( "\n      5      6      7\n" );
printf( "\t-----" );
printf( "-----\n" );
for( i=0 ; i<mx ; i++ )
{

```



```

    printf( "\t%2d", i );
    for( j=0 ; j<my ; j++ )
        printf( "%7.2lf", r2[i+lx2*j] );
    printf( "\n");
}
free( iwk );
free( wk );
free( r2 );
free( r1 );
return 0;
}

```

(d) Output results

```

*** ASL_qfcr2d ***

** Input **

isw =      1
(nx1, ny1) = ( 4,  4)
(nx2, ny2) = ( 4,  4)
(mx , my ) = ( 8,  8)

Data r1[i+ 10*j]
i/j      0      1      2      3
-----
0         0         0         0         0
1        0.5        0.5        0.5        0.5
2         1         1         1         1
3        1.5        1.5        1.5        1.5

Data r2[i+ 10*j]
i/j      0      1      2      3
-----
0         2         2         2         2
1        1.5        1.5        1.5        1.5
2         1         1         1         1
3        0.5        0.5        0.5        0.5

** Output **

ierr =      0
Correlation r2[i+ 10*j]
i/j      0      1      2      3      4      5      6      7
-----
0  3.00  6.00  9.00 12.00  9.00  6.00  3.00 -0.00
1  4.25  8.50 12.75 17.00 12.75  8.50  4.25 -0.00
2  4.00  8.00 12.00 16.00 12.00  8.00  4.00 -0.00
3  2.50  5.00  7.50 10.00  7.50  5.00  2.50 -0.00
4  1.00  2.00  3.00  4.00  3.00  2.00  1.00 -0.00
5  0.25  0.50  0.75  1.00  0.75  0.50  0.25  0.00
6  0.00  0.00  0.00 -0.00  0.00  0.00  0.00  0.00
7  0.00  0.00 -0.00 -0.00  0.00  0.00  0.00  0.00

```

## 6.12.2 ASL\_qfcr3d, ASL\_pfcr3d Three-Dimensional Correlations

### (1) Function

Assume that the two multiperiodic discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  of period  $(m_x, m_y, m_z)$  satisfying:

$$\begin{aligned} f(i_x, i_y, i_z) &= f(i_x + L_x m_x, i_y + L_y m_y, i_z + L_z m_z), \\ g(j_x, j_y, j_z) &= g(j_x + L_x m_x, j_y + L_y m_y, j_z + L_z m_z), \\ &(i_x, j_x = 0, \dots, m_x - 1; i_y, j_y = 0, \dots, m_y - 1; i_z, j_z = 0, \dots, m_z - 1) \end{aligned}$$

for arbitrary integers  $L_x, L_y$ , and  $L_z$  take nonzero values within their basic periods only for  $(i_x, i_y, i_z) \in [0, n_x^{(f)} - 1] \times [0, n_y^{(f)} - 1] \times [0, n_z^{(f)} - 1]$  and  $(j_x, j_y, j_z) \in [0, n_x^{(g)} - 1] \times [0, n_y^{(g)} - 1] \times [0, n_z^{(g)} - 1]$ . Here,  $[0, a] \times [0, b] \times [0, c]$  is the direct product region (region contained in the cube for which the point  $(0, 0, 0)$  and the point  $(a, b, c)$  are diagonal points) in the space in which the space coordinates  $(i, j, k)$  lie. At this time, ASL\_qfcr3d or ASL\_pfcr3d calculates the quantity  $\tilde{q}(k_x, k_y, k_z)$  obtained by shifting the discrete correlation  $q(k_x, k_y, k_z)$ , which is defined as follows:

$$\begin{aligned} q(k_x, k_y, k_z) &= \sum_{i_x=0}^{m_x-1} \sum_{i_y=0}^{m_y-1} \sum_{i_z=0}^{m_z-1} f(i_x, i_y, i_z) g(k_x + i_x, k_y + i_y, k_z + i_z) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1; k_z = 0, \dots, m_z - 1) \end{aligned}$$

by  $(n_x^{(f)} - 1, n_y^{(f)} - 1, n_z^{(f)} - 1)$  in the positive direction for  $(k_x, k_y, k_z)$ , respectively.  $\tilde{q}(k_x, k_y, k_z)$  is defined as follows:

$$\begin{aligned} \tilde{q}(k_x, k_y, k_z) &= q(k_x - (n_x^{(f)} - 1), k_y - (n_y^{(f)} - 1), k_z - (n_z^{(f)} - 1)) \\ &(k_x = 0, \dots, m_x - 1; k_y = 0, \dots, m_y - 1; k_z = 0, \dots, m_z - 1) \end{aligned}$$

Here,  $m_x = \min(n_x^{(f)} + n_x^{(g)} - 1, M_x)$ ,  $m_y = \min(n_y^{(f)} + n_y^{(g)} - 1, M_y)$ , and  $m_z = \min(n_z^{(f)} + n_z^{(g)} - 1, M_z)$  and  $M_x, M_y$ , and  $M_z$  are arbitrary integers satisfying  $M_x \geq \max(n_x^{(f)}, n_x^{(g)})$ ,  $M_y \geq \max(n_y^{(f)}, n_y^{(g)})$ , and  $M_z \geq \max(n_z^{(f)}, n_z^{(g)})$ , respectively. The three-dimensional real Fourier transform of  $q(k_x, k_y, k_z)$  can also be obtained.

### (2) Usage

Double precision:

```
ierr = ASL_qfcr3d (nx1, ny1, nz1, nx2, ny2, nz2, r1, lx1, ly1, lz1, r2, lx2, ly2, lz2, mx, my, mz,
                  isw, iwk, wk, nt);
```

Single precision:

```
ierr = ASL_pfcr3d (nx1, ny1, nz1, nx2, ny2, nz2, r1, lx1, ly1, lz1, r2, lx2, ly2, lz2, mx, my, mz,
                  isw, iwk, wk, nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx1	I	1	Input	Number of effective data in $i_x$ direction $n_x^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
2	ny1	I	1	Input	Number of effective data in $i_y$ direction $n_y^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
3	nz1	I	1	Input	Number of effective data in $i_z$ direction $n_z^{(f)}$ for discrete function $f(i_x, i_y, i_z)$
4	nx2	I	1	Input	Number of effective data in $j_x$ direction $n_x^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
5	ny2	I	1	Input	Number of effective data in $j_y$ direction $n_y^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
6	nz2	I	1	Input	Number of effective data in $j_z$ direction $n_z^{(g)}$ for discrete function $g(j_x, j_y, j_z)$
7	r1	$\begin{cases} D* \\ R* \end{cases}$	$lx1 \times ly1 \times lz1$	Input	Values of discrete function $f(i_x, i_y, i_z)$ (See Note (a))
				Output	When $isw \geq 1$ , result of three-dimensional real Fourier transform of discrete function $f(i_x, i_y, i_z)$ (period $(M_x, M_y, M_z)$ )
8	lx1	I	1	Input	Adjustable dimension of array r1
9	ly1	I	1	Input	Second dimension of array r1
10	lz1	I	1	Input	Third dimension of array r1
11	r2	$\begin{cases} D* \\ R* \end{cases}$	$lx2 \times ly2 \times lz2$	Input	Values of discrete function $g(j_x, j_y, j_z)$ (See Note (a))
				Output	Value of discrete function $\tilde{q}(k_x, k_y, k_z)$ or the three-dimensional real Fourier transform of $q(k_x, k_y, k_z)$ (See Note (b))
12	lx2	I	1	Input	Adjustable dimension of array r2
13	ly2	I	1	Input	Second dimension of array r2
14	lz2	I	1	Input	Third dimension of array r2
15	mx	I	1	Input	Parameter $M_x$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $\tilde{q}(k_x, k_y, k_z)$ (See Note (c))

No.	Argument and Return Value	Type	Size	Input/Output	Contents
16	my	I	1	Input	Parameter $M_y$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $\tilde{q}(k_x, k_y, k_z)$ (See Note (c))
17	mz	I	1	Input	Parameter $M_z$ corresponding to the period $(m_x, m_y, m_z)$ of discrete functions $f(i_x, i_y, i_z)$ , $g(j_x, j_y, j_z)$ , and $\tilde{q}(k_x, k_y, k_z)$ (See Note (c))
18	isw	I	1	Input	Processing switch (See Note (d)) isw= 0: Calculate the correlation according to the definition. isw= 1: Calculate the correlation according to the FFT method. isw= 2: Calculate the real Fourier transform of the correlation.
19	iwk	I*	See Contents	Work	Work area <b>Size:</b> 0 (When isw= 0) 60 (When isw $\geq$ 1)
20	wk	$\begin{Bmatrix} D^* \\ R^* \end{Bmatrix}$	See Contents	Work	Work area <b>Size:</b> $(nx2 + 1) \times (ny2 + 1) \times nz2$ (When isw= 0, nx2 is even and ny2 is even) $nx2 \times (ny2 + 1) \times nz2$ (When isw= 0, nx2 is odd and ny2 is even) $(nx2 + 1) \times ny2 \times nz2$ (When isw= 0, nx2 is even and ny2 is odd) $nx2 \times ny2 \times nz2$ (When isw= 0, nx2 is odd and ny2 is odd) $mx + 2 \times (my + mz) + \max(lx1 \times ly1 \times lz1, lx2 \times ly2 \times lz2)$ (When isw $\geq$ 1)
21	nt	I	1	Input	Number of tasks to be generated
22	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, 1, 2\}$
- (b)  $nx1 > 1$  and  $ny1 > 1$  and  $nz1 > 1$
- (c)  $nx2 > 1$  and  $ny2 > 1$  and  $nz2 > 1$
- (d)  $mx \geq \text{MAX}(nx1, nx2)$  and  $my \geq \text{MAX}(ny1, ny2)$  and  $mz \geq \text{MAX}(nz1, nz2)$
- (e)  $lx1 \geq nx1$  and  $ly1 \geq ny1$  and  $lz1 \geq nz1$  (when  $isw=0$ )  
 $lx1 \geq mx + 1$  and  $lx1$  is even and  $ly1 \geq my$  and  $lz1 \geq mz$  (when  $isw \geq 1$  and  $mx$  is odd)  
 $lx1 \geq mx + 2$  and  $lx1$  is even and  $ly1 \geq my$  and  $lz1 \geq mz$  (when  $isw \geq 1$  and  $mx$  is even)
- (f)  $lx2 \geq mx$  and  $ly2 \geq ny2$  and  $lz2 \geq nz2$  (when  $isw=0$ )  
 $lx2 \geq mx + 1$  and  $lx2$  is even and  $ly2 \geq my$  and  $lz2 \geq mz$  (when  $isw \geq 1$  and  $mx$  is odd)  
 $lx2 \geq mx + 2$  and  $lx2$  is even and  $ly2 \geq my$  and  $lz2 \geq mz$  (when  $isw \geq 1$  and  $mx$  is even)
- (g)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
1000	$mx < nx1 + nx2 - 1$ or $my < ny1 + ny2 - 1$ or $mz < nz1 + nz2 - 1$	Overlapping occurred during the correlation calculation.
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3020	Restriction (c) was not satisfied.	
3030	Restriction (d) was not satisfied.	
3040	Restriction (e) was not satisfied.	
3050	Restriction (f) was not satisfied.	
3060	Restriction (g) was not satisfied.	

(6) Notes

- (a) The values of the discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  and the elements of arrays r1 and r2 are associated as follows.

$$\begin{aligned} f(i_x, i_y, i_z) &\leftrightarrow r1[i_x + lx1 * (i_y + ly1 * i_z)] \\ g(j_x, j_y, j_z) &\leftrightarrow r2[j_x + lx2 * (j_y + ly2 * j_z)] \end{aligned}$$

Here,  $i_x = 0, \dots, n_x^{(f)} - 1$ ;  $i_y = 0, \dots, n_y^{(f)} - 1$ ;  $i_z = 0, \dots, n_z^{(f)} - 1$  and  $j_x = 0, \dots, n_x^{(g)} - 1$ ;  $j_y = 0, \dots, n_y^{(g)} - 1$ ;  $j_z = 0, \dots, n_z^{(g)} - 1$ , and no values need be entered in other elements. **The adjustable dimensions of arrays r1 and r2 should be set so that  $lx1/2$ ,  $ly1$ ,  $lz1$ ,  $lx2/2$ ,  $ly2$ , and  $lz2$  are odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within arrays r1 and r2. Usually, when  $mx$ , for example, is (a multiple of 4)+2,  $lx1=mx+4$  is set.**

- (b) The values of the discrete convolution  $\tilde{q}(k_x, k_y, k_z)$  and the elements of array r2 are associated as follows.

$$\tilde{q}(k_x, k_y, k_z) \leftrightarrow r2[k_x + lx2 * (k_y + ly2 * k_z)]$$

Here,  $k_x = 0, \dots, M_x - 1$ ;  $k_y = 0, \dots, M_y - 1$ ;  $k_z = 0, \dots, M_z - 1$ . When  $isw=2$  is set to obtain the three-dimensional real Fourier transform  $Q(j_x, j_y, j_z)$  of the discrete correlation  $q(k_x, k_y, k_z)$ , which is defined as follows ( $[x]$  represents the maximum integer that does not exceed  $x$ ):

$$\begin{aligned} Q(j_x, j_y, j_z) &= \frac{1}{M_x M_y M_z} \sum_{k_x=0}^{M_x-1} \sum_{k_y=0}^{M_y-1} \sum_{k_z=0}^{M_z-1} q(k_x, k_y, k_z) e^{-2\pi\sqrt{-1}(\frac{j_x k_x}{M_x} + \frac{j_y k_y}{M_y} + \frac{j_z k_z}{M_z})} \\ &\quad (j_x = 0, \dots, [\frac{M_x}{2}]; j_y = 0, \dots, [\frac{M_y}{2}]; j_z = 0, \dots, [\frac{M_z}{2}]) \end{aligned}$$

the following associations are made:

$$\begin{aligned} \Re\{Q(j_x, j_y, j_z)\} &\leftrightarrow r2[2 * j_x + lx2 * (j_y + ly2 * j_z)] \\ \Im\{Q(j_x, j_y, j_z)\} &\leftrightarrow r2[2 * j_x + 1 + lx2 * (j_y + ly2 * j_z)] \end{aligned}$$

In this case, note that the Fourier transform that is obtained is normalized. The remaining half period of the Fourier transform can be obtained from the symmetry of the real Fourier transform as follows:

$$\begin{aligned} Q(M_x - j_x, M_y - j_y, M_z - j_z)^* &= Q(j_x, j_y, j_z) \\ Q(M_x - j_x, j_y, j_z)^* &= Q(j_x, M_y - j_y, M_z - j_z) \\ Q(M_x - j_x, M_y - j_y, j_z)^* &= Q(j_x, j_y, M_z - j_z) \end{aligned}$$

(Here,  $z^*$  represents the conjugate complex number of the complex number  $z$ .) Now,  $Q(j_x, j_y, j_z)$  can be thought of as an estimate of the cross spectrum of the original two functions for which the correlation is to be calculated. In this case,  $M_x$ ,  $M_y$ , and  $M_z$  should be thought of as  $M_x = n_x^{(f)} + n_x^{(g)}$ ,  $M_y = n_y^{(f)} + n_y^{(g)}$ , and  $M_z = n_z^{(f)} + n_z^{(g)}$ . In particular, if the original two functions for which the correlation is to be calculated are the same function,  $Q(j_x, j_y, j_z)$  corresponds to the raw Fourier periodogram (estimate of the power spectrum), and  $Q(j_x, j_y, j_z)$  is a real number.

- (c) If  $mx \geq nx1 + nx2 - 1$  and  $my \geq ny1 + ny2 - 1$  and  $mz \geq nz1 + nz2 - 1$  are set, the correlation can be calculated without causing an overlap with the correlation of the next period. When  $mx > nx1 + nx2 - 1$  or  $my > ny1 + ny2 - 1$  or  $mz > nz1 + nz2 - 1$ , the following correspondences are made:

$$\tilde{q}(k_x, k_y) \leftrightarrow r2[k_x + lx2 * (k_y + ly2 * k_z)]$$

and values that match 0.0 within the error range are stored in elements corresponding to  $k_x = nx1 + nx2 - 1, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$ ;  $k_z = 0, \dots, mz - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y =$

$ny1 + ny2 - 1, \dots, my - 1$ ;  $k_z = 0, \dots, mz - 1$  or  $k_x = 0, \dots, mx - 1$ ;  $k_y = 0, \dots, my - 1$ ;  $k_z = nz1 + nz2 - 1, \dots, mz - 1$ . When  $isw=0$ ,  $mx = nx1 + nx2 - 1$ ,  $my = ny1 + ny2 - 1$ , and  $mz = nz1 + nz2 - 1$  should be set. When  $isw \geq 1$ , the calculations can be performed more efficiently by setting a value for  $mx$ ,  $my$  or  $mz$  for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, if  $nx1=nx2=145$ , then when  $isw=0$ ,  $mx = 289(=17^2)$  should be set. However, when  $isw \geq 1$ , it is usually more efficient to set  $mx = 300(=2^2 \times 3 \times 5^2)$ ,  $mx = 320(=2^6 \times 5)$ ,  $mx = 384(=2^7 \times 3)$  or the like.

- (d) **Usually, the calculations can be performed more efficiently by setting  $isw=1$  to calculate the FFT correlation.** However, to conserve work area or if there is a restriction on the method of selecting the parameter  $mx$ ,  $my$  or  $mz$ , the calculations should be performed by setting  $isw=0$ .
- (e) To calculate the correlation of discrete functions the starting position of the nonzero portions are separated from the origin, first perform the calculations by shifting the functions so that the starting positions are at the origin, and then shift the calculation results again to obtain the final results more efficiently. For example, when the nonzero portions of the discrete functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  are the intervals  $[i_0, i_0 + n_x^{(f)} - 1]$  and  $[j_0, j_0 + n_x^{(g)} - 1]$  for  $i_x$  and  $j_x$ , respectively, let  $\hat{f}(i_x, i_y, i_z)$  and  $\hat{g}(j_x, j_y, j_z)$  be defined as follows:

$$\hat{f}(i_x, i_y, i_z) = f(i_x - i_0, i_y, i_z), \quad \hat{g}(j_x, j_y, j_z) = g(j_x - j_0, j_y, j_z)$$

and apply this function to  $\hat{f}(i_x, i_y, i_z)$  and  $\hat{g}(j_x, j_y, j_z)$ . Let  $\tilde{q}(k_x, k_y, k_z)$  represent the result that was obtained, and the correlation  $q(k_x, k_y, k_z)$  of the original functions  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  is given as follows:

$$q(k_x, k_y, k_z) = \tilde{q}(k_x - (j_0 - i_0) + (n_x^{(f)} - 1), k_y, k_z)$$

Therefore, even when  $i_0 = j_0 = 0$ , to consider the correlation  $q(k_x, k_y, k_z)$  that conforms to the normal definition, you must consider shifting the result by  $n_x^{(f)} - 1$  in the negative direction of  $k_x$  after applying this function or if you shift  $f(i_x, i_y, i_z)$  and  $g(j_x, j_y, j_z)$  in the negative directions of  $i_x$  and  $j_x$  by  $i_0$  and  $j_0$ , respectively, before calculating the discrete correlation, you must then shift the calculation result again by  $j_0 - i_0$  in the positive direction of  $k_x$ . Similarly, for  $i_y, j_y$ , and  $k_y$  and  $i_z, j_z$ , and  $k_z$ .

- (f) The sampling interval cubed multiplied by the discrete correlation calculated by this function is the square approximation (or approximation by using the trapezoidal formula) of the continuous correlation integral of a bandwidth-limited function. Therefore, to raise the approximation precision, you must take a smaller sampling interval and a larger number of sample data. To associate these results with a continuous correlation it is easiest to let  $q(-n_x^{(f)}, k_y, k_z) = \tilde{q}(-1, k_y, k_z) = 0$ ,  $q(k_x, -n_y^{(f)}, k_z) = \tilde{q}(k_x, -1, k_z) = 0$ , and  $q(k_x, k_y, -n_z^{(f)}) = \tilde{q}(k_x, k_y, -1) = 0$  and consider  $(n_x^{(f)} + n_x^{(g)})(n_y^{(f)} + n_y^{(g)})(n_z^{(f)} + n_z^{(g)})$  data of  $q(k_x, k_y, k_z)$  ( $k_x = -n_x^{(f)}, \dots, -1, 0, 1, \dots, n_x^{(g)} - 1$ ;  $k_y = -n_y^{(f)}, \dots, -1, 0, 1, \dots, n_y^{(g)} - 1$ ;  $k_z = -n_z^{(f)}, \dots, -1, 0, 1, \dots, n_z^{(g)} - 1$ ). Of course, this is the same as letting  $q(n_x^{(f)} + n_x^{(g)}, k_y, k_z) = \tilde{q}(n_x^{(g)}, k_y, k_z) = 0$ ,  $q(k_x, n_y^{(f)} + n_y^{(g)}, k_z) = \tilde{q}(k_x, n_y^{(g)}, k_z) = 0$ , and  $q(k_x, k_y, n_z^{(f)} + n_z^{(g)}) = \tilde{q}(k_x, k_y, n_z^{(g)}) = 0$  and considering  $q(k_x, k_y, k_z)$  ( $k_x = -(n_x^{(f)} - 1), \dots, -1, 0, 1, \dots, n_x^{(g)}$ ;  $k_y = -(n_y^{(f)} - 1), \dots, -1, 0, 1, \dots, n_y^{(g)}$ ;  $k_z = -(n_z^{(f)} - 1), \dots, -1, 0, 1, \dots, n_z^{(g)}$ ). In this case, the coordinate  $(0, 0, 0)$  element usually is associated with  $q(0, 0, 0)$ . However,

When  $isw=0$ ,

then  $lx1 = nx1, ly1 = ny1, lz1 = nz1, lx2 = mx, ly2 = my, lz2 = mz$ , and  
 $nwk = (nx2 + 1) \times (ny2 + 1) \times nz2$  (when  $nx2$  is even and  $ny2$  is even) or  
 $nwk = nx2 \times (ny2 + 1) \times nz2$  (when  $nx2$  is odd and  $ny2$  is even) or  
 $nwk = (nx2 + 1) \times ny2 \times nz2$  (when  $nx2$  is even and  $ny2$  is odd) or  
 $nwk = nx2 \times ny2 \times nz2$  (when  $nx2$  is odd and  $ny2$  is odd)

When  $isw \geq 1$

$lx1=lx2=mx+1$  (when  $mx$  is odd) or

$lx1=lx2=mx+2$  (when  $mx$  is even),

$ly1=ly2=my$ ,  $lz1=lz2=mz$ , and  $nwk = mx + 2 \times (my + mz) + lx1 \times my \times mz$ .

(7) **Example**

(a) **Problem**

Use the sampling interval  $\Delta$  to discretize the two finite waveforms defined by the following equations and calculate the discrete correlation.

$$f(x, y, z) = \begin{cases} x & ((x, y, z) \in [0, x_f] \times [0, y_f] \times [0, z_f]) \\ 0 & \text{(Otherwise)} \end{cases}$$

$$g(x, y, z) = \begin{cases} x_g - x & ((x, y, z) \in [0, x_g] \times [0, y_g] \times [0, z_g]) \\ 0 & \text{(Otherwise)} \end{cases}$$

(b) **Input data**

Sampling data

$r1[i_x + lx1 * (i_y + ly1 * i_z)] = f(i_x \Delta, i_y \Delta, i_z \Delta)$  ( $i_x = 0, 1, \dots, nx1 - 1$ ;  $i_y = 0, 1, \dots, ny1 - 1$ ;  $i_z = 0, 1, \dots, nz1 - 1$ ),

$r2[j_x + lx2 * (j_y + ly2 * j_z)] = g(j_x \Delta, j_y \Delta, j_z \Delta)$  ( $j_x = 0, 1, \dots, nx2 - 1$ ;  $j_y = 0, 1, \dots, ny2 - 1$ ;  $j_z = 0, 1, \dots, nz2 - 1$ ).

Here,  $\Delta = 0.5$ .

$nx1, ny1, nz1, nx2, ny2, nz2, mx, my, mz$  and  $isw$ .

(c) **Main program**

```
/*      C interface example for ASL_qfcr3d */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    int nx1;
    int ny1;
    int nz1;
    int nx2;
    int ny2;
    int nz2;
    double *r1;
    int m0=8;
    int lx1;
    int ly1;
    int lz1;
    double *r2;
    int lx2;
    int ly2;
    int lz2;
    int mx;
    int my;
    int mz;
    int isw;
    int *iwk;
    int niwk=60;
    double *wk;
    int nwk;
    int ierr;
    int i,j,k;
    int nt = 2;
    double t;
    double dt=0.5;
    double xf=2.0,yf=2.0,zf=2.0;
    double xg=2.0,yg=2.0,zg=2.0;

    printf( "      *** ASL_qfcr3d ***\n" );
    printf( "\n      ** Input **\n\n" );

    isw=1;
    nx1=(int) xf/dt;
```



```

ny1=(int) yf/dt;
nz1=(int) zf/dt;
nx2=(int) xg/dt;
ny2=(int) yg/dt;
nz2=(int) zg/dt;
mx=my=mz=m0;
lx1=lx2=(m0+2)/2*2;
ly1=ly2=my;
lz1=lz2=mz;
nwk=mx+2*(my+mz)+lx2*my*mz;

r1 = ( double * )malloc((size_t)( sizeof(double) * (lx1*ly1*lz1) ));
if( r1 == NULL )
{
    printf( "no enough memory for array r1\n" );
    return -1;
}
r2 = ( double * )malloc((size_t)( sizeof(double) * (lx2*ly2*lz2) ));
if( r2 == NULL )
{
    printf( "no enough memory for array r2\n" );
    return -1;
}
wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}
iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
if( iwkw == NULL )
{
    printf( "no enough memory for array iwkw\n" );
    return -1;
}

printf( "\t isw = %6d\n\t (nx1, ny1, nz1) = (%3d,%3d,%3d)\n",
        isw, nx1, ny1, nz1);
printf( "\t (nx2, ny2, nz2) = (%3d,%3d,%3d)\n",
        nx2, ny2, nz2);
printf( "\t (mx , my , mz ) = (%3d,%3d,%3d)\n\n",
        mx, my, mz);

for( k=0 ; k<nz1 ; k++ )
    for( j=0 ; j<ny1 ; j++ )
        for( i=0 ; i<nx1 ; i++ )
            {
                t=i*dt;
                r1[i+lx1*(j+ly1*k)]=t;
            }
for( k=0 ; k<nz2 ; k++ )
    for( j=0 ; j<ny2 ; j++ )
        for( i=0 ; i<nx2 ; i++ )
            {
                t=i*dt;
                r2[i+lx2*(j+ly2*k)]=xg-t;
            }
for( k=0 ; k<nz1 ; k++ )
{
    printf( "\tData r1[i+%3d*(j+%3d*%3d)]\n", lx1, ly1, k );
    printf( "\t i/j      0      1      2      3\n" );
    printf( "\t-----\n" );
    for( i=0 ; i<nx1 ; i++ )
    {
        printf( "\t%3d", i );
        for( j=0 ; j<ny1 ; j++ )
            printf( "%8.3g", r1[i+lx1*(j+ly1*k)] );
        printf( "\n" );
    }
    printf( "\n" );
}
for( k=0 ; k<nz2 ; k++ )
{
    printf( "\tData r2[i+%3d*(j+%3d*%3d)]\n", lx2, ly2, k );
    printf( "\t i/j      0      1      2      3\n" );
    printf( "\t-----\n" );
    for( i=0 ; i<nx2 ; i++ )
    {
        printf( "\t%3d", i );
        for( j=0 ; j<ny2 ; j++ )
            printf( "%8.3g", r2[i+lx2*(j+ly2*k)] );
        printf( "\n" );
    }
    printf( "\n" );
}

ierr = ASL_qfcr3d(nx1, ny1, nz1, nx2, ny2, nz2,

```

```

    r1, lx1, ly1, lz1, r2, lx2, ly2, lz2,
    mx, my, mz, isw, iwk, wk, nt);

printf( "\n    ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

for( k=0 ; k<mz ; k++ )
{
    printf( "\tCorrelation r2[i+%3d*(j+%3d*%3d)]\n", lx2, ly2, k );
    printf( "\ti/j  0    1    2    3    4" );
    printf( "      5    6    7\n" );
    printf( "\t-----" );
    printf( "-----\n" );
    for( i=0 ; i<mx ; i++ )
    {
        printf( "\t%2d", i );
        for( j=0 ; j<my ; j++ )
            printf( "%7.2lf", r2[i+lx2*(j+ly2*k)] );
        printf( "\n" );
    }
    printf( "\n" );
}
free( iwk );
free( wk );
free( r2 );
free( r1 );

return 0;
}

```

(d) Output results

```

*** ASL_qfcr3d ***

** Input **

isw =      1
(nx1, ny1, nz1) = ( 4, 4, 4)
(nx2, ny2, nz2) = ( 4, 4, 4)
(mx , my , mz ) = ( 8, 8, 8)

Data r1[i+ 10*(j+ 8* 0)]
i/j      0      1      2      3
-----
0         0         0         0         0
1        0.5        0.5        0.5        0.5
2         1         1         1         1
3        1.5        1.5        1.5        1.5

Data r1[i+ 10*(j+ 8* 1)]
i/j      0      1      2      3
-----
0         0         0         0         0
1        0.5        0.5        0.5        0.5
2         1         1         1         1
3        1.5        1.5        1.5        1.5

Data r1[i+ 10*(j+ 8* 2)]
i/j      0      1      2      3
-----
0         0         0         0         0
1        0.5        0.5        0.5        0.5
2         1         1         1         1
3        1.5        1.5        1.5        1.5

Data r1[i+ 10*(j+ 8* 3)]
i/j      0      1      2      3
-----
0         0         0         0         0
1        0.5        0.5        0.5        0.5
2         1         1         1         1
3        1.5        1.5        1.5        1.5

Data r2[i+ 10*(j+ 8* 0)]
i/j      0      1      2      3
-----
0         2         2         2         2
1        1.5        1.5        1.5        1.5
2         1         1         1         1
3        0.5        0.5        0.5        0.5

Data r2[i+ 10*(j+ 8* 1)]
i/j      0      1      2      3
-----
0         2         2         2         2
1        1.5        1.5        1.5        1.5
2         1         1         1         1
3        0.5        0.5        0.5        0.5

```

Data r2[i+ 10\*(j+ 8\* 2)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

Data r2[i+ 10\*(j+ 8\* 3)]

i/j	0	1	2	3
0	2	2	2	2
1	1.5	1.5	1.5	1.5
2	1	1	1	1
3	0.5	0.5	0.5	0.5

\*\* Output \*\*

ierr = 0  
 Correlation r2[i+ 10\*(j+ 8\* 0)]

i/j	0	1	2	3	4	5	6	7
0	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
1	4.25	8.50	12.75	17.00	12.75	8.50	4.25	-0.00
2	4.00	8.00	12.00	16.00	12.00	8.00	4.00	-0.00
3	2.50	5.00	7.50	10.00	7.50	5.00	2.50	-0.00
4	1.00	2.00	3.00	4.00	3.00	2.00	1.00	-0.00
5	0.25	0.50	0.75	1.00	0.75	0.50	0.25	0.00
6	0.00	0.00	-0.00	-0.00	0.00	0.00	0.00	0.00
7	0.00	0.00	-0.00	-0.00	-0.00	0.00	0.00	0.00

Correlation r2[i+ 10\*(j+ 8\* 1)]

i/j	0	1	2	3	4	5	6	7
0	6.00	12.00	18.00	24.00	18.00	12.00	6.00	-0.00
1	8.50	17.00	25.50	34.00	25.50	17.00	8.50	-0.00
2	8.00	16.00	24.00	32.00	24.00	16.00	8.00	-0.00
3	5.00	10.00	15.00	20.00	15.00	10.00	5.00	-0.00
4	2.00	4.00	6.00	8.00	6.00	4.00	2.00	-0.00
5	0.50	1.00	1.50	2.00	1.50	1.00	0.50	-0.00
6	0.00	0.00	0.00	-0.00	0.00	0.00	0.00	0.00
7	0.00	0.00	-0.00	-0.00	0.00	0.00	0.00	0.00

Correlation r2[i+ 10\*(j+ 8\* 2)]

i/j	0	1	2	3	4	5	6	7
0	9.00	18.00	27.00	36.00	27.00	18.00	9.00	-0.00
1	12.75	25.50	38.25	51.00	38.25	25.50	12.75	-0.00
2	12.00	24.00	36.00	48.00	36.00	24.00	12.00	-0.00
3	7.50	15.00	22.50	30.00	22.50	15.00	7.50	-0.00
4	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
5	0.75	1.50	2.25	3.00	2.25	1.50	0.75	-0.00
6	0.00	0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
7	0.00	-0.00	0.00	-0.00	0.00	0.00	-0.00	0.00

Correlation r2[i+ 10\*(j+ 8\* 3)]

i/j	0	1	2	3	4	5	6	7
0	12.00	24.00	36.00	48.00	36.00	24.00	12.00	-0.00
1	17.00	34.00	51.00	68.00	51.00	34.00	17.00	-0.00
2	16.00	32.00	48.00	64.00	48.00	32.00	16.00	-0.00
3	10.00	20.00	30.00	40.00	30.00	20.00	10.00	-0.00
4	4.00	8.00	12.00	16.00	12.00	8.00	4.00	0.00
5	1.00	2.00	3.00	4.00	3.00	2.00	1.00	0.00
6	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	0.00
7	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00

Correlation r2[i+ 10\*(j+ 8\* 4)]

i/j	0	1	2	3	4	5	6	7
0	9.00	18.00	27.00	36.00	27.00	18.00	9.00	-0.00
1	12.75	25.50	38.25	51.00	38.25	25.50	12.75	-0.00
2	12.00	24.00	36.00	48.00	36.00	24.00	12.00	-0.00
3	7.50	15.00	22.50	30.00	22.50	15.00	7.50	-0.00
4	3.00	6.00	9.00	12.00	9.00	6.00	3.00	-0.00
5	0.75	1.50	2.25	3.00	2.25	1.50	0.75	0.00
6	0.00	0.00	0.00	0.00	0.00	0.00	0.00	0.00
7	0.00	0.00	0.00	-0.00	0.00	-0.00	0.00	0.00

Correlation r2[i+ 10\*(j+ 8\* 5)]

i/j	0	1	2	3	4	5	6	7
0	6.00	12.00	18.00	24.00	18.00	12.00	6.00	-0.00
1	8.50	17.00	25.50	34.00	25.50	17.00	8.50	-0.00
2	8.00	16.00	24.00	32.00	24.00	16.00	8.00	-0.00
3	5.00	10.00	15.00	20.00	15.00	10.00	5.00	-0.00
4	2.00	4.00	6.00	8.00	6.00	4.00	2.00	-0.00
5	0.50	1.00	1.50	2.00	1.50	1.00	0.50	0.00
6	0.00	0.00	0.00	-0.00	0.00	0.00	0.00	0.00
7	0.00	0.00	-0.00	-0.00	0.00	0.00	0.00	0.00

Correlation r2[i+ 10\*(j+ 8\* 6)]

i/j	0	1	2	3	4	5	6	7
0	3.00	6.00	9.00	12.00	9.00	6.00	3.00	0.00
1	4.25	8.50	12.75	17.00	12.75	8.50	4.25	-0.00
2	4.00	8.00	12.00	16.00	12.00	8.00	4.00	-0.00
3	2.50	5.00	7.50	10.00	7.50	5.00	2.50	-0.00
4	1.00	2.00	3.00	4.00	3.00	2.00	1.00	0.00
5	0.25	0.50	0.75	1.00	0.75	0.50	0.25	0.00
6	0.00	-0.00	-0.00	0.00	0.00	0.00	0.00	0.00
7	0.00	-0.00	-0.00	-0.00	0.00	0.00	0.00	0.00

Correlation  $r2[i+10*(j+8*7)]$

i/j	0	1	2	3	4	5	6	7
0	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	0.00	0.00
1	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
2	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
3	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
4	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00	-0.00
5	-0.00	-0.00	0.00	0.00	0.00	0.00	-0.00	-0.00
6	-0.00	-0.00	0.00	0.00	0.00	0.00	0.00	0.00
7	-0.00	0.00	-0.00	0.00	0.00	0.00	0.00	0.00

## 6.13 POWER SPECTRUM ANALYSIS

### 6.13.1 ASL\_qfps2d, ASL\_pfps2d

#### Two-Dimensional Fourier Periodograms

(1) **Function**

ASL\_qfps2d or ASL\_pfps2d obtains the (modified) Fourier periodogram of the series  $u_{j_x, j_y}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ). The Fourier periodogram  $p_{k_x, k_y}$  is defined by the following equation.

$$p_{k_x, k_y} = \frac{\left| \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} w_{j_x}^{(x)} w_{j_y}^{(y)} u_{j_x, j_y} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y}\right)} \right|^2}{n_x n_y \beta} \quad (k_x = 0, 1, \dots, \lfloor \frac{n_x}{2} \rfloor; k_y = 0, 1, \dots, n_y - 1)$$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .  $w_{j_x}^{(x)}$  and  $w_{j_y}^{(y)}$  are the truncation functions (data windows). For a raw Fourier periodogram,  $w_{j_x}^{(x)} = w_{j_y}^{(y)} = 1$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ) and  $\beta = n_x n_y$  are set, and for a modified periodogram  $\beta$  is set as follows:

$$\beta = \begin{cases} \left( \sum_{j_x=0}^{n_x-1} (w_{j_x}^{(x)})^2 \right) \left( \sum_{j_y=0}^{n_y-1} (w_{j_y}^{(y)})^2 \right) & \text{(when a power modification expression according} \\ & \text{to a data window is used)} \\ n_x n_y & \text{(Otherwise)} \end{cases}$$

The periodogram  $p_{k_x, k_y}$  corresponds to a half period (period  $(n_x, n_y)$ ) and the remainder is obtained from the relationship as follows.

$$\begin{aligned} p_{n_x - k_x, n_y - k_y} &= p_{k_x, k_y} \\ p_{n_x - k_x, k_y} &= p_{k_x, n_y - k_y} \end{aligned}$$

Also, the total power of the corresponding series is as follows.

$$\frac{\sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \{u_{j_x, j_y}\}^2}{n_x n_y}$$

(2) **Usage**

Double precision:

ierr = ASL\_qfps2d (nx, ny, r, lx, ly, isw, iwk, wk, nt);

Single precision:

ierr = ASL\_pfps2d (nx, ny, r, lx, ly, isw, iwk, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Length $n_x$ in the $j_x$ direction of series $u_{j_x, j_y}$ (See Note (d))
2	ny	I	1	Input	Length $n_y$ in the $j_y$ direction of series $u_{j_x, j_y}$ (See Note (d))
3	r	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	lx×ly	Input	Values of series $u_{j_x, j_y}$ (See Note (a))
				Output	Values of Fourier periodogram $p_{k_x, k_y}$ of series $u_{j_x, j_y}$ (See Notes (b) and (c))
4	lx	I	1	Input	Adjustable dimension of array r
5	ly	I	1	Input	Second dimension of array r
6	isw	I	1	Input	Processing switch (See Note (e)) isw= 0: Calculate the raw Fourier periodogram isw=±1: Calculate the periodogram using a user-defined data window isw=±2: Calculate the periodogram using the Hanning window isw=±3: Calculate the periodogram using the Bartlett window isw=±4: Calculate the periodogram using the Welch window isw=±5: Calculate the periodogram using the Parzen window To use a power modification expression according to a data window, set isw > 0; otherwise, set isw < 0.
7	iwk	I*	40	Work	Work area
8	wk	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	See Contents	Work	Work area When isw=±1, enter the values of the user-defined data window. (See Note (e)) <b>Size:</b> $nx + 2 \times ny + lx \times ly$
9	nt	I	1	Input	Number of tasks to be generated
10	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, \pm 1, \pm 2, \pm 3, \pm 4, \pm 5\}$
- (b)  $nx > 1$  and  $ny > 1$
- (c)  $lx \geq nx + 1$  and  $ly \geq ny$  (when  $nx$  is odd)  
 $lx \geq nx + 2$  and  $ly \geq ny$  (when  $nx$  is even)
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3030	Restriction (c) was not satisfied.	
3040	Restriction (d) was not satisfied.	
4000	When $isw=1$ , the user-defined data window was $w_{j_x}^{(x)} = 0$ ( $j_x = 0, \dots, n_x - 1$ ).	
4010	When $isw=1$ , the user-defined data window was $w_{j_y}^{(y)} = 0$ ( $j_y = 0, \dots, n_y - 1$ ).	

(6) **Notes**

- (a) The elements of array  $r$  and the values of the series  $u_{j_x, j_y}$  are associated as follows.

$$u_{j_x, j_y} \leftrightarrow r[j_x + lx * j_y]$$

Here,  $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ , and no values need be entered in other elements. **The adjustable dimensions of array  $r$  should be set so that  $lx/2$  and  $ly$  are odd numbers to avoid bank conflict of main memory. Usually, when  $nx$ , for example, is (a multiple of 4)+2,  $lx=nx+4$  is set.**

- (b) The values of the Fourier periodogram  $p_{k_x, k_y}$  are associated as follows with the elements of array  $r$ .

$$p(k_x, k_y) \leftrightarrow r[k_x + lx * k_y] \quad (k_x = 0, \dots, \lfloor \frac{nx}{2} \rfloor; k_y = 0, \dots, n_y - 1)$$

$\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

- (c) The frequencies  $(\xi_{k_x}, \eta_{k_y})$  corresponding to obtained Fourier periodogram  $p_{k_x, k_y}$  ( $k_x = 0, 1, \dots, n_x - 1$ ;  $k_y = 0, 1, \dots, n_y - 1$ ) are given as follows.

$$\xi_{k_x} = \frac{k_x}{n_x \Delta} \quad (k_x = 0, 1, \dots, \lfloor \frac{n_x}{2} \rfloor)$$

$$\eta_{k_y} = \begin{cases} \frac{k_y}{n_y \Delta} & (k_y = 0, 1, \dots, \lfloor \frac{n_y}{2} \rfloor) \\ \frac{k_y - n_y}{n_y \Delta} & (k_y = \lfloor \frac{n_y}{2} \rfloor + 1, \dots, n_y - 1) \end{cases}$$

Where  $\Delta$  is the sampling interval.

- (d) The calculations can be performed more efficiently by setting the length  $nx$  and  $ny$  of the series  $u_{j_x, j_y}$  to a value for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, rather than setting  $nx = 289(=17^2)$ , it is usually more efficient to set  $nx = 300(=2^2 \times 3 \times 5^2)$ ,  $nx = 320(=2^6 \times 5)$ ,  $nx = 384(=2^7 \times 3)$  or the

like. When the number of data cannot be increased, adjust  $n_x$  by supplying the required number of zeros at the end of the data to perform the calculations.

- (e) The truncation function (data window) can be changed as follows according to the value of the processing switch  $isw$ .

$$w_j = \begin{cases} \begin{cases} \sin^2(\pi v_j) & isw = \pm 2 \text{ (Hanning window)} \\ 1 - |2v_j - 1| & isw = \pm 3 \text{ (Bartlett window)} \\ 1 - (2v_j - 1)^2 & isw = \pm 4 \text{ (Welch window)} \end{cases} \\ \begin{cases} \begin{cases} 16v_j^3 & 0 \leq v_j < \frac{1}{4} \\ 1 - 6v_j(v_j - 1)^2 & \frac{1}{4} \leq v_j \leq \frac{1}{2} \\ 1 - 6v_j(v_{n-j+1} - 1)^2 & \frac{1}{2} \leq v_j \leq \frac{3}{4} \\ 16v_{n-j+1}^3 & \frac{3}{4} \leq v_j < 1 \end{cases} \end{cases} \end{cases} \quad isw = \pm 5 \text{ (Parzen window)}$$

Here,  $v_j = \frac{j}{n}$ , and  $j = j_x$  and  $n = n_x$  are set for  $w_{j_x}^{(x)}$  and  $j = j_y$  and  $n = n_y$  are set for  $w_{j_y}^{(y)}$ . Therefore, when the data windows shown above are used, the elements  $u_{0,j_y}$  and  $u_{j_x,0}$  of the series  $u_{j_x,j_y}$  do not affect the calculation of the modified periodogram. To avoid this situation, you should specify values for  $n_x$  and  $n_y$  that are larger by 1 than the lengths of the series for which you actually want to calculate the periodogram and set the effective data in elements corresponding to 1 and after for  $j_x$  and  $j_y$ . The data windows are represented as follows as time (or space) domain functions that are nonzero only for  $|x| \leq 1$ .

$$w(x) = \begin{cases} \frac{1 + \cos \pi x}{2} = \cos^2 \frac{\pi x}{2} & \text{Hanning window} \\ 1 - |x| & \text{Bartlett window} \\ 1 - x^2 & \text{Welch window} \\ \begin{cases} 1 - 6x^2 + 6|x|^3 & |x| \leq \frac{1}{2} \\ 2(1 - |x|)^3 & \frac{1}{2} \leq |x| \leq 1 \end{cases} & \text{Parzen window} \end{cases}$$

Also, to use user-defined data window values  $w_{j_x}^{(x)}$  and  $w_{j_y}^{(y)}$ , set  $isw = \pm 1$ , set the values in work array  $wk$  as follows:

$$wk[j_x] = w_{j_x}^{(x)} \quad (j_x = 0, \dots, n_x - 1), \quad wk[n_x + j_y] = w_{j_y}^{(y)} \quad (j_y = 0, \dots, n_y - 1)$$

and then call this function.

- (f) From its definition, the raw periodogram should be regarded as an approximation of a discrete Fourier transform of the autocorrelation function. Since the effective data length of the autocorrelation function of a discrete function having effective number of data  $n$  is  $2n - 1$ , approximating the power spectrum of a general function by a raw periodogram corresponds to truncating the function by using a square truncation function  $w(k)$  for which one period is given as follows.

$$w(k) = \begin{cases} 1 & k = 0, 1, \dots, n - 1 \\ 0 & \text{Otherwise} \end{cases}$$

When the frequency is  $f$  for the Fourier transform of the square function, a  $\frac{\sin f}{f}$  type function form is assumed having a sidelobe that is not small around the central frequency. Therefore, when a periodic function is sampled, for example, by simply truncating it using a width that is not an integer multiple of one period, since the raw periodogram will be the convolution of the Fourier transform of the periodic function for which the power spectrum is to be obtained and the  $\frac{\sin f}{f}$  type function in the frequency domain, an excess frequency component called **leakage** occurs. To suppress this kind of leakage, simple truncation is not performed, and a truncation function having a small sidelobe in the frequency domain, such as the Hanning window, is used. However, in general, the more the leakage is suppressed, the



more the result of the discrete Fourier transform widens and blurs. Therefore, when estimating the power spectrum, you must select a suitable truncation function according to your objectives, that is, according to whether the spectral width or the central frequency is to be the problem, for example.

- (g) To raise the resolution (sampling interval in the frequency domain)  $\frac{1}{nT}$  of the discrete Fourier transform, you should increase the number of sample data  $n$  or increase the sampling interval  $T$ . However, to raise the precision of the power spectrum estimate while holding the sampling interval and resolution fixed, a technique is often used of taking  $m$  groups of samples for which the number of samples is  $n$ , obtaining the modified periodogram for each of those  $m$  groups, and then taking the average of those values. In this case, a technique is also proposed in which the  $m$  groups of sample data are taken from the series so that they overlap. For details, refer to the Reference Bibliography.
- (h) When obtaining the power spectrum, the property related to the frequency transition of the Fourier transform, that is, the multiplication by  $e^{2\pi\sqrt{-1}f_0t}$  in the time (or space) domain, is associated with the shifting of the frequency by  $f_0$  in the frequency domain, and a technique is often used of reducing the number of data points required for the calculation in which the central frequency of the power spectrum is shifted in advance, using the property that the function shape does not change. This kind of operation is known as **modulation**. However,  $lx=nx+1$  (when  $nx$  is odd) or  $lx=nx+2$  (when  $nx$  is even) and  $ly=ny$ .

#### (7) Example

##### (a) Problem

Use the sampling interval  $\Delta$  to discretize the waveform defined by the following equation and estimate the power spectrum by calculating the Fourier periodogram.

$$f(x, y) = \cos 2\pi f_1 x + \cos 2\pi f_2 y$$

##### (b) Input data

Sampling data

$$r[j_x + lx * j_y] = f(j_x \Delta, j_y \Delta) \quad (j_x = 0, 1, \dots, nx - 1; j_y = 0, 1, \dots, ny - 1).$$

Here,  $\Delta = 0.5$ .

$nx$ ,  $ny$  and  $isw$ .

##### (c) Main program

```
/*      C interface example for ASL_qfps2d */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    int n0=8, isw0=4;
    int nx;
    int ny;
    double *r;
    int lx;
    int ly;
    int isw;
    int *iwk;
    int niwk=40;
    double *wk;
    int nwk;
    int ierr;
    int i, j, m, nd2, is;
    int nt = 2;
    double *p, t, tx, ty, dt, dfx, dfy, f0, f1, f2;

    printf( "      *** ASL_qfps2d ***\n" );
    printf( "\n      ** Input **\n\n" );
```

```

nx=n0;
ny=n0;
lx=n0+2;
ly=ny;
nwk=nx+2*ny+lx*ly;

r = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*(isw0+2)) ));
if( r == NULL )
{
    printf( "no enough memory for array r\n" );
    return -1;
}
wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}
p = ( double * )malloc((size_t)( sizeof(double) * (isw0+2) ));
if( p == NULL )
{
    printf( "no enough memory for array p\n" );
    return -1;
}
iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
if( iwk == NULL )
{
    printf( "no enough memory for array iwk\n" );
    return -1;
}

printf( "\t isw=0, 2 to %6d\n", isw0+1 );
printf( "\t nx=%6d\n\t ny=%6d\n\n", nx,ny );
dt=0.5;
f0=1.0/(2.0*dt);
f1=0.62*f0;
f2=0.14*f0;
nd2=(int) (nx+1)/2;
dfx=1.0/(dt*nx);
dfy=1.0/(dt*ny);
p[isw0+1]=0.0;
for( j=0 ; j<ny ; j++ )
{
    ty=(double) j*dt;
    for( i=0 ; i<nx ; i++ )
    {
        tx=(double) i*dt;
        t=cos(2.0*M_PI*f1*tx)+cos(2.0*M_PI*f2*ty);
        r[i+lx*(j+ly*(isw0+1))]=t;
        p[isw0+1] += (t*t);
    }
}
p[isw0+1] /= (double) (nx*ny);
printf( "\tTime series data r[i+%3d*j]\n", lx);
printf( " i/j\n");
for( j=0 ; j<ny ; j++ )
    printf( "%9d", j);
printf( "\n" );
printf( " -----");
printf( "-----\n");
for( i=0 ; i<nx ; i++ )
{
    printf( "%5d", i );
    for( j=0 ; j<ny ; j++ )
        printf( "%9.4lf", r[i+lx*(j+ly*(isw0+1))] );
    printf( "\n" );
}
printf( "\n");
printf( "\tTime domain power =%9.4lf\n", p[isw0+1]);
printf( "\tSignal frequency =( %9.4lf, %9.4lf)\n", f1, f2);
is=0;
for( isw=0 ; isw<=isw0 ; isw++ )
{
    for( j=0 ; j<ny ; j++ )
        for( i=0 ; i<nx ; i++ )
            r[i+lx*(j+ly*isw)]=r[i+lx*(j+ly*(isw0+1))];
    if ( isw != 0 )
        is=isw+1;

    ierr = ASL_qfps2d(nx, ny, &r[lx*ly*isw], lx, ly, is, iwk, wk, nt);
    p[isw]=0.0;
    if (nx%2==0)
    {
        m=nd2-1;
        for( j=0 ; j<ny ; j++ )
            for( i=1 ; i<m ; i++ )

```

```

        p[isw]+=2.0*r[i+lx*(j+ly*isw)];
    for( j=0 ; j<ny ; j++ )
        p[isw]+=r[lx*(j+ly*isw)]+r[m+lx*(j+ly*isw)];
    }
    else
    {
        m=nd2;
        for( j=0 ; j<ny ; j++ )
            for( i=1 ; i<m ; i++ )
                p[isw]+=2.0*r[i+lx*(j+ly*isw)];
        for( j=0 ; j<ny ; j++ )
            p[isw]+=r[lx*(j+ly*isw)];
    }
}

printf( "\n      ** Output **\n\n" );
printf( "\tierr = %6d\n", ierr );

isw=0;
printf( "\t(Modified) periodogram (Raw)\n" );
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
printf( "  x/y-freq");
for( j=(ny+1)/2 ; j<ny ; j++ )
    printf( "%8.2lf", (j-ny)*dfy );
for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( "  -----");
printf( "  -----");
for( i=0 ; i<nd2 ; i++ )
{
    printf( "  %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    printf( "\n" );
}
printf( "\n");

isw=1;
printf( "\t(Modified) periodogram (Hanning)\n" );
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
printf( "  x/y-freq");
for( j=(ny+1)/2 ; j<ny ; j++ )
    printf( "%8.2lf", (j-ny)*dfy );
for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( "  -----");
printf( "  -----");
for( i=0 ; i<nd2 ; i++ )
{
    printf( "  %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    printf( "\n" );
}
printf( "\n");

isw=2;
printf( "\t(Modified) periodogram (Bartlett)\n" );
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
printf( "  x/y-freq");
for( j=(ny+1)/2 ; j<ny ; j++ )
    printf( "%8.2lf", (j-ny)*dfy );
for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( "  -----");
printf( "  -----");
for( i=0 ; i<nd2 ; i++ )
{
    printf( "  %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    printf( "\n" );
}

```

```

}
printf( "\n");

isw=3;
printf( "\t(Modified) periodogram (Welch)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
printf( "  x/y-freq");
for( j=(ny+1)/2 ; j<ny ; j++ )
    printf( "%8.2lf", (j-ny)*dfy );
for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( "  -----");
printf( "-----\n");
for( i=0 ; i<nd2 ; i++ )
{
    printf( " %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    printf( "\n" );
}
printf( "\n");

isw=4;
printf( "\t(Modified) periodogram (Parzen)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
printf( "  x/y-freq");
for( j=(ny+1)/2 ; j<ny ; j++ )
    printf( "%8.2lf", (j-ny)*dfy );
for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( "  -----");
printf( "-----\n");
for( i=0 ; i<nd2 ; i++ )
{
    printf( " %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*isw)] );
    printf( "\n" );
}

free( iwk );
free( p );
free( wk );
free( r );

return 0;
}

```

(d) Output results

```

*** ASL_qfps2d ***

** Input **

isw=0, 2 to      5
nx=             8
ny=             8

Time series data r[i+ 10*j]
-----
i/j   0      1      2      3      4      5      6      7
0  2.0000  1.9048  1.6374  1.2487  0.8126  0.4122  0.1237  0.0020
1  0.6319  0.5367  0.2693 -0.1194 -0.5555 -0.9559 -1.2444 -1.3662
2  0.2710  0.1759 -0.0915 -0.4803 -0.9163 -1.3168 -1.6053 -1.7270
3  1.9048  1.8097  1.5423  1.1535  0.7174  0.3170  0.0285 -0.0932
4  1.0628  0.9676  0.7002  0.3115 -0.1246 -0.5250 -0.8135 -0.9352
5  0.0489 -0.0462 -0.3136 -0.7024 -1.1384 -1.5388 -1.8274 -1.9491
6  1.6374  1.5423  1.2748  0.8861  0.4500  0.0496 -0.2389 -0.3606
7  1.4818  1.3866  1.1192  0.7304  0.2944 -0.1060 -0.3946 -0.5163

Time domain power = 1.0626
Signal frequency =( 0.6200, 0.1400)

** Output **

ierr = 0
(Modified) periodogram (Raw)
Frequency domain power= 0.9717
x/y-freq  -1.00  -0.75  -0.50  -0.25  0.00  0.25  0.50  0.75
-----
0.00  0.0158  0.0188  0.0350  0.2150  0.0218  0.2150  0.0350  0.0188

```

0.25	0.0000	0.0000	0.0000	0.0000	0.0239	0.0000	0.0000	0.0000
0.50	0.0000	0.0000	0.0000	0.0000	0.1352	0.0000	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0781	0.0000	0.0000	0.0000
(Modified) periodogram (Hanning)								
Frequency domain power= 0.5980								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0001	0.0054	0.0632	0.0105	0.0632	0.0054	0.0001
0.25	0.0000	0.0000	0.0013	0.0095	0.0056	0.0236	0.0013	0.0000
0.50	0.0000	0.0000	0.0000	0.0204	0.0814	0.0204	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0205	0.0820	0.0205	0.0000	0.0000
(Modified) periodogram (Bartlett)								
Frequency domain power= 0.5835								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0009	0.0820	0.0109	0.0820	0.0009	0.0000
0.25	0.0000	0.0000	0.0002	0.0122	0.0025	0.0178	0.0002	0.0000
0.50	0.0000	0.0005	0.0000	0.0156	0.0855	0.0156	0.0000	0.0005
0.75	0.0000	0.0004	0.0000	0.0095	0.0762	0.0191	0.0000	0.0004
(Modified) periodogram (Welch)								
Frequency domain power= 0.7072								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0001	0.1263	0.0124	0.1263	0.0001	0.0000
0.25	0.0000	0.0000	0.0000	0.0140	0.0014	0.0127	0.0000	0.0000
0.50	0.0002	0.0003	0.0010	0.0195	0.1065	0.0054	0.0009	0.0003
0.75	0.0002	0.0003	0.0008	0.0064	0.0941	0.0142	0.0009	0.0003
(Modified) periodogram (Parzen)								
Frequency domain power= 0.4909								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0002	0.0093	0.0253	0.0070	0.0253	0.0093	0.0002
0.25	0.0000	0.0001	0.0022	0.0022	0.0127	0.0279	0.0064	0.0001
0.50	0.0000	0.0000	0.0006	0.0171	0.0558	0.0323	0.0030	0.0000
0.75	0.0000	0.0000	0.0013	0.0206	0.0485	0.0214	0.0014	0.0000

### 6.13.2 ASL\_qfps3d, ASL\_pfps3d Three-Dimensional Fourier Periodograms

(1) **Function**

ASL\_qfps3d or ASL\_pfps3d obtains the (modified) Fourier periodogram of the series  $u_{j_x, j_y, j_z}$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ). The Fourier periodogram  $p_{k_x, k_y, k_z}$  is defined by the following equation.

$$p_{k_x, k_y, k_z} = \frac{\left| \sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \sum_{j_z=0}^{n_z-1} w_{j_x}^{(x)} w_{j_y}^{(y)} w_{j_z}^{(z)} u_{j_x, j_y, j_z} e^{-2\pi\sqrt{-1}\left(\frac{j_x k_x}{n_x} + \frac{j_y k_y}{n_y} + \frac{j_z k_z}{n_z}\right)} \right|^2}{n_x n_y n_z \beta}$$

$(k_x = 0, 1, \dots, \lfloor \frac{n_x}{2} \rfloor; k_y = 0, 1, \dots, n_y - 1; k_z = 0, 1, \dots, n_z - 1)$

Here,  $\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .  $w_{j_x}^{(x)}$ ,  $w_{j_y}^{(y)}$  and  $w_{j_z}^{(z)}$  are the truncation functions (data windows). For a raw Fourier periodogram,  $w_{j_x}^{(x)} = w_{j_y}^{(y)} = w_{j_z}^{(z)} = 1$  ( $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$ ) and  $\beta = n_x n_y n_z$  are set, and for a modified periodogram  $\beta$  is set as follows:

$$\beta = \begin{cases} \left( \sum_{j_x=0}^{n_x-1} (w_{j_x}^{(x)})^2 \right) \left( \sum_{j_y=0}^{n_y-1} (w_{j_y}^{(y)})^2 \right) \left( \sum_{j_z=0}^{n_z-1} (w_{j_z}^{(z)})^2 \right) & \text{(when a power modification expression according} \\ & \text{to a data window is used)} \\ n_x n_y n_z & \text{(Otherwise)} \end{cases}$$

The periodogram  $p_{k_x, k_y, k_z}$  corresponds to a half period (period  $(n_x, n_y, n_z)$ ) and the remainder is obtained from the relationship as follows.

$$\begin{aligned} p_{n_x - k_x, n_y - k_y, n_z - k_z} &= p_{k_x, k_y, k_z} \\ p_{n_x - k_x, k_y, k_z} &= p_{k_x, n_y - k_y, n_z - k_z} \\ p_{n_x - k_x, n_y - k_y, k_z} &= p_{k_x, k_y, n_z - k_z} \end{aligned}$$

Also, the total power of the corresponding series is as follows.

$$\frac{\sum_{j_x=0}^{n_x-1} \sum_{j_y=0}^{n_y-1} \sum_{j_z=0}^{n_z-1} \{u_{j_x, j_y, j_z}\}^2}{n_x n_y n_z}$$

(2) **Usage**

Double precision:

ierr = ASL\_qfps3d (nx, ny, nz, r, lx, ly, lz, isw, iwk, wk, nt);

Single precision:

ierr = ASL\_pfps3d (nx, ny, nz, r, lx, ly, lz, isw, iwk, wk, nt);

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
 R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	nx	I	1	Input	Length $n_x$ in the $j_x$ direction of series $u_{j_x, j_y, j_z}$ (See Note (d))
2	ny	I	1	Input	Length $n_y$ in the $j_y$ direction of series $u_{j_x, j_y, j_z}$ (See Note (d))
3	nz	I	1	Input	Length $n_z$ in the $j_z$ direction of series $u_{j_x, j_y, j_z}$ (See Note (d))
4	r	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	lx×ly×lz	Input	Values of series $u_{j_x, j_y, j_z}$ (See Note (a))
				Output	Values of Fourier periodogram $p_{k_x, k_y, k_z}$ of series $u_{j_x, j_y, j_z}$ (See Notes (b) and (c))
5	lx	I	1	Input	Adjustable dimension of array r
6	ly	I	1	Input	Second dimension of array r
7	lz	I	1	Input	Third dimension of array r
8	isw	I	1	Input	Processing switch (See Note (e)) isw= 0: Calculate the raw Fourier periodogram isw=±1: Calculate the periodogram using a user-defined data window isw=±2: Calculate the periodogram using the Hanning window isw=±3: Calculate the periodogram using the Bartlett window isw=±4: Calculate the periodogram using the Welch window isw=±5: Calculate the periodogram using the Parzen window To use a power modification expression according to a data window, set isw > 0; otherwise, set isw < 0.
9	iwk	I*	60	Work	Work area
10	wk	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	See Contents	Work	Work area When isw=±1, enter the values of the user-defined data window. (See Note (e)) <b>Size:</b> $n_x + 2 \times (n_y + n_z) + l_x \times l_y \times l_z$
11	nt	I	1	Input	Number of tasks to be generated
12	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

- (a)  $isw \in \{0, \pm 1, \pm 2, \pm 3, \pm 4, \pm 5\}$
- (b)  $nx > 1$  and  $ny > 1$  and  $nz > 1$
- (c)  $lx \geq nx + 1$  and  $lx$  is even and  $ly \geq ny$  and  $lz \geq nz$  (when  $nx$  is odd)  
 $lx \geq nx + 2$  and  $lx$  is even and  $ly \geq ny$  and  $lz \geq nz$  (when  $nx$  is even)
- (d)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ier value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3010	Restriction (b) was not satisfied.	
3030	Restriction (c) was not satisfied.	
3040	Restriction (d) was not satisfied.	
4000	When $isw=1$ , the user-defined data window was $w_{j_x}^{(x)} = 0$ ( $j_x = 0, \dots, n_x - 1$ ).	
4010	When $isw=1$ , the user-defined data window was $w_{j_y}^{(y)} = 0$ ( $j_y = 0, \dots, n_y - 1$ ).	
4020	When $isw=1$ , the user-defined data window was $w_{j_z}^{(z)} = 0$ ( $j_z = 0, \dots, n_z - 1$ ).	

(6) **Notes**

- (a) The elements of array  $r$  and the values of the series  $u_{j_x, j_y, j_z}$  are associated as follows.

$$u_{j_x, j_y, j_z} \leftrightarrow r[j_x + lx * (j_y + ly * j_z)]$$

Here,  $j_x = 0, \dots, n_x - 1$ ;  $j_y = 0, \dots, n_y - 1$ ;  $j_z = 0, \dots, n_z - 1$  and no values need be entered in other elements. **The adjustable dimensions of array  $r$  should be set so that  $lx/2$ ,  $ly$ , and  $lz$  are odd numbers to avoid bank conflict of main memory. Also, to increase speed, calculations are executed even for elements outside areas where data is set within array  $r$ . Usually, when  $nx$ , for example, is (a multiple of 4)+2,  $lx=nx+4$  is set.**

- (b) The values of the Fourier periodogram  $p_{k_x, k_y, k_z}$  are associated as follows with the elements of array  $r$ .

$$p(k_x, k_y, k_z) \leftrightarrow r[k_x + lx * (k_y + ly * k_z)]$$

$$(k_x = 0, \dots, \lfloor \frac{nx}{2} \rfloor; k_y = 0, \dots, n_y - 1; k_z = 0, \dots, n_z - 1)$$

$\lfloor x \rfloor$  represents the maximum integer that does not exceed  $x$ .

- (c) The frequencies  $(\xi_{k_x}, \eta_{k_y}, \zeta_{k_z})$  corresponding to obtained Fourier periodogram  $p_{k_x, k_y, k_z}$  ( $k_x = 0, 1, \dots, \lfloor \frac{nx}{2} \rfloor$ ;  $k_y = 0, 1, \dots, n_y - 1$ ;  $k_z = 0, \dots, n_z - 1$ ) are given as follows.

$$\xi_{k_x} = \frac{k_x}{n_x \Delta} \quad (k_x = 0, 1, \dots, \lfloor \frac{n_x}{2} \rfloor)$$

$$\eta_{k_y} = \begin{cases} \frac{k_y}{n_y \Delta} & (k_y = 0, 1, \dots, \lfloor \frac{n_y}{2} \rfloor) \\ \frac{k_y - n_y}{n_y \Delta} & (k_y = \lfloor \frac{n_y}{2} \rfloor + 1, \dots, n_y - 1) \end{cases}$$

$$\zeta_{k_z} = \begin{cases} \frac{k_z}{n_z \Delta} & (k_z = 0, 1, \dots, \lfloor \frac{n_z}{2} \rfloor) \\ \frac{k_z - n_z}{n_z \Delta} & (k_z = \lfloor \frac{n_z}{2} \rfloor + 1, \dots, n_z - 1) \end{cases}$$



Where  $\Delta$  is the sampling interval.

- (d) The calculations can be performed more efficiently by setting the length  $n_x$ ,  $n_y$  and  $n_z$  of the series  $u_{j_x, j_y, j_z}$  to a value for which the mixed radix FFT algorithm operates effectively (multiples of 2, 3, 5, etc., which are the mixed radix values of FFT). For example, rather than setting  $n_x = 289 (=17^2)$ , it is usually more efficient to set  $n_x = 300 (=2^2 \times 3 \times 5^2)$ ,  $n_x = 320 (=2^6 \times 5)$ ,  $n_x = 384 (=2^7 \times 3)$  or the like. When the number of data cannot be increased, adjust  $n_x$  by supplying the required number of zeros at the end of the data to perform the calculations.
- (e) The truncation function (data window) can be changed as follows according to the value of the processing switch  $isw$ .

$$w_j = \begin{cases} \begin{cases} \sin^2(\pi v_j) & isw = \pm 2 \text{ (Hanning window)} \\ 1 - |2v_j - 1| & isw = \pm 3 \text{ (Bartlett window)} \\ 1 - (2v_j - 1)^2 & isw = \pm 4 \text{ (Welch window)} \end{cases} \\ \begin{cases} \begin{cases} 16v_j^3 & 0 \leq v_j < \frac{1}{4} \\ 1 - 6v_j(v_j - 1)^2 & \frac{1}{4} \leq v_j \leq \frac{1}{2} \\ 1 - 6v_j(v_{n-j+1} - 1)^2 & \frac{1}{2} \leq v_j \leq \frac{3}{4} \\ 16v_{n-j+1}^3 & \frac{3}{4} \leq v_j < 1 \end{cases} \end{cases} \end{cases} \quad isw = \pm 5 \text{ (Parzen window)}$$

Here,  $v_j = \frac{j}{n}$ , and  $j = j_x$  and  $n = n_x$  are set for  $w_{j_x}^{(x)}$ ,  $j = j_y$  and  $n = n_y$  are set for  $w_{j_y}^{(y)}$ , and  $j = j_z$  and  $n = n_z$  are set for  $w_{j_z}^{(z)}$ . Therefore, when the data windows shown above are used, the elements  $u_{0, j_y, j_z}$ ,  $u_{j_x, 0, j_z}$  and  $u_{j_x, j_y, 0}$  of the series  $u_{j_x, j_y, j_z}$  do not affect the calculation of the modified periodogram. To avoid this situation, you should specify values for  $n_x$ ,  $n_y$  and  $n_z$  that are larger by 1 than the lengths of the series for which you actually want to calculate the periodogram and set the effective data in elements corresponding to 1 and after for  $j_x$ ,  $j_y$  and  $j_z$ . The data windows are represented as follows as time (or space) domain functions that are nonzero only for  $|x| \leq 1$ .

$$w(x) = \begin{cases} \frac{1 + \cos \pi x}{2} = \cos^2 \frac{\pi x}{2} & \text{Hanning window} \\ 1 - |x| & \text{Bartlett window} \\ 1 - x^2 & \text{Welch window} \\ \begin{cases} 1 - 6x^2 + 6|x|^3 & |x| \leq \frac{1}{2} \\ 2(1 - |x|)^3 & \frac{1}{2} \leq |x| \leq 1 \end{cases} & \text{Parzen window} \end{cases}$$

Also, to use user-defined data window values  $w_{j_x}^{(x)}$ ,  $w_{j_y}^{(y)}$  and  $w_{j_z}^{(z)}$ , set  $isw = \pm 1$ , set the values in work array  $wk$  as follows:

$$wk[j_x] = w_{j_x}^{(x)} \quad (j_x = 0, \dots, n_x - 1), \quad wk[n_x + j_y] = w_{j_y}^{(y)} \quad (j_y = 0, \dots, n_y - 1), \\ wk[n_x + n_y + j_z] = w_{j_z}^{(z)} \quad (j_z = 0, \dots, n_z - 1)$$

and then call this function.

- (f) From its definition, the raw periodogram should be regarded as an approximation of a discrete Fourier transform of the autocorrelation function. Since the effective data length of the autocorrelation function of a discrete function having effective number of data  $n$  is  $2n - 1$ , approximating the power spectrum of a general function by a raw periodogram corresponds to truncating the function by using a square truncation function  $w(k)$  for which one period is given as follows.

$$w(k) = \begin{cases} 1 & k = 0, 1, \dots, n - 1 \\ 0 & \text{Otherwise} \end{cases}$$

When the frequency is  $f$  for the Fourier transform of the square function, a  $\frac{\sin f}{f}$  type function form is assumed having a sidelobe that is not small around the central frequency. Therefore, when a periodic

function is sampled, for example, by simply truncating it using a width that is not an integer multiple of one period, since the raw periodogram will be the convolution of the Fourier transform of the periodic function for which the power spectrum is to be obtained and the  $\frac{\sin f}{f}$  type function in the frequency domain, an excess frequency component called **leakage** occurs. To suppress this kind of leakage, simple truncation is not performed, and a truncation function having a small sidelobe in the frequency domain, such as the Hanning window, is used. However, in general, the more the leakage is suppressed, the more the result of the discrete Fourier transform widens and blurs. Therefore, when estimating the power spectrum, you must select a suitable truncation function according to your objectives, that is, according to whether the spectral width or the central frequency is to be the problem, for example.

- (g) To raise the resolution (sampling interval in the frequency domain)  $\frac{1}{nT}$  of the discrete Fourier transform, you should increase the number of sample data  $n$  or increase the sampling interval  $T$ . However, to raise the precision of the power spectrum estimate while holding the sampling interval and resolution fixed, a technique is often used of taking  $m$  groups of samples for which the number of samples is  $n$ , obtaining the modified periodogram for each of those  $m$  groups, and then taking the average of those values. In this case, a technique is also proposed in which the  $m$  groups of sample data are taken from the series so that they overlap. For details, refer to the Reference Bibliography.
- (h) When obtaining the power spectrum, the property related to the frequency transition of the Fourier transform, that is, the multiplication by  $e^{2\pi\sqrt{-1}f_0t}$  in the time (or space) domain, is associated with the shifting of the frequency by  $f_0$  in the frequency domain, and a technique is often used of reducing the number of data points required for the calculation in which the central frequency of the power spectrum is shifted in advance, using the property that the function shape does not change. This kind of operation is known as **modulation**. However,  $lx=nx+1$  (when  $nx$  is odd) or  $lx=nx+2$  (when  $nx$  is even)  
 $ly=ny$  and  $lz=nz$ .

(7) **Example**

(a) Problem

Use the sampling interval  $\Delta$  to discretize the waveform defined by the following equation and estimate the power spectrum by calculating the Fourier periodogram.

$$f(x, y, z) = \cos 2\pi f_1 x + \cos 2\pi f_2 y + \cos 2\pi f_3 z$$

(b) Input data

Sampling data

$$r[j_x + lx*(j_y + ly*j_z)] = f(j_x\Delta, j_y\Delta, j_z\Delta) \quad (j_x = 0, 1, \dots, nx-1; j_y = 0, 1, \dots, ny-1; j_z = 0, 1, \dots, nz-1).$$

Here,  $\Delta = 0.5$ .

$nx, ny, nz$  and  $isw$ .

(c) Main program

```
/*      C interface example for ASL_qfps3d */
#include <stdio.h>
#include <stdlib.h>
#include <math.h>
#include <asl.h>

int main()
{
    int n0=8, isw0=4;
    int nx;
    int ny;
    int nz;
    double *r;
    int lx;
    int ly;
```

```

int lz;
int isw;
int *iwk;
int niwk=60;
double *wk;
int nwk;
int ierr;
int i,j,k,m,nd2,is;
int nt = 2;
double *p,t,tx,ty,tz,dt,dfx,dfy,dfz,f0,f1,f2,f3;

printf( "      *** ASL_qfps3d ***\n" );
printf( "\n      ** Input **\n\n" );

nx=n0;
ny=n0;
nz=n0;
lx=(n0+2)/2*2;
ly=ny;
lz=nz;
nwk=nx+2*(ny+nz)+lx*ly*lz;

r = ( double * )malloc((size_t)( sizeof(double) * (lx*ly*lz*(isw0+2)) ));
if( r == NULL )
{
    printf( "no enough memory for array r\n" );
    return -1;
}
wk = ( double * )malloc((size_t)( sizeof(double) * nwk ));
if( wk == NULL )
{
    printf( "no enough memory for array wk\n" );
    return -1;
}
p = ( double * )malloc((size_t)( sizeof(double) * (isw0+2) ));
if( p == NULL )
{
    printf( "no enough memory for array p\n" );
    return -1;
}
iwk = ( int * )malloc((size_t)( sizeof(int) * niwk ));
if( iwkw == NULL )
{
    printf( "no enough memory for array iwkw\n" );
    return -1;
}

printf( "\t isw=0, 2 to %6d\n", isw0+1 );
printf( "\t nx=%6d\n\t ny=%6d\n\t nz=%6d\n\n", nx,ny,nz );
dt=0.5;
f0=1.0/(2.0*dt);
f1=0.62*f0;
f2=0.14*f0;
f3=0.55*f0;
nd2=(int) (nx+1)/2;
dfx=1.0/(dt*nx);
dfy=1.0/(dt*ny);
dfz=1.0/(dt*nz);
p[isw0+1]=0.0;
for( k=0 ; k<nz ; k++ )
{
    tz=(double) k*dt;
    for( j=0 ; j<ny ; j++ )
    {
        ty=(double) j*dt;
        for( i=0 ; i<nx ; i++ )
        {
            tx=(double) i*dt;
            t=cos(2.0*M_PI*f1*tx)+cos(2.0*M_PI*f2*ty)
                +cos(2.0*M_PI*f3*tz);
            r[i+lx*(j+ly*(k+lz*(isw0+1)))]=t;
            p[isw0+1] += (t*t);
        }
    }
}
p[isw0+1] /= (double) (nx*ny*nz);
printf( "\tTime series data\n");
for( k=0 ; k<nz ; k++ )
{
    printf( "\t r[i+%3d*(j+%3d*%3d)]\n", lx,ly,k);
    printf( "  i/j\n");
    for( j=0 ; j<ny ; j++ )
        printf( "%9d", j);
    printf( "\n" );
    printf( "  -----");
    printf( "-----\n");
    for( i=0 ; i<nx ; i++ )
    {
        printf( "%5d", i );
        for( j=0 ; j<ny ; j++ )

```

```

                printf( "%9.4lf", r[i+lx*(j+ly*(k+lz*(isw0+1)))] );
                printf( "\n" );
            }
            printf( "\n" );
        }
        printf( "\n");
        printf( "\tTime domain power =%9.4lf\n", p[isw0+1]);
        printf( "\tSignal frequency =( %9.4lf, %9.4lf, %9.4lf)\n", f1, f2, f3);
        is=0;

        for( isw=0 ; isw<=isw0 ; isw++ )
        {
            for( k=0 ; k<nz ; k++ )
                for( j=0 ; j<ny ; j++ )
                    for( i=0 ; i<nz ; i++ )
                        r[i+lx*(j+ly*(k+lz*isw))]=
                            r[i+lx*(j+ly*(k+lz*(isw0+1)))]);
            if ( isw != 0 )
                is=isw+1;

            ierr = ASL_qfps3d(nx, ny, nz, &r[lx*ly*lz*isw],
                lx, ly, lz, is, iwk, wk, nt);
            p[isw]=0.0;
            if (nx%2==0)
            {
                m=nd2-1;
                for( k=0 ; k<nz ; k++ )
                    for( j=0 ; j<ny ; j++ )
                        for( i=1 ; i<m ; i++ )
                            p[isw]+=2.0
                                *r[i+lx*(j+ly*(k+lz*isw))];
                for( k=0 ; k<nz ; k++ )
                    for( j=0 ; j<ny ; j++ )
                        p[isw]+=r[lx*(j+ly*(k+lz*isw))]
                            +r[m+lx*(j+ly*(k+lz*isw))];
            }
            else
            {
                m=nd2;
                for( k=0 ; k<nz ; k++ )
                    for( j=0 ; j<ny ; j++ )
                        for( i=1 ; i<m ; i++ )
                            p[isw]+=2.0
                                *r[i+lx*(j+ly*(k+lz*isw))];
                for( k=0 ; k<nz ; k++ )
                    for( j=0 ; j<ny ; j++ )
                        p[isw]+=r[lx*(j+ly*(k+lz*isw))];
            }
        }
    }

    printf( "\n    ** Output **\n\n" );
    printf( "\tierr = %6d\n", ierr );

    isw=0;
    printf( "\t(Modified) periodogram (Raw)\n");
    printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
    for( k=(nz+1)/2 ; k<nz ; k++ )
    {
        printf( "\tz-frq=%8.2lf\n", (k-nz)*dfz );
        printf( "  x/y-freq");
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.2lf", (j-ny)*dfy );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.2lf", j*dfy );
        printf( "\n");
        printf( "  -----");
        printf( "-----\n");
        for( i=0 ; i<nd2 ; i++ )
        {
            printf( " %8.2lf", i*dfx );
            for( j=(ny+1)/2 ; j<ny ; j++ )
                printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
            for( j=0 ; j<(ny+1)/2 ; j++ )
                printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
            printf( "\n" );
        }
        printf( "\n");
    }
    for( k=0 ; k<(nz+1)/2 ; k++ )
    {
        printf( "\tz-frq=%8.2lf\n", k*dfz );
        printf( "  x/y-freq");
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.2lf", (j-ny)*dfy );
    }

```

```

for( j=0 ; j<(ny+1)/2 ; j++ )
    printf( "%8.2lf", j*dfy );
printf( "\n");
printf( " -----");
printf( " -----\n");
for( i=0 ; i<nd2 ; i++ )
{
    printf( " %8.2lf", i*dfx );
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
    printf( "\n" );
}
printf( "\n");
}
printf( "\n");

isw=1;
printf( "\t(Modified) periodogram (Hanning)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
for( k=(nz+1)/2 ; k<nz ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", (k-nz)*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n");
    printf( " -----");
    printf( " -----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n");
}
for( k=0 ; k<(nz+1)/2 ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", k*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n");
    printf( " -----");
    printf( " -----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n");
}
printf( "\n");

isw=2;
printf( "\t(Modified) periodogram (Bartlett)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
for( k=(nz+1)/2 ; k<nz ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", (k-nz)*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n");
    printf( " -----");
    printf( " -----\n");
    for( i=0 ; i<nd2 ; i++ )
    {

```

```

        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n" );
}
for( k=0 ; k<(nz+1)/2 ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", k*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n" );
    printf( " -----");
    printf( "-----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n" );
}
printf( "\n");

isw=3;
printf( "\t(Modified) periodogram (Welch)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
for( k=(nz+1)/2 ; k<nz ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", (k-nz)*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n" );
    printf( " -----");
    printf( "-----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n" );
}
for( k=0 ; k<(nz+1)/2 ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", k*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n" );
    printf( " -----");
    printf( "-----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n" );
}
printf( "\n");

isw=4;

```

```

printf( "\t(Modified) periodogram (Parzen)\n");
printf( "\tFrequency domain power=%9.4lf\n", p[isw] );
for( k=(nz+1)/2 ; k<nz ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", (k-nz)*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n");
    printf( " -----");
    printf( "-----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n");
}
for( k=0 ; k<(nz+1)/2 ; k++ )
{
    printf( "\tz-frq=%8.2lf\n", k*dfz );
    printf( " x/y-freq");
    for( j=(ny+1)/2 ; j<ny ; j++ )
        printf( "%8.2lf", (j-ny)*dfy );
    for( j=0 ; j<(ny+1)/2 ; j++ )
        printf( "%8.2lf", j*dfy );
    printf( "\n");
    printf( " -----");
    printf( "-----\n");
    for( i=0 ; i<nd2 ; i++ )
    {
        printf( " %8.2lf", i*dfx );
        for( j=(ny+1)/2 ; j<ny ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        for( j=0 ; j<(ny+1)/2 ; j++ )
            printf( "%8.4lf", r[i+lx*(j+ly*(k+lz*isw))] );
        printf( "\n" );
    }
    printf( "\n");
}
printf( "\n");

free( iwk );
free( p );
free( wk );
free( r );

return 0;
}

```

(d) Output results

```

*** ASL_qfps3d ***

** Input **

isw=0, 2 to      5
nx=      8
ny=      8
nz=      8

Time series data
r[i+ 10*(j+ 8* 0)]
i/j      0      1      2      3      4      5      6      7
-----
0  3.0000  2.9048  2.6374  2.2487  1.8126  1.4122  1.1237  1.0020
1  1.6319  1.5367  1.2693  0.8806  0.4445  0.0441 -0.2444 -0.3662
2  1.2710  1.1759  0.9085  0.5197  0.0837 -0.3168 -0.6053 -0.7270
3  2.9048  2.8097  2.5423  2.1535  1.7174  1.3170  1.0285  0.9068
4  2.0628  1.9676  1.7002  1.3115  0.8754  0.4750  0.1865  0.0648
5  1.0489  0.9538  0.6864  0.2976 -0.1384 -0.5388 -0.8274 -0.9491
6  2.6374  2.5423  2.2748  1.8861  1.4500  1.0496  0.7611  0.6394
7  2.4818  2.3866  2.1192  1.7304  1.2944  0.8940  0.6054  0.4837

r[i+ 10*(j+ 8* 1)]
i/j      0      1      2      3      4      5      6      7
-----
0  1.8436  1.7484  1.4810  1.0923  0.6562  0.2558 -0.0327 -0.1545
1  0.4754  0.3803  0.1129 -0.2759 -0.7119 -1.1123 -1.4009 -1.5226
2  0.1146  0.0194 -0.2480 -0.6367 -1.0728 -1.4732 -1.7617 -1.8834
3  1.7484  1.6532  1.3858  0.9971  0.5610  0.1606 -0.1279 -0.2496

```

```

4  0.9064  0.8112  0.5438  0.1550 -0.2810 -0.6814 -0.9700 -1.0917
5 -0.1075 -0.2027 -0.4701 -0.8588 -1.2949 -1.6953 -1.9838 -2.1055
6  1.4810  1.3858  1.1184  0.7297  0.2936 -0.1068 -0.3953 -0.5170
7  1.3253  1.2301  0.9627  0.5740  0.1379 -0.2625 -0.5510 -0.6727

  r[i+ 10*(j+ 8* 2)]
i/j  0      1      2      3      4      5      6      7
-----
0  1.0489  0.9538  0.6864  0.2976 -0.1384 -0.5388 -0.8274 -0.9491
1 -0.3192 -0.4144 -0.6818 -1.0705 -1.5066 -1.9070 -2.1955 -2.3172
2 -0.6800 -0.7752 -1.0426 -1.4313 -1.8674 -2.2678 -2.5563 -2.6781
3  0.9538  0.8586  0.5912  0.2025 -0.2336 -0.6340 -0.9225 -1.0443
4  0.1117  0.0166 -0.2508 -0.6396 -1.0756 -1.4761 -1.7646 -1.8863
5 -0.9021 -0.9973 -1.2647 -1.6534 -2.0895 -2.4899 -2.7784 -2.9001
6  0.6864  0.5912  0.3238 -0.0649 -0.5010 -0.9014 -1.1899 -1.3117
7  0.5307  0.4355  0.1681 -0.2206 -0.6567 -1.0571 -1.3456 -1.4673

  r[i+ 10*(j+ 8* 3)]
i/j  0      1      2      3      4      5      6      7
-----
0  2.4540  2.3588  2.0914  1.7027  1.2666  0.8662  0.5777  0.4560
1  1.0859  0.9907  0.7233  0.3346 -0.1015 -0.5019 -0.7904 -0.9122
2  0.7250  0.6298  0.3624 -0.0263 -0.4624 -0.8628 -1.1513 -1.2730
3  2.3588  2.2636  1.9962  1.6075  1.1714  0.7710  0.4825  0.3608
4  1.5168  1.4216  1.1542  0.7655  0.3294 -0.0710 -0.3595 -0.4812
5  0.5029  0.4078  0.1404 -0.2484 -0.6844 -1.0849 -1.3734 -1.4951
6  2.0914  1.9962  1.7288  1.3401  0.9040  0.5036  0.2151  0.0934
7  1.9357  1.8406  1.5732  1.1844  0.7484  0.3480  0.0594 -0.0623

  r[i+ 10*(j+ 8* 4)]
i/j  0      1      2      3      4      5      6      7
-----
0  2.8090  2.7138  2.4464  2.0577  1.6216  1.2212  0.9327  0.8110
1  1.4409  1.3457  1.0783  0.6896  0.2535 -0.1469 -0.4354 -0.5571
2  1.0800  0.9849  0.7175  0.3287 -0.1073 -0.5077 -0.7963 -0.9180
3  2.7138  2.6187  2.3513  1.9625  1.5265  1.1261  0.8375  0.7158
4  1.8718  1.7766  1.5092  1.1205  0.6844  0.2840 -0.0045 -0.1262
5  0.8580  0.7628  0.4954  0.1067 -0.3294 -0.7298 -1.0183 -1.1401
6  2.4464  2.3513  2.0839  1.6951  1.2591  0.8587  0.5701  0.4484
7  2.2908  2.1956  1.9282  1.5395  1.1034  0.7030  0.4145  0.2927

  r[i+ 10*(j+ 8* 5)]
i/j  0      1      2      3      4      5      6      7
-----
0  1.2929  1.1977  0.9303  0.5416  0.1055 -0.2949 -0.5834 -0.7051
1 -0.0752 -0.1704 -0.4378 -0.8265 -1.2626 -1.6630 -1.9515 -2.0733
2 -0.4361 -0.5312 -0.7987 -1.1874 -1.6235 -2.0239 -2.3124 -2.4341
3  1.1977  1.1025  0.8351  0.4464  0.0103 -0.3901 -0.6786 -0.8003
4  0.3557  0.2605 -0.0069 -0.3956 -0.8317 -1.2321 -1.5206 -1.6423
5 -0.6582 -0.7533 -1.0207 -1.4095 -1.8455 -2.2459 -2.5345 -2.6562
6  0.9303  0.8351  0.5677  0.1790 -0.2571 -0.6575 -0.9460 -1.0677
7  0.7746  0.6795  0.4121  0.0233 -0.4127 -0.8131 -1.1017 -1.2234

  r[i+ 10*(j+ 8* 6)]
i/j  0      1      2      3      4      5      6      7
-----
0  1.4122  1.3170  1.0496  0.6609  0.2248 -0.1756 -0.4641 -0.5858
1  0.0441 -0.0511 -0.3185 -0.7072 -1.1433 -1.5437 -1.8322 -1.9539
2 -0.3168 -0.4119 -0.6793 -1.0681 -1.5041 -1.9045 -2.1931 -2.3148
3  1.3170  1.2219  0.9545  0.5657  0.1297 -0.2707 -0.5593 -0.6810
4  0.4750  0.3798  0.1124 -0.2763 -0.7124 -1.1128 -1.4013 -1.5230
5 -0.5388 -0.6340 -0.9014 -1.2902 -1.7262 -2.1266 -2.4151 -2.5369
6  1.0496  0.9545  0.6871  0.2983 -0.1377 -0.5381 -0.8267 -0.9484
7  0.8940  0.7988  0.5314  0.1427 -0.2934 -0.6938 -0.9823 -1.1041

  r[i+ 10*(j+ 8* 7)]
i/j  0      1      2      3      4      5      6      7
-----
0  2.8910  2.7958  2.5284  2.1397  1.7036  1.3032  1.0147  0.8930
1  1.5229  1.4277  1.1603  0.7716  0.3355 -0.0649 -0.3534 -0.4751
2  1.1620  1.0669  0.7995  0.4107 -0.0253 -0.4257 -0.7143 -0.8360
3  2.7958  2.7007  2.4333  2.0445  1.6085  1.2080  0.9195  0.7978
4  1.9538  1.8586  1.5912  1.2025  0.7664  0.3660  0.0775 -0.0442
5  0.9400  0.8448  0.5774  0.1886 -0.2474 -0.6478 -0.9364 -1.0581
6  2.5284  2.4333  2.1659  1.7771  1.3410  0.9406  0.6521  0.5304
7  2.3728  2.2776  2.0102  1.6215  1.1854  0.7850  0.4965  0.3747

Time domain power = 1.6439
Signal frequency =( 0.6200, 0.1400, 0.5500)

** Output **

ierr = 0
(Modified) periodogram (Raw)
Frequency domain power= 1.5531
z-frq= -1.00
x/y-freq -1.00 -0.75 -0.50 -0.25 0.00 0.25 0.50 0.75
-----
0.00 0.0000 0.0000 0.0000 0.0000 0.0007 0.0000 0.0000 0.0000
0.25 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000
0.50 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000
0.75 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000 0.0000
    
```



z-freq=	-0.75								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.0071	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	-0.50								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.2513	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	-0.25								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.0136	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	0.00								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0158	0.0188	0.0350	0.2150	0.0583	0.2150	0.0350	0.0188	
0.25	0.0000	0.0000	0.0000	0.0000	0.0239	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.1352	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0781	0.0000	0.0000	0.0000	
z-freq=	0.25								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.0136	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	0.50								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.2513	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	0.75								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.0071	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
(Modified) periodogram (Hanning)									
Frequency domain power= 1.0699									
z-freq=	-1.00								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0001	0.0004	0.0001	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0001	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	-0.75								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0078	0.0310	0.0078	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0019	0.0078	0.0019	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	-0.50								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0176	0.0704	0.0176	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0044	0.0176	0.0044	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq=	-0.25								
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0009	0.0164	0.0045	0.0053	0.0009	0.0000	
0.25	0.0000	0.0000	0.0002	0.0027	0.0004	0.0023	0.0002	0.0000	
0.50	0.0000	0.0000	0.0000	0.0034	0.0136	0.0034	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0034	0.0137	0.0034	0.0000	0.0000	

z-freq= 0.00									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0001	0.0036	0.0427	0.0087	0.0427	0.0036	0.0001	
0.25	0.0000	0.0000	0.0009	0.0064	0.0041	0.0158	0.0009	0.0000	
0.50	0.0000	0.0000	0.0000	0.0034	0.0136	0.0543	0.0136	0.0000	
0.75	0.0000	0.0000	0.0000	0.0137	0.0547	0.0137	0.0000	0.0000	
z-freq= 0.25									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0009	0.0053	0.0045	0.0164	0.0009	0.0000	
0.25	0.0000	0.0000	0.0002	0.0006	0.0031	0.0058	0.0002	0.0000	
0.50	0.0000	0.0000	0.0000	0.0034	0.0136	0.0034	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0034	0.0137	0.0034	0.0000	0.0000	
z-freq= 0.50									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0176	0.0704	0.0176	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0044	0.0176	0.0044	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq= 0.75									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0078	0.0310	0.0078	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0019	0.0078	0.0019	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
(Modified) periodogram (Bartlett)									
Frequency domain power= 1.0593									
z-freq= -1.00									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0000	0.0000	0.0001	0.0000	0.0000	0.0000	
0.25	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
z-freq= -0.75									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0002	0.0000	0.0079	0.0346	0.0050	0.0000	0.0002	
0.25	0.0000	0.0000	0.0000	0.0014	0.0062	0.0009	0.0000	0.0000	
0.50	0.0000	0.0000	0.0000	0.0001	0.0003	0.0001	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0003	0.0001	0.0000	0.0000	
z-freq= -0.50									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0005	0.0000	0.0165	0.0907	0.0165	0.0000	0.0005	
0.25	0.0000	0.0001	0.0000	0.0030	0.0165	0.0030	0.0000	0.0001	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0001	0.0005	0.0001	0.0000	0.0000	
z-freq= -0.25									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0001	0.0141	0.0037	0.0074	0.0001	0.0000	
0.25	0.0000	0.0000	0.0000	0.0022	0.0005	0.0017	0.0000	0.0000	
0.50	0.0000	0.0001	0.0000	0.0021	0.0113	0.0021	0.0000	0.0001	
0.75	0.0000	0.0001	0.0000	0.0012	0.0096	0.0024	0.0000	0.0001	
z-freq= 0.00									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0007	0.0594	0.0070	0.0594	0.0007	0.0000	
0.25	0.0000	0.0000	0.0001	0.0089	0.0017	0.0129	0.0001	0.0000	
0.50	0.0000	0.0003	0.0000	0.0113	0.0622	0.0113	0.0000	0.0003	
0.75	0.0000	0.0003	0.0000	0.0069	0.0554	0.0139	0.0000	0.0003	
z-freq= 0.25									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0001	0.0074	0.0037	0.0141	0.0001	0.0000	
0.25	0.0000	0.0000	0.0000	0.0011	0.0011	0.0030	0.0000	0.0000	
0.50	0.0000	0.0001	0.0000	0.0021	0.0113	0.0021	0.0000	0.0001	
0.75	0.0000	0.0001	0.0000	0.0014	0.0108	0.0027	0.0000	0.0001	
z-freq= 0.50									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0005	0.0000	0.0165	0.0907	0.0165	0.0000	0.0005	
0.25	0.0000	0.0001	0.0000	0.0030	0.0165	0.0030	0.0000	0.0001	
0.50	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0001	0.0005	0.0001	0.0000	0.0000	
z-freq= 0.75									

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0002	0.0000	0.0050	0.0346	0.0079	0.0000	0.0002
0.25	0.0000	0.0000	0.0000	0.0009	0.0065	0.0015	0.0000	0.0000
0.50	0.0000	0.0000	0.0000	0.0001	0.0003	0.0001	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0001	0.0008	0.0002	0.0000	0.0000

(Modified) periodogram (Welch)  
 Frequency domain power= 1.2154  
 z-freq= -1.00

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0000	0.0005	0.0030	0.0005	0.0000	0.0000
0.25	0.0000	0.0000	0.0000	0.0001	0.0003	0.0001	0.0000	0.0000
0.50	0.0000	0.0000	0.0000	0.0000	0.0002	0.0000	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0002	0.0000	0.0000	0.0000

z-freq= -0.75

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0001	0.0001	0.0004	0.0051	0.0357	0.0028	0.0003	0.0001
0.25	0.0000	0.0000	0.0000	0.0005	0.0038	0.0003	0.0000	0.0000
0.50	0.0000	0.0000	0.0000	0.0001	0.0009	0.0001	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0002	0.0000	0.0000	0.0000

z-freq= -0.50

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0003	0.0004	0.0011	0.0103	0.1247	0.0186	0.0011	0.0004
0.25	0.0000	0.0000	0.0001	0.0011	0.0131	0.0020	0.0001	0.0000
0.50	0.0000	0.0000	0.0000	0.0001	0.0009	0.0001	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0001	0.0017	0.0002	0.0000	0.0000

z-freq= -0.25

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0001	0.0122	0.0012	0.0090	0.0000	0.0000
0.25	0.0000	0.0000	0.0000	0.0013	0.0002	0.0009	0.0000	0.0000
0.50	0.0000	0.0000	0.0001	0.0017	0.0095	0.0005	0.0001	0.0000
0.75	0.0000	0.0000	0.0001	0.0005	0.0078	0.0012	0.0001	0.0000

z-freq= 0.00

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0000	0.1034	0.0186	0.1034	0.0000	0.0000
0.25	0.0000	0.0000	0.0000	0.0115	0.0021	0.0104	0.0000	0.0000
0.50	0.0002	0.0003	0.0008	0.0158	0.0862	0.0044	0.0007	0.0003
0.75	0.0002	0.0002	0.0007	0.0052	0.0760	0.0115	0.0007	0.0002

z-freq= 0.25

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0000	0.0090	0.0012	0.0122	0.0001	0.0000
0.25	0.0000	0.0000	0.0000	0.0010	0.0001	0.0012	0.0000	0.0000
0.50	0.0000	0.0000	0.0001	0.0016	0.0086	0.0004	0.0001	0.0000
0.75	0.0000	0.0000	0.0000	0.0006	0.0083	0.0012	0.0001	0.0000

z-freq= 0.50

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0003	0.0004	0.0011	0.0186	0.1247	0.0103	0.0011	0.0004
0.25	0.0000	0.0000	0.0001	0.0020	0.0133	0.0011	0.0001	0.0000
0.50	0.0000	0.0000	0.0000	0.0004	0.0031	0.0003	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0004	0.0001	0.0000	0.0000

z-freq= 0.75

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0001	0.0001	0.0003	0.0028	0.0357	0.0051	0.0004	0.0001
0.25	0.0000	0.0000	0.0000	0.0003	0.0037	0.0005	0.0000	0.0000
0.50	0.0000	0.0000	0.0000	0.0000	0.0003	0.0000	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0005	0.0001	0.0000	0.0000

(Modified) periodogram (Parzen)  
 Frequency domain power= 0.9132  
 z-freq= -1.00

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0001	0.0023	0.0053	0.0023	0.0001	0.0000
0.25	0.0000	0.0000	0.0001	0.0010	0.0023	0.0010	0.0001	0.0000
0.50	0.0000	0.0000	0.0000	0.0001	0.0001	0.0001	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000

z-freq= -0.75

x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75
0.00	0.0000	0.0000	0.0006	0.0092	0.0209	0.0089	0.0006	0.0000
0.25	0.0000	0.0000	0.0003	0.0039	0.0090	0.0038	0.0002	0.0000
0.50	0.0000	0.0000	0.0000	0.0002	0.0005	0.0002	0.0000	0.0000
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000

z-freq= -0.50									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0014	0.0162	0.0313	0.0112	0.0005	0.0000	
0.25	0.0000	0.0000	0.0005	0.0062	0.0120	0.0044	0.0002	0.0000	
0.50	0.0000	0.0000	0.0000	0.0003	0.0007	0.0004	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0003	0.0007	0.0003	0.0000	0.0000	
z-freq= -0.25									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0001	0.0030	0.0125	0.0054	0.0016	0.0012	0.0000	
0.25	0.0000	0.0000	0.0008	0.0019	0.0010	0.0030	0.0010	0.0000	
0.50	0.0000	0.0000	0.0001	0.0032	0.0106	0.0063	0.0006	0.0000	
0.75	0.0000	0.0000	0.0003	0.0046	0.0108	0.0048	0.0003	0.0000	
z-freq= 0.00									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0001	0.0047	0.0117	0.0007	0.0117	0.0047	0.0001	
0.25	0.0000	0.0000	0.0011	0.0006	0.0055	0.0140	0.0033	0.0001	
0.50	0.0000	0.0000	0.0003	0.0089	0.0291	0.0168	0.0016	0.0000	
0.75	0.0000	0.0000	0.0007	0.0107	0.0253	0.0111	0.0007	0.0000	
z-freq= 0.25									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0012	0.0016	0.0054	0.0125	0.0030	0.0001	
0.25	0.0000	0.0000	0.0002	0.0005	0.0086	0.0110	0.0019	0.0000	
0.50	0.0000	0.0000	0.0002	0.0048	0.0151	0.0085	0.0008	0.0000	
0.75	0.0000	0.0000	0.0003	0.0047	0.0110	0.0049	0.0003	0.0000	
z-freq= 0.50									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0005	0.0112	0.0313	0.0162	0.0014	0.0000	
0.25	0.0000	0.0000	0.0003	0.0055	0.0155	0.0081	0.0007	0.0000	
0.50	0.0000	0.0000	0.0001	0.0010	0.0028	0.0014	0.0001	0.0000	
0.75	0.0000	0.0000	0.0000	0.0003	0.0008	0.0003	0.0000	0.0000	
z-freq= 0.75									
x/y-freq	-1.00	-0.75	-0.50	-0.25	0.00	0.25	0.50	0.75	
0.00	0.0000	0.0000	0.0006	0.0089	0.0209	0.0092	0.0006	0.0000	
0.25	0.0000	0.0000	0.0002	0.0039	0.0091	0.0040	0.0003	0.0000	
0.50	0.0000	0.0000	0.0000	0.0003	0.0006	0.0003	0.0000	0.0000	
0.75	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	0.0000	

## Chapter 7

---

# SORTING

## 7.1 INTRODUCTION

This chapter describes the functions for sorting data.

**Function described in this chapter divides up and allocates internal processing among threads and executes allocated processing in parallel.**

This library provides functions having the following operations.

- (1) Sorting a list of data
- (2) Sorting a list of pairwise data

### 7.1.1 Notes

Since the parallel processing overhead significantly affects the computation cost if the number of data to be sorted is small, performance may be worse than when a non-parallel function is used.

## 7.1.2 Algorithms Used

The algorithms for sorting in ascending order are shown below. The algorithms for sorting in descending order, which differ only in terms of the relative magnitudes, are similar.

(1) Shell sort

- (1) Set the spacing  $h$ .
- (2) Take all subsequences of spacing  $h$  from the data sequence.
- (3) Compare adjacent pairs within each subsequence to arrange them in ascending order. If they are in the reverse order, exchange their positions and confirm again the relative order with the preceding data. If they are again in the reverse order, exchange the positions and work back toward the beginning.
- (4) Decrease the spacing  $h$  and repeat steps (2) and (3). When the processing for  $h = 1$  ends, sorting is completed.

(2) Heap sort

- (1) Organize the assigned data into a heap tree (well-ordered binary tree in which parents have value greater than or equal to those of children).
- (2) Exchange the root and the data at the very end of the tree.
- (3) Let the position excluding the very last data be  $A$ .
- (4) Consider  $A$  to be a new tree, and organize this into a heap tree again.
- (5) Repeat steps (2) to (4). When the remaining data is only the root, sorting is completed.

(3) Quick sort

- (1) Count the number of data values within the sort interval.
- (2) Do the following depending on the number of data values.
  - If the number of data values is less than or equal to one:  
Do nothing.
  - If the number of data values is 2:  
If they are in the reverse order, exchange their positions.
  - If the number of data values is greater than or equal to three:
    - ① Select one pivot value from within the interval.
    - ② Divide the data within the interval into two intervals consisting of values greater than the pivot value and values less than the pivot value.
- (3) Repeat steps (1) and (2). When the number of data values in all data intervals is less than or equal to two, sorting is completed.

(4) Merge sort

- (1) Count the number of data values within the sort interval.
- (2) Do the following depending on the number of data values.
  - If the number of data value is one:  
Do nothing.
  - If the number of data values is two:  
If they are in the reverse order, exchange the positions.
  - If the number of data values is greater than or equal to three:
    - ① Divide the data within the interval in half into the top half and bottom half.

- ② Recursively merge sort the top half. Recursively merge sort the bottom half.
- ③ Merge the sorted top half and bottom half.

### **7.1.3 Reference Bibliography**

- (1) Niklaus Wirth, “ALGORITHMS + DATA STRUCTURES = PROGRAMS”, Prentice–Hall Inc. (1976).
- (2) Hiroto Namihira, “Sorting and Searching”, CQ Publishing Co. Ltd.
- (3) Yoshiyuki Kondo, “Algorithms and Data Structures”, Softbank Publishing Inc.



## 7.2 SORTING

### 7.2.1 ASL\_qssta1, ASL\_pssta1 Sorting a List of Data

(1) **Function**

Given  $n$  data values  $a_{i_k}$  ( $k = 1, 2, \dots, n$ ), the ASL\_qssta1 or ASL\_pssta1 obtains the data sequence  $a_{j_k}$  ( $k = 1, 2, \dots, n$ ) consisting of the original data values  $a_i$  rearranged in ascending or descending order. Here,  $a_j$  satisfies the following relationship.

For ascending order :  $a_{j_1} \leq a_{j_2} \leq \dots \leq a_{j_n}$

For descending order :  $a_{j_1} \geq a_{j_2} \geq \dots \geq a_{j_n}$

(2) **Usage**

Double precision:

ierr = ASL\_qssta1 (a,n,isw,wk,iwk,nt);

Single precision:

ierr = ASL\_pssta1 (a,n,isw,wk,iwk,nt);

(3) **Arguments and Return Value**

D:Double precision real    Z:Double precision complex    I:  $\begin{cases} \text{int} & \text{as for 32bit Integer} \\ \text{long} & \text{as for 64bit Integer} \end{cases}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\begin{cases} D* \\ R* \end{cases}$	n	Input	Data to be sorted $a_i$
				Output	Sorted data $a_j$
2	n	I	1	Input	Size of array a
3	isw	I	1	Input	Sort method selection switch (See Note (a))
4	wk	$\begin{cases} D* \\ R* \end{cases}$	n	Work	Work area
5	iwk	I*	2×n	Work	Work area
6	nt	I	1	Input	Number of tasks to be generated
7	ierr	I	1	Output	Error indicator (Return Value)

(4) **Restrictions**

(a)  $n \geq 1$

(b)  $isw \in \{\pm 1, \pm 2, \pm 3, \pm 4\}$

(c)  $nt \geq 1$

(5) **Error indicator (Return Value)**

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	

(6) **Notes**

(a) The sort methods to be selected by isw are as follows.

isw	Sort Method	isw	Sort Method
1	Shell sort (ascending order)	-4	Shell sort (descending order)
2	Heap sort (ascending order)	-2	Heap sort (descending order)
3	Quick sort (ascending order)	-3	Quick sort (descending order)
4	Merge sort (ascending order)	-4	Merge sort (descending order)

The user should select an appropriate sort method according to the properties of the input data. The features of each sort method are shown below.

· Shell sort

The average of the amount of calculation is on the order of  $O(n^{1.5})$ . A fast, stable sort is performed for any kind of data. Sorting is faster if part of the data sequence has been sorted.

Retention of the ordinal relationships among data having the same value is not guaranteed between before and after sorting.

No work area is necessary.

· Heap sort

Although the amount of calculation is on the order of  $O(n \log n)$ , the constant term portion is large. The sort time does not change much according to the properties of the input data.

Retention of the ordinal relationships among data having the same value is not guaranteed between before and after sorting.

No work area is necessary.

· Quick sort

Although the average of the amount of calculation is on the order of  $O(n \log n)$ , this is an extremely inefficient sort for data having certain types of regularities such as data that has been partially sorted to begin with. This is the fastest sort method for random data.

Retention of the ordinal relationships among data having the same value is not guaranteed between before and after sorting.

· Merge sort

Although the amount of calculation is on the order of  $O(n \log n)$ , the constant term portion is large. The sort time does not change much according to the properties of the input data.

The ordinal relationships before sorting among data having the same value is retained after sorting.

(7) **Example**

(a) Problem

Sort the following data in ascending order by using a shell sort.

a[0] = 5.0

a[1] = 4.0

a[2] = 9.0  
a[3] = 6.0  
a[4] = 2.0  
a[5] = 5.0

(b) Input data

Array a, n=6 and isw=1.

(c) Main program

```
/*      C interface example for ASL_qssta1 */
#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a,*wk;
    int n,isw,*iwk,ierr;
    int nt;
    int i;
    FILE *fp;

    fp = fopen( "qssta1.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    isw=1;
    n=6;
    nt=4;

    printf( "      *** ASL_qssta1 ***\n" );
    printf( "\n      ** Input **\n\n" );
    printf( "\tn=%3d\n\n", n );

    a = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    wk = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    iwkw = ( int * )malloc((size_t)( sizeof(int) * (2*n) ));
    if( iwkw == NULL )
    {
        printf( "no enough memory for array iwkw\n" );
        return -1;
    }

    printf( "\tArray a" );
    for( i=0 ; i<n ; i++ )
    {
        fscanf( fp, "%lf", &a[i] );
        printf( "%8.3g", a[i] );
    }
    printf( "\n" );
    fclose( fp );

    ierr = ASL_qssta1(a, n, isw, wk, iwkw, nt);

    printf( "\n      ** Output **\n\n" );
    printf( "\tierr = %6d\n\n", ierr );

    printf( "\tArray a" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "%8.3g", a[i] );
    }
    printf( "\n" );

    free( a );
}
```

```
    free( wk );  
    free( iwk );  
    return 0;  
}
```

(d) Output results

```
*** ASL_qssta1 ***  
** Input **  
n= 6  
Array a      5      4      9      6      2      5  
** Output **  
ierr =      0  
Array a      2      4      5      5      6      9
```

## 7.2.2 ASL\_qssta2, ASL\_pssta2 Sorting a List of Pairwise Data

### (1) Function

Given two sets of  $n$  data values  $a_{i_k} (k = 1, 2, \dots, n)$ ,  $b_{i_k} (k = 1, 2, \dots, n)$ , the ASL\_qssta2 or ASL\_pssta2 obtains the data sequence  $a_{j_k} (k = 1, 2, \dots, n)$  consisting of the original  $a_i$  data values rearranged in ascending or ascending order and the data sequence  $b_{j_k} (k = 1, 2, \dots, n)$  corresponding to it.

Here,  $a_j$  satisfies the following relationship.

For ascending order :  $a_{j_1} \leq a_{j_2} \leq \dots \leq a_{j_n}$

For descending order :  $a_{j_1} \geq a_{j_2} \geq \dots \geq a_{j_n}$

If a second order sort is specified, the function determines  $j = j_1, j_2, \dots, j_n$  so that the following relationship is satisfied for any  $k$  for which  $a_{j_k} = a_{j_{k+1}}$  is satisfied.

For ascending order :  $b_{j_k} \leq b_{j_{k+1}}$

For descending order :  $b_{j_k} \geq b_{j_{k+1}}$

### (2) Usage

Double precision:

```
ierr = ASL_qssta2 (a,n,b,isw1,isw2,wk,iwk,nt);
```

Single precision:

```
ierr = ASL_pssta2 (a,n,b,isw1,isw2,wk,iwk,nt);
```

(3) Arguments and Return Value

D:Double precision real    Z:Double precision complex    I:  $\left\{ \begin{array}{l} \text{int as for 32bit Integer} \\ \text{long as for 64bit Integer} \end{array} \right\}$   
R:Single precision real    C:Single precision complex

No.	Argument and Return Value	Type	Size	Input/Output	Contents
1	a	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	n	Input	Data to be sorted $a_i$
				Output	Sorted data $a_j$
2	n	I	1	Input	Size of array a
3	b	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	n	Input	Data $b_i$ corresponding to $a_i$
				Output	Data $b_j$ corresponding to sorted $a_j$
4	isw1	I	1	Input	Sort method selection switch (See Note (a))
5	isw2	I	1	Input	Second order sort switch isw2=0 : Do not perform second order sort isw2=1 : Perform second order sort
6	wk	$\left\{ \begin{array}{l} D^* \\ R^* \end{array} \right\}$	2×n	Work	Work area
7	iwk	I*	2×n	Work	Work area
8	nt	I	1	Input	Number of tasks to be generated
9	ierr	I	1	Output	Error indicator (Return Value)

(4) Restrictions

- (a)  $n \geq 1$
- (b)  $isw1 \in \{\pm 1, \pm 2, \pm 3, \pm 4\}$
- (c)  $isw2 \in \{0, 1\}$
- (d)  $nt \geq 1$

(5) Error indicator (Return Value)

ierr value	Meaning	Processing
0	Normal termination.	
3000	Restriction (a) was not satisfied.	Processing is aborted.
3100	Restriction (b) was not satisfied.	
3200	Restriction (c) was not satisfied.	
3300	Restriction (d) was not satisfied.	

(6) **Notes**

(a) The sort methods to be selected by isw1 are as follows.

isw1	Sort Method	isw1	Sort Method
1	Shell sort (ascending order)	-1	Shell sort (descending order)
2	Heap sort (ascending order)	-2	Heap sort (descending order)
3	Quick sort (ascending order)	-3	Quick sort (descending order)
4	Merge sort (ascending order)	-4	Merge sort (descending order)

The user should select an appropriate sort method according to the properties of the input data. The features of each sort method are shown below.

· Shell sort

The average of the amount of calculation is on the order of  $O(n^{1.5})$ . A fast, stable sort is performed for any kind of data. Sorting is faster if part of the data sequence has been sorted.

It is not guaranteed that ordinal relations among plural values of the second set having the same value of the first set keep unchanged between before and after sorting.

No work area is necessary.

· Heap sort

Although the amount of calculation is on the order of  $O(n \log n)$ , the constant term portion is large. The sort time does not change much according to the properties of the input data.

It is not guaranteed that ordinal relations among plural values of the second set having the same value of the first set keep unchanged between before and after sorting.

No work area is necessary.

· Quick sort

Although the average of the amount of calculation is on the order of  $O(n \log n)$ , this is an extremely inefficient sort for data having certain types of regularities such as data that has been partially sorted to begin with. This is the fastest sort method for random data.

It is not guaranteed that ordinal relations among plural values of the second set having the same value of the first set keep unchanged between before and after sorting.

· Merge sort

Although the amount of calculation is on the order of  $O(n \log n)$ , the constant term portion is large. The sort time does not change much according to the properties of the input data.

The ordinal relationships before sorting among data having the same value is retained after sorting.

(7) **Example**

(a) Problem

Sort the following data for a in ascending order by using a shell sort and rearrange the corresponding data for b according to the sorted data for a. Also perform a second order sort.

a[0] = 5.0, b[0] = 3.0

a[1] = 4.0, b[1] = 4.0

a[2] = 9.0, b[2] = 2.0

a[3] = 6.0, b[3] = 3.0

a[4] = 2.0, b[4] = 8.0

a[5] = 5.0, b[5] = 1.0

(b) Input data

Arrays a and b, n=6, isw1=1 and isw2=1.

(c) Main program

```

/*      C interface example for ASL_qssta2 */

#include <stdio.h>
#include <stdlib.h>
#include <asl.h>

int main()
{
    double *a,*b,*wk;
    int n,isw1,isw2,*iwk,ierr;
    int nt;
    int i;
    FILE *fp;

    fp = fopen( "qssta2.dat", "r" );
    if( fp == NULL )
    {
        printf( "file open error\n" );
        return -1;
    }

    isw1=1;
    isw2=1;
    n=6;
    nt=4;

    printf( "      *** ASL_qssta2 ***\n" );
    printf( "\n      ** Input **\n\n" );
    printf( "\tn=%3d\n", n );

    a = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( a == NULL )
    {
        printf( "no enough memory for array a\n" );
        return -1;
    }

    b = ( double * )malloc((size_t)( sizeof(double) * n ));
    if( b == NULL )
    {
        printf( "no enough memory for array b\n" );
        return -1;
    }

    wk = ( double * )malloc((size_t)( sizeof(double) * (2*n) ));
    if( wk == NULL )
    {
        printf( "no enough memory for array wk\n" );
        return -1;
    }

    iw = ( int * )malloc((size_t)( sizeof(int) * (2*n) ));
    if( iw == NULL )
    {
        printf( "no enough memory for array iw\n" );
        return -1;
    }

    for( i=0 ; i<n ; i++ )
    {
        fscanf( fp, "%lf %lf", &a[i], &b[i] );
    }
    printf( "\t <Array a>      <Array b>\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t%8.3g      %8.3g\n", a[i], b[i] );
    }

    fclose( fp );

    ierr = ASL_qssta2(a, n, b, isw1, isw2, wk, iw, nt);

    printf( "\n      ** Output **\n\n" );
    printf( "\t ierr = %6d\n", ierr );

    printf( "\t <Array a>      <Array b>\n" );
    for( i=0 ; i<n ; i++ )
    {
        printf( "\t%8.3g      %8.3g\n", a[i], b[i] );
    }

    free( a );
    free( b );
    free( wk );
    free( iw );

    return 0;
}

```



(d) Output results

```
*** ASL_qssta2 ***  
** Input **  
n= 6  
  <Array a>  <Array b>  
    5         3  
    4         4  
    9         2  
    6         3  
    2         8  
    5         1  
  
** Output **  
ierr = 0  
  <Array a>  <Array b>  
    2         8  
    4         4  
    5         1  
    5         3  
    6         3  
    9         2
```

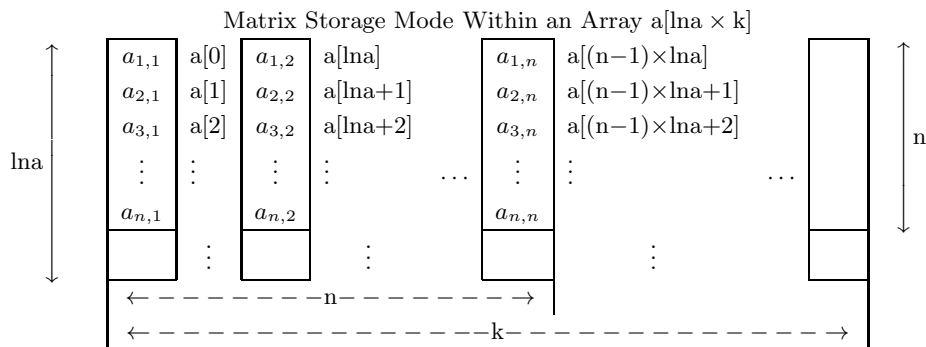
# Appendix A

## METHODS OF HANDLING ARRAY DATA

### A.1 Methods of handling array data corresponding to matrix

Since the ASL C interface function library uses array data corresponding to matrix, this section describes various methods of handling arrays.

To call a function that uses array data, you must declare that array in advance in the calling program. If the declared array is  $a[\text{lna} \times k]$ , then  $n \times n$  matrix  $A = (a_{i,j})$  ( $i = 1, 2, \dots, n; j = 1, 2, \dots, n$ ) is stored in array  $a$  as shown in the figure below.

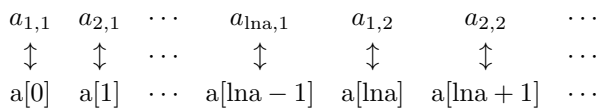


#### Remarks

- $\text{lna} \geq n$  and  $k \geq n$  must hold.
- Matrix element  $a_{i,j}$  corresponds to the array element  $a[(i-1) + \text{lna} \times (j-1)]$ .

Figure A-1 Matrix Storage Mode Within an Array  $a[\text{lna} \times k]$

$\text{lna}$  is called an adjustable dimension. If a two-dimensional array is used as an argument, the adjustable must be passed to the function as an argument in addition to the array name and order of the array. Because the matrix storage mode in ASL C interface corresponds to that of FORTRAN and the matrix elements  $a_{i,j}$  ( $i = 1, 2, \dots, \text{lna}; j = 1, 2, \dots, k$ ) must correspond to the array element  $a[\ell]$  ( $\ell = 0, 1, 2, \dots, \text{lna} \times k - 1$ ), as follows on the main memory.



#### Example ASL\_dam1ad (Real matrix addition)

Add  $3 \times 2$  matrices  $A$  and  $B$  placing the sum in matrix  $C$ . If you declare arrays of size  $[5 \times 4]$ , the declaration is as follows.

```

/*      C interface example for ASL_dam1ad */
#include <stdio.h>
#include <stdlib.h>

#include <asl.h>

int main()
{
    double *a, *b, *c;
    int lma, lmb, lmc;
    int m, n, ierr;

```

```

int k;
lma = lmb = lmc = 5;
k = 4;
m = 3;
n = 2;
a = (double *)malloc((size_t) sizeof(double) * lma*k);
if(a == NULL)
{
    printf("no enough memory for array a\n");
    return -1;
}
b = (double *)malloc((size_t) sizeof(double) * lmb*k);
if(b == NULL)
{
    printf("no enough memory for array b\n");
    return -1;
}
c = (double *)malloc((size_t) sizeof(double) * lmc*k);
if(c == NULL)
{
    printf("no enough memory for array c\n");
    return -1;
}

    :
ierr = ASL_dam1ad(a, lma , m, n, b, lmb, c, lmc);
    :

free(a);
free(b);
free(c);
return 0;
}

```

Data is stored in a as follows. Data are stored in b and c in the same way.

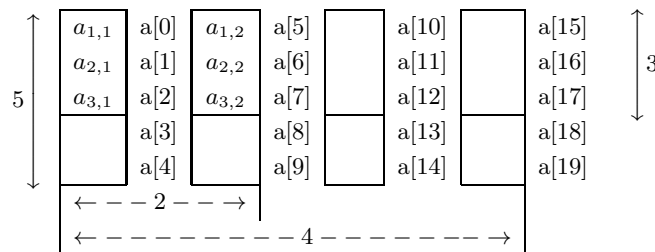


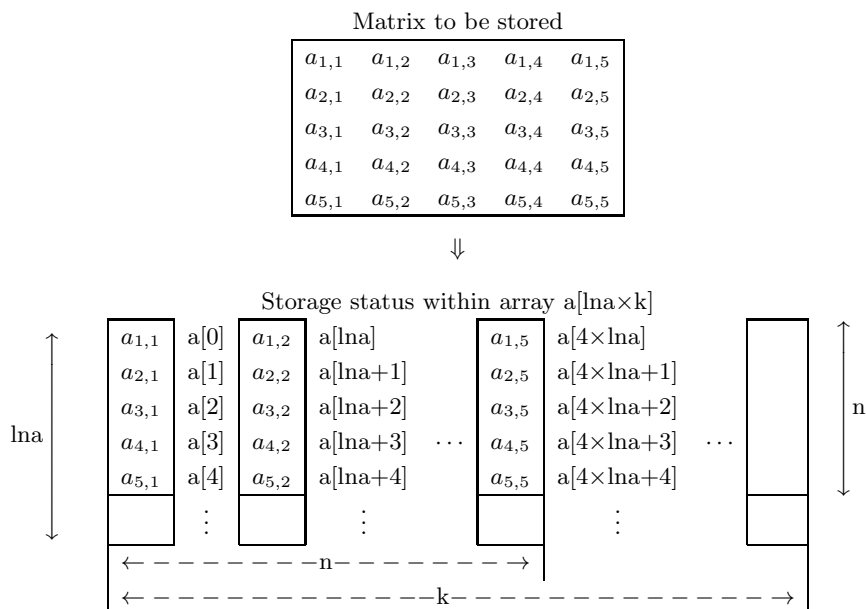
Figure A-2 Matrix Storage Mode Within an Array  $a[lma \times k]$

If you will be manipulating several arrays having different orders as data, you can prepare one array having lna equal to the largest order and use that array successively for each array. However, you must always assign the lna value as an adjustable dimension.

## A.2 Data storage modes

Matrix data storage modes differ according to the matrix type. Storage modes for each type of matrix are shown below.

### A.2.1 Real matrix (two-dimensional array type)



#### Remarks

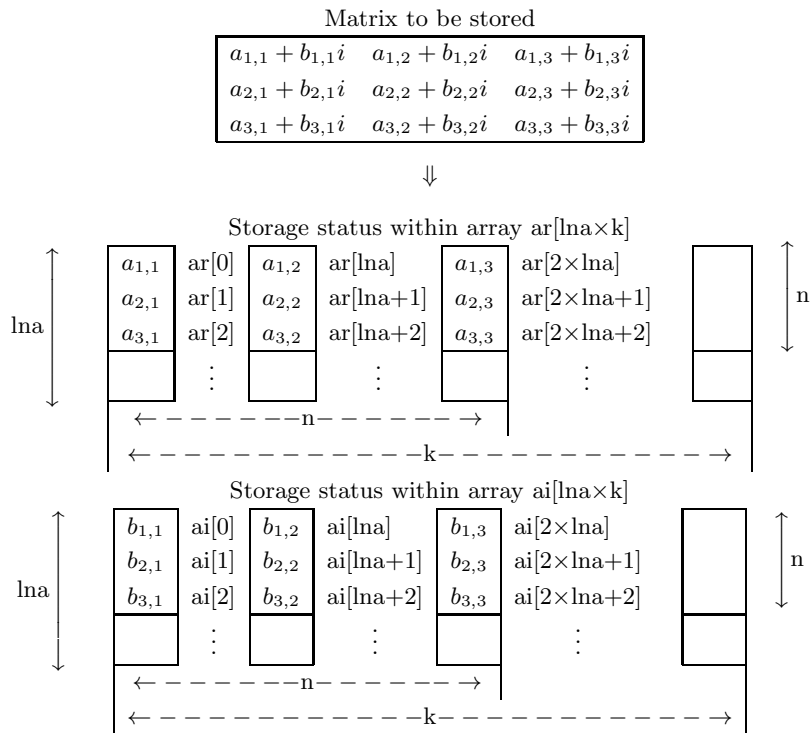
- a.  $lna \geq n$  and  $k \geq n$  must hold.

Figure A-3 Real Matrix (Two-Dimensional Array Type) Storage Mode

### A.2.2 Complex matrix

(1) **Two-dimensional array type, real argument type**

Real and imaginary parts are stored in separate arrays.

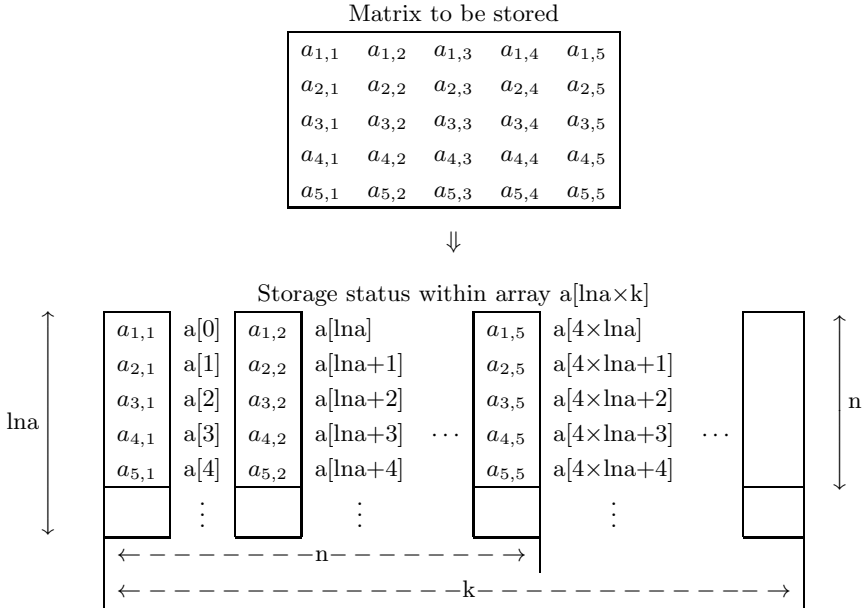


**Remarks**

- a.  $l_{na} \geq n$  and  $k \geq n$  must hold.

Figure A-4 Complex Matrix (Two-dimensional Array Type) (Real Argument Type) Storage Mode

(2) Two-dimensional array type, complex argument type

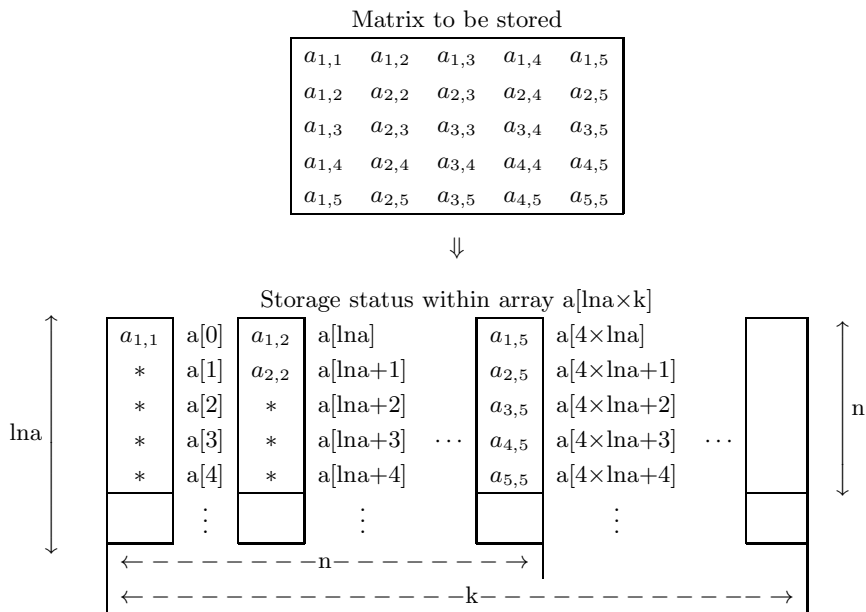


**Remarks**  
 a.  $lna \geq n$  and  $k \geq n$  must hold.

Figure A-5 Complex Matrix (Two-dimensional Array Type)(Complex Argument Type) Storage Mode

### A.2.3 Real symmetric matrix and positive symmetric matrix

(1) Two-dimensional array type, upper triangular type



**Remarks**

- a. The asterisk (\*) indicates an arbitrary value.
- b.  $lna \geq n$  and  $k \geq n$  must hold.

Figure A-6 Real Symmetric Matrix (Two-dimensional Array Type) (Upper Triangular Type) Storage mode

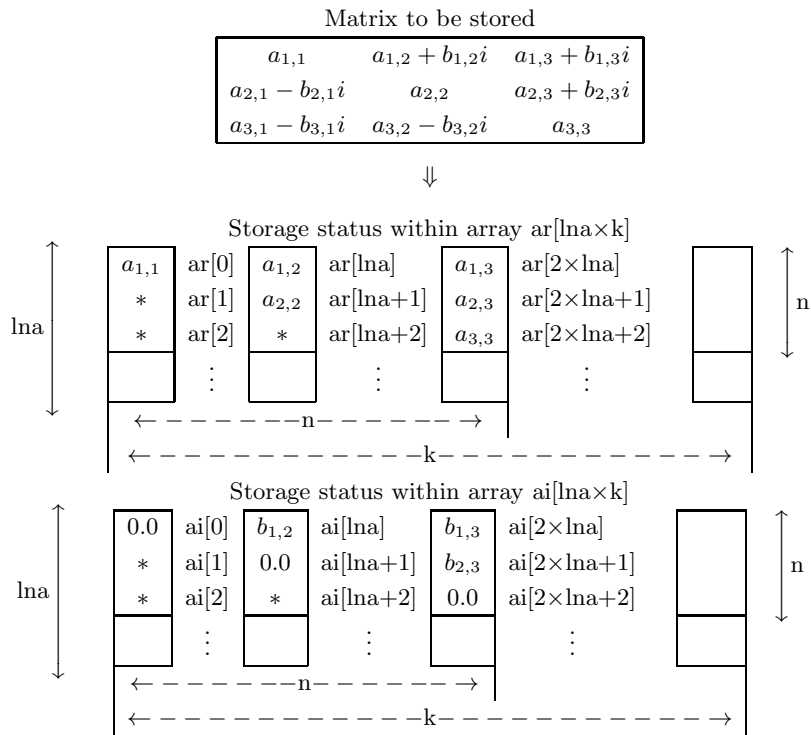




### A.2.4 Hermitian matrix

(1) **Two-dimensional array type, real argument type, upper triangular type**

Upper triangular portions of the real and imaginary parts are stored in separate arrays.

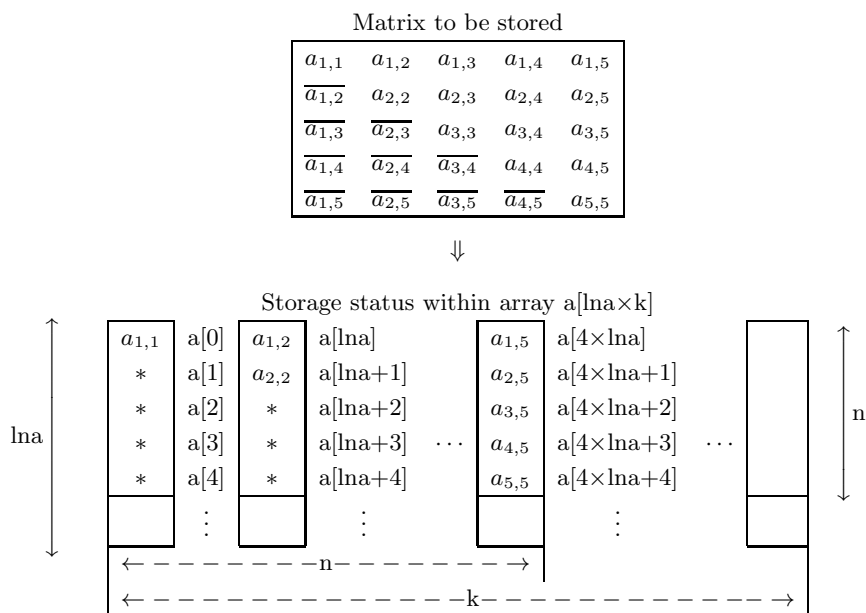


**Remarks**

- a. The asterisk (\*) indicates an arbitrary value.
- b.  $l_{na} \geq n$  and  $k \geq n$  must hold.

Figure A–8 Hermitian Matrix (Two-dimensional Array Type) (Real Argument Type) (Upper Triangular Type) Storage Mode

(2) Two-dimensional array type, complex argument type, upper triangular type



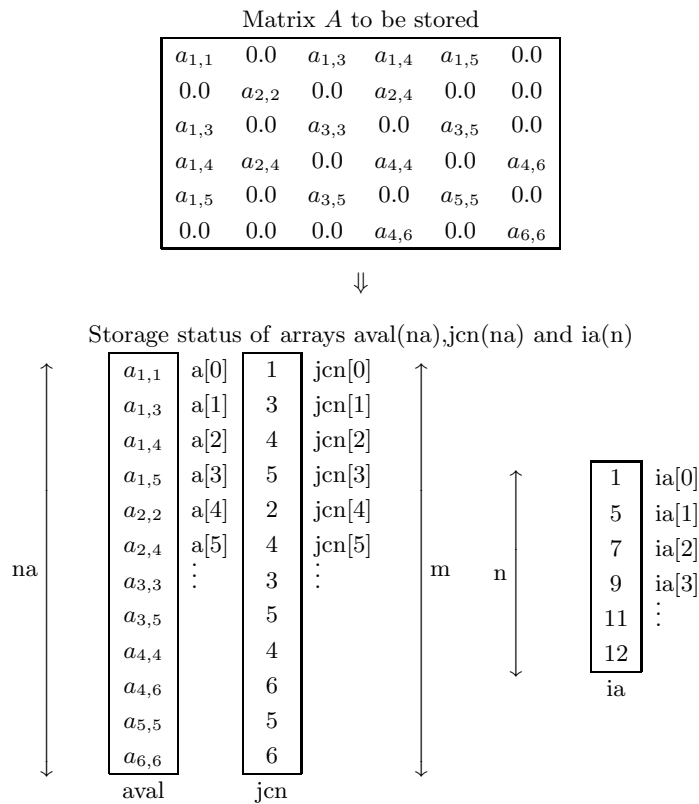
**Remarks**

- a. The  $\overline{x}$  indicates the complex conjugate of  $x$ .
- b. The asterisk  $*$  indicates an arbitrary value.
- c.  $lna \geq n$  and  $k \geq n$  must hold.

Figure A-9 Hermitian Matrix (Two-dimensional Array Type) (Complex Argument Type) (Upper Triangular Type) Storage Mode

### A.2.5 Random sparse matrix (For symmetric matrix only)

(1) Sparse format (Symmetric case)



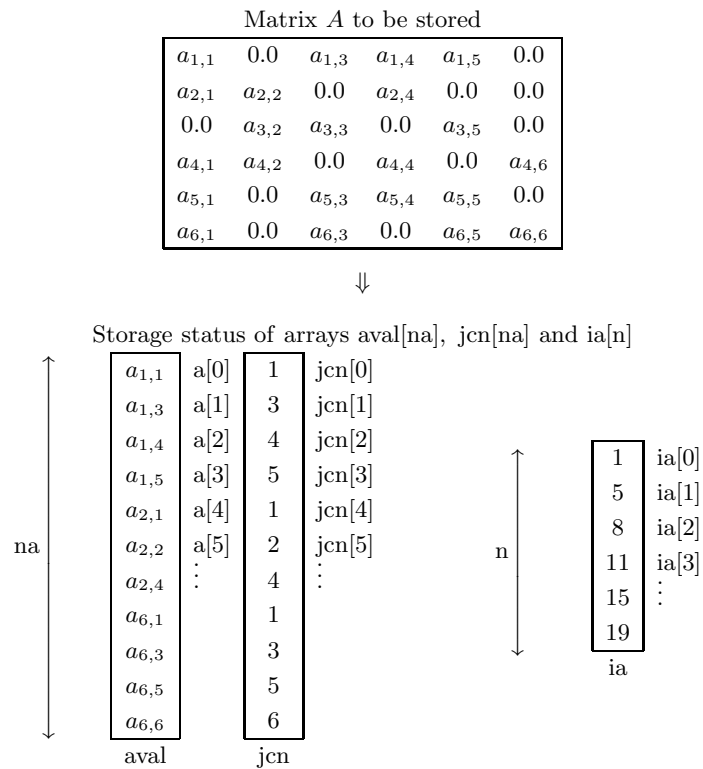
**Remarks**

- a.  $m$  is the number of nonzero elements in the upper triangular part of the original matrix  $A$  including the diagonal.
- b. Array  $aval$  contains the nonzero upper triangular elements of the original matrix  $A$ , stored sequentially beginning with the first row.
- c. Array  $jcn$  contains the column numbers in the original matrix  $A$  of the elements stored in array  $aval$ .
- d. Array  $ia$  contains values equal to one added to the positions in array  $aval$  of the diagonal elements.
- e.  $n \leq m < na$  must hold.

Figure A–10 Storage of Random Symmetric Sparse Matrix (Sparse format)

### A.2.6 Random sparse matrix

(1) Sparse format



**Remarks**

- a.  $na$  is the number of nonzero elements of the original matrix  $A$ .
- b. Array  $aval$  contains the nonzero elements of the original matrix  $A$ , stored sequentially beginning with the first row.
- c. Array  $jcn$  contains the column indices in the original matrix  $A$  of the elements stored in array  $aval$ .
- d. Array  $ia$  contains values equal to one added to the positions in array  $aval$  of the first nonzero element in each row.
- e.  $n < na$  must hold.

Figure A–11 Storage of Real Asymmetric Random Sparse Matrix (Sparse Format)

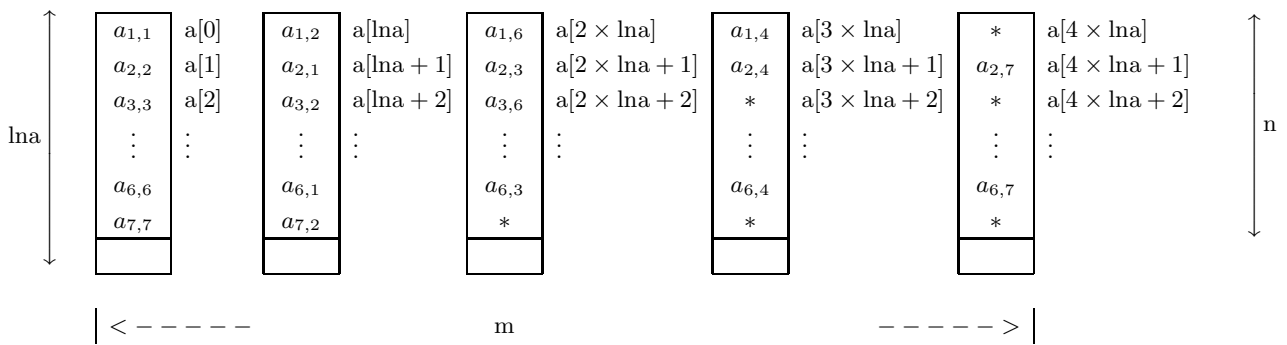
(2) ELLPACK format

Matrix  $A$  to be stored

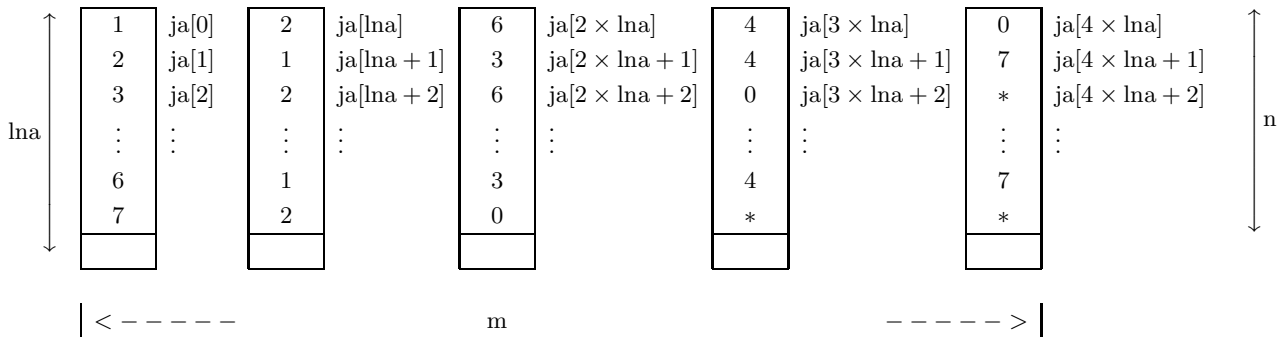
$a_{1,1}$	$a_{1,2}$	0.0	$a_{1,4}$	0.0	$a_{1,6}$	0.0
$a_{2,1}$	$a_{2,2}$	$a_{2,3}$	$a_{2,4}$	0.0	0.0	$a_{2,7}$
0.0	$a_{3,2}$	$a_{3,3}$	0.0	0.0	$a_{3,6}$	0.0
$a_{4,1}$	$a_{4,2}$	0.0	$a_{4,4}$	$a_{4,5}$	$a_{4,6}$	0.0
0.0	0.0	0.0	$a_{5,4}$	$a_{5,5}$	0.0	0.0
$a_{6,1}$	0.0	$a_{6,3}$	$a_{6,4}$	0.0	$a_{6,6}$	$a_{6,7}$
0.0	$a_{7,2}$	0.0	0.0	0.0	0.0	$a_{7,7}$

↓

Storage status of array  $a[lna \times m]$



Storage status of array  $ja[l na \times m]$



**Remarks**

- $n$  is order of Matrix  $A$ .
- $lna \geq n$  must hold.
- $m$  is the column number of Array  $a$ , which contains the nonzero elements of Matrix  $A$ .
- Array  $a$  should contain nonzero elements of Matrix  $A$  so that:
  - Diagonal elements are stored in the first column.
  - Nonzero elements in the lower triangular part and the upper triangular part are stored in the second through  $m$ -th columns, with the first one in the second column, one adjacent to the next in each row. Here, it is unnecessary that nonzero elements in each row are stored sequentially.
  - Arbitrary values can be stored in the remaining positions that are marked with '\*'.
- Array  $ja$  should contain the column indices in Matrix  $A$  of those elements that correspond to the elements contained in Array  $a$ .

For those rows in which  $m - 1$  becomes greater than the number of nonzero elements in the lower and upper triangular part, value 0 should be stored in the right neighbor of the rightmost position of the region

in `ja` in which the column indices of nonzero elements in Matrix  $A$  are stored. Arbitrary values can be stored in the remaining positions that are marked with `'*'`.

Figure A–12 Storage format for Asymmetric Random Sparse Matrix (ELLPACK Format)

## Appendix B

# MACHINE CONSTANTS USED IN ASL C INTERFACE

### B.1 Units for Determining Error

The table below shows values in ASL C interface as units for determining error in floating point calculations. The units shown in the table are numeric values determined by the internal representation of floating point data. ASL C interface uses these units for determining convergence and zeros.

Table B–1 Units for Determining Error

Single-precision	Double-precision
$2^{-23} (\simeq 1.19 \times 10^{-7})$	$2^{-52} (\simeq 2.22 \times 10^{-16})$

**Remark:** The unit for determining error  $\varepsilon$ , which is also called the machine  $\varepsilon$ , is usually defined as the smallest positive constant for which the calculation result of  $1 + \varepsilon$  differs from 1 in the corresponding floating point mode. Therefore, seeing the unit for determining error enables you to know the maximum number of significant digits of an operation (on the mantissa) in that floating point mode.

### B.2 Maximum and Minimum Values of Floating Point Data

The table below shows maximum and minimum values of floating point data defined within ASL C interface. Note that the maximum and minimum values shown below may differ from the maximum and minimum values that are actually used by the hardware for each floating point mode.

Table B–2 Maximum and Minimum Values of Floating Point Data

	Single-precision	Double-precision
Maximum value	$2^{127}(2 - 2^{-23}) (\simeq 3.40 \times 10^{38})$	$2^{1023}(2 - 2^{-52}) (\simeq 1.80 \times 10^{308})$
Positive minimum value	$2^{-126} (\simeq 1.17 \times 10^{-38})$	$2^{-1022} (\simeq 2.23 \times 10^{-308})$
Negative maximum value	$-2^{-126} (\simeq -1.17 \times 10^{-38})$	$-2^{-1022} (\simeq -2.23 \times 10^{-308})$
Minimum value	$-2^{127}(2 - 2^{-23}) (\simeq -3.40 \times 10^{38})$	$-2^{1023}(2 - 2^{-52}) (\simeq -1.80 \times 10^{308})$

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